## Subject: [patch -mm 00/17] new namespaces and related syscalls Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:27:52 GMT

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All,

Here's a patch serie we've been working on for a while. I'd like to know which could make it to -mm and which need refinements. Depending on the feedback, i'll send them on lkml@.

It feels like that the empty net namespace framework could make it there already.

## Patch menu:

\* start with user and net namespace. no real feature and no unshare. mostly solves conflicts in nsproxy which are becoming painful. tasty and harmless

[patch 01/17] net namespace: empty framework [patch 02/17] user namespace: add the framework

\* follow some cleanups and code movement preparing the next drop. less tasty but still harmless.

[patch 03/17] namespace: export unshare of namespace and fs\_struct [patch 04/17] nsproxy: externalizes exit\_task\_namespaces [patch 05/17] ipc namespace: externalizes unshare\_ipcs [patch 06/17] nsproxy: add extern to nsproxy functions [patch 07/17] nsproxy: make put\_nsproxy an extern

\* nsproxy syscalls. real meat.

[patch 08/17] nsproxy: add hashtable

[patch 09/17] nsproxy: add namespace flags

[patch 10/17] nsproxy: add unshare ns and bind ns syscalls

\* use the user namespace for real. wine and cheese

[patch 11/17] user namespace: add user\_namespace ptr to vfsmount

[patch 12/17] user namespace: hook permission

[patch 13/17] user namespace: implement shared mounts

[patch 14/17] user namespace: maintain user ns for priv\_userns mounts to vfsmount

\* finish with the stuff that make a real meal: unsharing namespaces

dessert without sugar :) [patch 15/17] pid namespace: add unshare [patch 16/17] net namespace: add unshare [patch 17/17] user namespace: add unshare Cheers, C. Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers Subject: [patch -mm 01/17] net namespace: empty framework Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:27:53 GMT View Forum Message <> Reply to Message From: Cedric Le Goater <clg@fr.ibm.com> This patch adds an empty net namespace framework with an API compatible with the other available namespaces. It doesn't provide any feature by itself but prepares the ground for work on L2 and L3 network isolation. It also solves patch conflict issues in nsproxy, which are painful and boring to resolve. Signed-off-by: Cedric Le Goater <clg@fr.ibm.com> include/linux/init task.h 2+ include/linux/nsproxy.h 2 +

