Subject: Re: MCR
Posted by Daniel Lezcano on Thu, 14 Dec 2006 10:01:42 GMT
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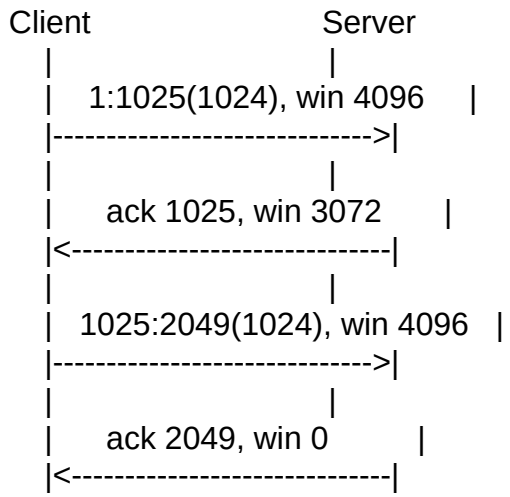
Masahiko Takahashi wrote:
> Thank you Daniel. I'm now getting much understanding about
> what MCR does. Then, the next step you are going to do is
> to implement sk_filter and check its performance ?

Yes after network isolation is finished.

>
> On Wed, 2006-12-13 at 11:34 +0100, Daniel Lezcano wrote:
>> There are 2 aspects:
>> - TCP buffer
>> - TCP re-transmission
>>
>> When you checkpoint, the outgoing packets are dropped when they are sent
>> but the buffer is still there until the packets are acked by the peer
>> (and that never happens because the traffic is dropped). These buffers
>> are checkpointed. When you restart, you restore these buffers, you
>> release the traffic and the TCP layer continue to send the packets. It
>> is like you unplugged the network cable, wait a little and plugged it,
>> the TCP traffic is resumed. We rely on the TCP mechanisms.
> :
>> In the TCP communication process, the receiver send the remaining opened
>> windows he has. The sender rely on that to send the nb bytes
>> corresponding. If you send a zero window while the sender is expecting
>> to have at least "last windows - nb bytes sent", depending on the TCP
>> implementation, nothing can happen (traffic will block), sender can
>> decide to drop the connection because it thinks it is inconsistent or
>> you can have other behaviors like retransmissions ...
>
> I understand MCR relies on TCP's robustness and zero-window
> advertisement has a delicate issue, but I, as a network
> amateur, still couldn't believe zero-window's inconsistency
> is a serious problem...

Both client and server have a 4096 window size. The server application
is blocked (control+Z, data processing, stucked, ...) so it never reads
the incoming traffic.
  1. Client send 1024 bytes to server
  2. Server ack these data and the window is 3072 (4096-1024)
  3. Client send 1024 bytes again to server
  4. Client ack these data but with a window advertisement of zero,
normally this one should be between 2048 and 4096, no less than 2048.



Perhaps, it could work, perhaps not. To use the zero window we should
check that all TCP stacks are compatible with that. If we have a server
running on linux and we checkpoint it, it will not be cool if all
windows client lose their connecions while all linux client or mac
clients stay blocked.

>
>> When an application creates a socket, write to it and close it, the
>> connection stay alive until all data are sent and wait for a moment in a
>> specific tcp state (LAST_ACK, CLOSE_WAIT, TIME_WAIT, ...). After a time
>> the socket is destroyed, freezing the sockets timer will avoid the
>> socket to be destroyed. Do netstat -tano command...
>
> It seems we should check there is no packet in receive_queue
> before closing the socket. If there is, the kernel tries to
> send a RESET packet.

The receive and the send queue are checkpointed. If the receive queue is
not flushed before destroying the socket, there is no issue because the
RST packet will be dropped because of the blocked traffic.

Regards.

-- Daniel

Subject: TCP checkpoint/restart (Re: MCR)
Posted by Cedric Le Goater on Thu, 14 Dec 2006 12:57:29 GMT
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Thanks Daniel for moving that thread on the containers@ list.

When you have some time, could you just recap the main topics
of this discussion on tcp stack checkpoint/restart. I'm pretty
sure the openvz team as plenty to say.

Thanks,

C.

