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Subject: Re: [patch -mm 08/17] nsproxy: add hashtable  
Posted by [Herbert Poetzl](#) on Tue, 12 Dec 2006 23:22:22 GMT  
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On Mon, Dec 11, 2006 at 04:01:15PM -0600, Serge E. Hallyn wrote:  
> Quoting Eric W. Biederman (ebiederm@xmission.com):  
> > "Serge E. Hallyn" <serue@us.ibm.com> writes:  
> >  
> > > Quoting Eric W. Biederman (ebiederm@xmission.com):  
> > >  
> > > Yeah, that occurred to me, but it doesn't seem like we can possibly make  
> > > sufficient guarantees to the client to make this worthwhile.  
> > >  
> > > I'd love to be wrong about that, but if nothing else we can't prove to  
> > > the client that they're running on an unhacked host. So the host admin  
> > > will always have to be trusted.  
> >  
> > To some extent that is true. Although all security models we have  
> > currently fall down if you hack the kernel, or run your kernel  
> > in a hacked virtual environment. It would be nice if under normal  
> > conditions you could mount an encrypted filesystem only in a container  
> > and not have concerns of those files escaping.  
>  
> Hmm, well perhaps I'm being overly pessimistic - IBM research did have a  
> demo based on TPM of remote attestation, which may be usable for  
> ensuring that you're connecting to a service on your virtual machine on  
> a certain (unhacked) kernel on particular hardware, in which case what  
> you're talking about may be possible - given a stringent initial  
> environment (i.e. not the 'gimme \$20/month for a hosted partition in  
> arizona' environment).

interesting, how would you ensure from inside  
such an environment, that nobody tampered with  
the kernel you are running on?

> Given that, perhaps having a virtual machine with access to encrypted  
> storage - safe from the host machine admins - may not be unattainable  
> after all. And given that, it would be worth designing the ns\_enter()  
> system call so that a parent cannot enter some child namespace.

we currently call this Context Privacy, and it  
is partially implemented, but of course, it  
does only work if the kernel is known good

> > Which would probably be a matter of having a separate uid\_ns and not  
> > allowing process outside of your container to have any permissions in  
> > that filesystem.  
>

> Yup. Or even just a separate uid\_ns and an ecryptfs partition, so  
> that the host can back up the encrypted data incrementally (per file,  
> i.e. not just the whole dmccrypted loop file).

it's simple to avoid access to certain 'tagged'  
devices and/or filesystems, it's hard to handle  
kernel modifications or even simple things like  
reading the kernel memory ...

best,  
Herbert

> -serge

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Containers mailing list  
Containers@lists.osdl.org  
<https://lists.osdl.org/mailman/listinfo/containers>

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