
Subject: Re: [ckrm-tech] [RFC] [PATCH 0/3] Add group fairness to CFS
Posted by [William Lee Irwin III](#) on Thu, 31 May 2007 08:43:29 GMT
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On Wed, May 30, 2007 at 11:36:47PM -0700, William Lee Irwin III wrote:
>> Temporarily, yes. All this only works when averaged out.

On Thu, May 31, 2007 at 02:03:53PM +0530, Srivatsa Vaddagiri wrote:
> So essentially when we calculate delta_mine component for each of those
> 1000 tasks, we will find that it has executed for 1 tick (4 ms say) but
> its fair share was very very low.
> $\text{fair_share} = \text{delta_exec} * p \rightarrow \text{load_weight} / \text{total_weight}$
> If $p \rightarrow \text{load_weight}$ has been calculated after factoring in hierarchy (as
> you outlined in a previous mail), then $p \rightarrow \text{load_weight}$ of those 1000 tasks
> will be far less compared to the $p \rightarrow \text{load_weight}$ of one task belonging to
> other user, correct? Just to make sure I get all this correct:

You've got it all correct.

On Thu, May 31, 2007 at 02:03:53PM +0530, Srivatsa Vaddagiri wrote:
> User U1 has tasks T0 - T999
> User U2 has task T1000
> assuming each task's weight is 1 and each user's weight is 1 then:
> $WT0 = (WU1 / WU1 + WU2) * (WT0 / WT0 + WT1 + \dots + WT999)$
> $= (1 / 1 + 1) * (1 / 1000)$
> $= 1/2000$
> $= 0.0005$
> WT1 ..WT999 will be same as WT0
> whereas, weight of T1000 will be:
> $WT1000 = (WU1 / WU1 + WU2) * (WT1000 / WT1000)$
> $= (1 / 1 + 1) * (1/1)$
> $= 0.5$
> ?

Yes, these calculations are correct.

On Thu, May 31, 2007 at 02:03:53PM +0530, Srivatsa Vaddagiri wrote:
> So when T0 (or T1 ..T999) executes for 1 tick (4ms), their fair share would
> be:
> T0's fair_share (delta_mine)
> $= 4 \text{ ms} * 0.0005 / (0.0005 * 1000 + 0.5)$
> $= 4 \text{ ms} * 0.0005 / 1$
> $= 0.002 \text{ ms} (2000 \text{ ns})$
> This would cause T0's $\rightarrow \text{wait_runtime}$ to go negative sharply, causing it to be
> inserted back in rb-tree well ahead in future. One change I can foresee
> in CFS is with regard to `limit_wait_runtime()` ..We will have to change

> its default limit, atleast when group fairness thingy is enabled.
> Compared to this when T1000 executes for 1 tick, its fair share would be
> calculated as:
> T1000's fair_share (delta_mine)
> = $4 \text{ ms} * 0.5 / (0.0005 * 1000 + 0.5)$
> = $4 \text{ ms} * 0.5 / 1$
> = 2 ms (2000000 ns)
> Its ->wait_runtime will drop less significantly, which lets it be
> inserted in rb-tree much to the left of those 1000 tasks (and which indirectly
> lets it gain back its fair share during subsequent schedule cycles).

This analysis is again entirely correct.

On Thu, May 31, 2007 at 02:03:53PM +0530, Srivatsa Vaddagiri wrote:

> Hmm ..is that the theory?
> Ingo, do you have any comments on this approach?
> /me is tempted to try this all out.

Yes, this is the theory behind using task weights to flatten the task group hierarchies. My prior post assumed all this and described a method to make nice numbers behave as expected in the global context atop it.

-- wli

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