some will surely find these less in sync with their taste, like

Index: 2.6.19-rc6-mm2/include/linux/init\_task.h

```
--- 2.6.19-rc6-mm2.orig/include/linux/init_task.h
+++ 2.6.19-rc6-mm2/include/linux/init_task.h
@ @ -8,6 +8,7 @ @
#include <linux/lockdep.h>
#include <linux/ipc.h>
```

```
#include linux/pid_namespace.h>
+#include linux/net namespace.h>
#define INIT_FDTABLE \
@@ -78,6 +79,7 @@ extern struct nsproxy init_nsproxy;
 .id = 0,
 .uts_ns = &init_uts_ns,
 .mnt ns = NULL,
+ INIT NET NS(net ns)
 INIT_IPC_NS(ipc_ns)
Index: 2.6.19-rc6-mm2/include/linux/net_namespace.h
--- /dev/null
+++ 2.6.19-rc6-mm2/include/linux/net namespace.h
@ @ -0,0 +1,49 @ @
+#ifndef LINUX NET NAMESPACE H
+#define LINUX NET NAMESPACE H
+#include ux/kref.h>
+#include linux/nsproxy.h>
+struct net_namespace {
+ struct kref kref;
+};
+extern struct net namespace init net ns;
+#ifdef CONFIG NET NS
+#define INIT_NET_NS(net_ns) .net_ns = &init_net_ns,
+static inline void get_net_ns(struct net_namespace *ns)
+ kref_get(&ns->kref);
+}
+extern int copy_net_ns(int flags, struct task_struct *tsk);
+extern void free_net_ns(struct kref *kref);
+static inline void put_net_ns(struct net_namespace *ns)
+ kref_put(&ns->kref, free_net_ns);
+}
```

```
+#else
+#define INIT_NET_NS(net_ns)
+static inline void get_net_ns(struct net_namespace *ns)
+{
+}
+static inline int copy_net_ns(int flags, struct task_struct *tsk)
+{
+ return 0;
+}
+static inline void put_net_ns(struct net_namespace *ns)
+{
+}
+#endif
+#endif /* LINUX NET NAMESPACE H */
Index: 2.6.19-rc6-mm2/include/linux/nsproxy.h
--- 2.6.19-rc6-mm2.orig/include/linux/nsproxy.h
+++ 2.6.19-rc6-mm2/include/linux/nsproxy.h
@ @ -8,6 +8,7 @ @ struct mnt_namespace;
struct uts_namespace;
struct ipc_namespace;
struct pid_namespace;
+struct net_namespace;
 * A structure to contain pointers to all per-process
@ @ -29,6 +30,7 @ @ struct nsproxy {
 struct ipc_namespace *ipc_ns;
 struct mnt_namespace *mnt_ns;
 struct pid_namespace *pid_ns;
+ struct net_namespace *net_ns;
extern struct nsproxy init_nsproxy;
Index: 2.6.19-rc6-mm2/kernel/nsproxy.c
--- 2.6.19-rc6-mm2.orig/kernel/nsproxy.c
+++ 2.6.19-rc6-mm2/kernel/nsproxy.c
@ @ -20,6 +20,7 @ @
#include linux/mnt_namespace.h>
#include linux/utsname.h>
#include linux/pid namespace.h>
+#include linux/net namespace.h>
```

```
struct nsproxy init_nsproxy = INIT_NSPROXY(init_nsproxy);
@@ -71,6 +72,8 @@ struct nsproxy *dup_namespaces(struct ns
  get_ipc_ns(ns->ipc_ns);
 if (ns->pid_ns)
  get_pid_ns(ns->pid_ns);
+ if (ns->net_ns)
+ get_net_ns(ns->net_ns);
 return ns:
@ @ -118,10 +121,17 @ @ int copy_namespaces(int flags, struct ta
 if (err)
 goto out_pid;
+ err = copy_net_ns(flags, tsk);
+ if (err)
+ goto out_net;
out:
 put_nsproxy(old_ns);
 return err;
+out_net:
+ if (new_ns->pid_ns)
+ put_pid_ns(new_ns->pid_ns);
out pid:
if (new ns->ipc ns)
 put_ipc_ns(new_ns->ipc_ns);
@ @ -147,5 +157,7 @ @ void free nsproxy(struct nsproxy *ns)
 put_ipc_ns(ns->ipc_ns);
 if (ns->pid_ns)
 put_pid_ns(ns->pid_ns);
+ if (ns->net_ns)
+ put_net_ns(ns->net_ns);
 kfree(ns);
}
Index: 2.6.19-rc6-mm2/net/core/Makefile
--- 2.6.19-rc6-mm2.orig/net/core/Makefile
+++ 2.6.19-rc6-mm2/net/core/Makefile
@@ -3,7 +3,7 @@
#
obj-y := sock.o request_sock.o skbuff.o iovec.o datagram.o stream.o scm.o \
- gen stats.o gen estimator.o
+ gen stats.o gen estimator.o net namespace.o
```

Index: 2.6.19-rc6-mm2/net/core/net\_namespace.c

--- /dev/null +++ 2.6.19-rc6-mm2/net/core/net\_namespace.c @ @ -0,0 +1,41 @ @ +/\* + \* This program is free software; you can redistribute it and/or + \* modify it under the terms of the GNU General Public License as + \* published by the Free Software Foundation, version 2 of the + \* License. + \*/ +#include linux/module.h> +#include linux/version.h> +#include linux/nsproxy.h> +#include linux/net namespace.h> +struct net\_namespace init\_net\_ns = { + .kref = { + .refcount = ATOMIC\_INIT(2), + }, +}; +#ifdef CONFIG\_NET\_NS +int copy net ns(int flags, struct task struct \*tsk) + struct net namespace \*old ns = tsk->nsproxy->net ns; + int err = 0; + if (!old\_ns) + return 0; + get\_net\_ns(old\_ns); + return err; +} +void free net ns(struct kref \*kref) +{ + struct net\_namespace \*ns; + ns = container\_of(kref, struct net\_namespace, kref); + kfree(ns); +}

+#endif /\* CONFIG\_NET\_NS \*/
Index: 2.6.19-rc6-mm2/net/Kconfig

\_\_\_\_\_\_

--- 2.6.19-rc6-mm2.orig/net/Kconfig +++ 2.6.19-rc6-mm2/net/Kconfig @@ -67,6 +67,14 @@ source "net/netlabel/Kconfig"

endif # if INET

- +config NET NS
- + bool "Network Namespaces"
- + help
- + This option enables multiple independent network namespaces,
- + each having own network devices, IP addresses, routes, and so on.
- + If unsure, answer N.

+

config NETWORK\_SECMARK bool "Security Marking" help

--

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: [patch -mm 02/17] user namespace: add the framework Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:27:54 GMT View Forum Message <> Reply to Message

From: Cedric Le Goater <clg@fr.ibm.com>

This patch adds the user namespace struct and framework

Basically, it will allow a process to unshare its user\_struct table, resetting at the same time its own user\_struct and all the associated accounting.