Daniel Lezcano wrote:
> Masahiko Takahashi wrote:
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> 1. Client send 1024 bytes to server
> 2. Server ack these data and the window is 3072 (4096-1024)
> 3. Client send 1024 bytes again to server
> 4. Client ack these data but with a window advertisement of zero,
> normally this one should be between 2048 and 4096, no less than 2048.
>
>
> Client   Server
> | |
> | 1:1025(1024), win 4096 |
> |--------------------------->|
> | |
> | ack 1025, win 3072 |
> |<---------------------------|
> | |
> | 1025:2049(1024), win 4096 |
> |--------------------------->|
> | |
> | ack 2049, win 0 |
> |<---------------------------|
>
>
> Perhaps, it could work, perhaps not. To use the zero window we should
> check that all TCP stacks are compatible with that. If we have a server
> running on linux and we checkpoint it, it will not be cool if all
> windows client lose their connecions while all linux client or mac
> clients stay blocked.
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>>> connection stay alive until all data are sent and wait for a moment in a
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>>> the socket is destroyed, freezing the sockets timer will avoid the
>>> socket to be destroyed. Do netstat -tano command...
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> RST packet will be dropped because of the blocked traffic.
>
> Regards.
>

> -- Daniel
>
>
>
>

_____
Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: MCR
Posted by Masahiko Takahashi on Fri, 15 Dec 2006 00:56:10 GMT
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Hi folks,

On Thu, 2006-12-14 at 11:01 +0100, Daniel Lezcano wrote:
> Perhaps, it could work, perhaps not. To use the zero window we should
> check that all TCP stacks are compatible with that. If we have a server
> running on linux and we checkpoint it, it will not be cool if all
> windows client lose their connecions while all linux client or mac
> clients stay blocked.

I understand what you mean. We should take a conservative
way, even if dropping packets results in shrinking the
window size.

> > It seems we should check there is no packet in receive_queue
> > before closing the socket. If there is, the kernel tries to
> > send a RESET packet.
>
> The receive and the send queue are checkpointed. If the receive queue is
> not flushed before destroying the socket, there is no issue because the
> RST packet will be dropped because of the blocked traffic.

You are right. I forgot that all packets would be dropped
after checkpointing in MCR.

Thanks,

Masahiko.

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Containers mailing list

Subject: Re: TCP checkpoint/restart (Re: MCR)
Posted by Daniel Lezcano on Fri, 15 Dec 2006 10:56:45 GMT
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Cedric Le Goater wrote:
> Thanks Daniel for moving that thread on the containers@ list.
>
> When you have some time, could you just recap the main topics
> of this discussion on tcp stack checkpoint/restart. I'm pretty
> sure the openvz team as plenty to say.

Sure.

Actually we are working on the network isolation. There are 2 aspects:

* Full network isolation/virtualization acting at the layer 2 (device)
* Network isolation at IP layer, we call it layer 3

The network isolation is the mandatory mechanism to ensure the checkpoint/restart because we must identify the network ressources associated to a container and avoid these ressources to overlap with other containers.

To be able to take a snapshot of the network container, we must ensure it is freezed during the checkpoint, because we must ensure the consistency in the host and with the peers network stack.

We began the checkpoint/restart discussion with this point: how do we do container's network freeze ?

* The first step is to drop the traffic
  - shall it be done with the sk_filter fields of the socket ?
  - or with the netfilter NF_DROP/NF_STOLEN ?

* The second step is to stop tcp timers to avoid socket destruction while checkpointing it

-- Daniel

_____
Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: TCP checkpoint/restart (Re: MCR)
Posted by Masahiko Takahashi on Tue, 30 Jan 2007 00:54:43 GMT
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Hi Daniel,

On Fri, 2006-12-15 at 11:56 +0100, Daniel Lezcano wrote:
> To be able to take a snapshot of the network container, we must ensure
> it is freezed during the checkpoint, because we must ensure the
> consistency in the host and with the peers network stack.
>
> We began the checkpoint/restart discussion with this point: how do we do
> container's network freeze ?
>
> * The first step is to drop the traffic
> - shall it be done with the sk_filter fields of the socket ?
> - or with the netfilter NF_DROP/NF_STOLEN ?

I've tested with a very simple sk_filter and found that sk_filter
cannot drop outgoing packets whereas it can drop incoming packets.
Therefore, if the implementation is carefully done so as not to
send outgoing ones when checkpointing, it can freeze container's
network. But if there may be some accidental packet sending, maybe
it fails to freeze the network. Or, we had better implement
sk_filter for outgoing packets.

The following function is my simple sk_filter set in user level.
I have tested for both TCP and UDP socket.

Thanks,

Masahiko.

```
void filter_socket ( int fd )
{
        struct sock_fprog prg;
        struct sock_filter skf[1] = { 0, };
```

```
        prg.filter = skf;

        skf[0].code = BPF_RET|BPF_A;
        prg.len = 1;

        if ( setsockopt ( fd, SOL_SOCKET, SO_ATTACH_FILTER,
                        (void *) &prg, sizeof(prg)) < 0 )
                perror( "setsockopt SO_ATTACH_FILTER" );
}
```

_____

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