A new root user (uid == 0) is added to the user namespace upon creation. Such root users have full privileges and it seems that theses privileges should be controlled through some means (process capabilities?)

The unshare is not included in this patch.

Changes since [try #3]:

```
- moved struct user namespace to files user namespace.{c,h}
Changes since [try #2]:

    removed struct user_namespace* argument from find_user()

Changes since [try #1]:
- removed struct user namespace* argument from find user()
- added a root_user per user namespace
Signed-off-by: Cedric Le Goater <clg@fr.ibm.com>
include/linux/init_task.h
                          2 +
                       | 2+
include/linux/nsproxy.h
include/linux/sched.h
                       | 3+-
init/Kconfig
                    8 +++++
kernel/Makefile
                      | 2 -
kernel/fork.c
                    | 2 -
kernel/nsproxy.c
                      | 11 +++++++
kernel/sys.c
                       5 ++--
kernel/user.c
                       18 ++++++
kernel/user namespace.c
                          11 files changed, 132 insertions(+), 14 deletions(-)
Index: 2.6.19-rc6-mm2/kernel/user.c
--- 2.6.19-rc6-mm2.oria/kernel/user.c
+++ 2.6.19-rc6-mm2/kernel/user.c
@@ -14,20 +14,19 @@
#include linux/bitops.h>
#include linux/key.h>
#include linux/interrupt.h>
+#include linux/module.h>
+#include linux/user_namespace.h>
 * UID task count cache, to get fast user lookup in "alloc uid"
 * when changing user ID's (ie setuid() and friends).
-#define UIDHASH_BITS (CONFIG_BASE_SMALL ? 3 : 8)
-#define UIDHASH_SZ (1 << UIDHASH_BITS)
#define UIDHASH_MASK (UIDHASH_SZ - 1)
#define __uidhashfn(uid) (((uid >> UIDHASH_BITS) + uid) & UIDHASH_MASK)
-#define uidhashentry(uid) (uidhash_table + __uidhashfn((uid)))
+#define uidhashentry(ns, uid) ((ns)->uidhash table + uidhashfn((uid)))
```

```
static kmem cache t *uid cachep;
-static struct list head uidhash table[UIDHASH SZ];
/*
 * The uidhash lock is mostly taken from process context, but it is
@ @ -94,9 +93,10 @ @ struct user_struct *find_user(uid_t uid)
{
 struct user_struct *ret;
 unsigned long flags;
+ struct user namespace *ns = current->nsproxy->user ns;
 spin lock irgsave(&uidhash lock, flags);
- ret = uid_hash_find(uid, uidhashentry(uid));
+ ret = uid_hash_find(uid, uidhashentry(ns, uid));
 spin_unlock_irgrestore(&uidhash_lock, flags);
 return ret:
@ @ -120,9 +120,9 @ @ void free uid(struct user struct *up)
}
-struct user struct * alloc uid(uid t uid)
+struct user_struct * alloc_uid(struct user_namespace *ns, uid_t uid)
- struct list_head *hashent = uidhashentry(uid);
+ struct list_head *hashent = uidhashentry(ns, uid);
 struct user_struct *up;
 spin lock irg(&uidhash lock);
@ @ -211,11 +211,11 @ @ static int __init uid_cache_init(void)
  0, SLAB HWCACHE ALIGN|SLAB PANIC, NULL, NULL);
for(n = 0; n < UIDHASH_SZ; ++n)

    INIT_LIST_HEAD(uidhash_table + n);

+ INIT_LIST_HEAD(init_user_ns.uidhash_table + n);
 /* Insert the root user immediately (init already runs as root) */
 spin lock irg(&uidhash lock);
- uid hash insert(&root user, uidhashentry(0));
+ uid hash insert(&root user, uidhashentry(&init user ns, 0));
 spin unlock irg(&uidhash lock);
 return 0:
Index: 2.6.19-rc6-mm2/include/linux/nsproxy.h
--- 2.6.19-rc6-mm2.orig/include/linux/nsproxy.h
+++ 2.6.19-rc6-mm2/include/linux/nsproxy.h
@ @ -9,6 +9,7 @ @ struct uts namespace;
```

```
struct ipc namespace;
struct pid namespace;
struct net_namespace;
+struct user_namespace;
 * A structure to contain pointers to all per-process
@ @ -31,6 +32,7 @ @ struct nsproxy {
 struct mnt namespace *mnt ns;
 struct pid namespace *pid ns;
 struct net namespace *net ns;
+ struct user namespace *user ns;
};
extern struct nsproxy init_nsproxy;
Index: 2.6.19-rc6-mm2/include/linux/sched.h
--- 2.6.19-rc6-mm2.orig/include/linux/sched.h
+++ 2.6.19-rc6-mm2/include/linux/sched.h
@@ -252,6 +252,7 @@ extern signed long schedule timeout unin
asmlinkage void schedule(void);
struct nsproxy;
+struct user_namespace;
/* Maximum number of active map areas.. This is a random (large) number */
#define DEFAULT_MAX_MAP_COUNT 65536
@ @ -1286,7 +1287,7 @ @ extern struct task struct *find task by
extern void set special pids(pid t session, pid t pgrp);
/* per-UID process charging. */
-extern struct user_struct * alloc_uid(uid_t);
+extern struct user_struct * alloc_uid(struct user_namespace *, uid_t);
static inline struct user_struct *get_uid(struct user_struct *u)
{
 atomic_inc(&u->__count);
Index: 2.6.19-rc6-mm2/include/linux/init_task.h
--- 2.6.19-rc6-mm2.orig/include/linux/init_task.h
+++ 2.6.19-rc6-mm2/include/linux/init_task.h
@@ -9,6 +9,7 @@
#include linux/ipc.h>
#include linux/pid namespace.h>
#include linux/net_namespace.h>
+#include linux/user namespace.h>
#define INIT FDTABLE \
     \
```

```
@ @ -81,6 +82,7 @ @ extern struct nsproxy init_nsproxy;
 .mnt ns = NULL.
 INIT_NET_NS(net_ns)
 INIT_IPC_NS(ipc_ns)
+ .user_ns = &init_user_ns, \
#define INIT_SIGHAND(sighand) {
Index: 2.6.19-rc6-mm2/kernel/sys.c
--- 2.6.19-rc6-mm2.orig/kernel/sys.c
+++ 2.6.19-rc6-mm2/kernel/svs.c
@ @ -33,6 +33,7 @ @
#include linux/compat.h>
#include linux/syscalls.h>
#include linux/kprobes.h>
+#include linux/user namespace.h>
#include <asm/uaccess.h>
#include <asm/io.h>
@ @ -1062,13 +1063,13 @ @ static int set user(uid t new ruid, int
 struct user_struct *new_user;
- new_user = alloc_uid(new_ruid);
+ new_user = alloc_uid(current->nsproxy->user_ns, new_ruid);
 if (!new_user)
 return -EAGAIN;
 if (atomic_read(&new_user->processes) >=
  current->signal->rlim[RLIMIT NPROC].rlim cur &&
- new user!= &root user) {
+ new_user != current->nsproxy->user_ns->root_user) {
 free_uid(new_user);
 return -EAGAIN;
Index: 2.6.19-rc6-mm2/kernel/fork.c
--- 2.6.19-rc6-mm2.orig/kernel/fork.c
+++ 2.6.19-rc6-mm2/kernel/fork.c
@ @ -996,7 +996,7 @ @ static struct task struct *copy process(
 if (atomic_read(&p->user->processes) >=
  p->signal->rlim[RLIMIT_NPROC].rlim_cur) {
 if (!capable(CAP_SYS_ADMIN) && !capable(CAP_SYS_RESOURCE) &&
  p->user!= &root user)
    p->user!= current->nsproxy->user_ns->root_user)
  goto bad_fork_free;
 }
```

```
--- 2.6.19-rc6-mm2.orig/kernel/nsproxy.c
+++ 2.6.19-rc6-mm2/kernel/nsproxy.c
@@ -74,6 +74,8 @@ struct nsproxy *dup_namespaces(struct ns
  get pid ns(ns->pid ns);
 if (ns->net_ns)
  get_net_ns(ns->net_ns);
+ if (ns->user ns)
+ get_user_ns(ns->user_ns);
 return ns:
@ @ -125,10 +127,17 @ @ int copy_namespaces(int flags, struct ta
 if (err)
 goto out_net;
+ err = copy_user_ns(flags, tsk);
+ if (err)
+ goto out_user;
out:
 put_nsproxy(old_ns);
 return err;
+out_user:
+ if (new ns->net ns)
+ put net ns(new ns->net ns);
out net:
if (new ns->pid ns)
 put_pid_ns(new_ns->pid_ns);
@ @ -159,5 +168,7 @ @ void free_nsproxy(struct nsproxy *ns)
 put_pid_ns(ns->pid_ns);
 if (ns->net_ns)
 put_net_ns(ns->net_ns);
+ if (ns->user_ns)
+ put_user_ns(ns->user_ns);
 kfree(ns);
Index: 2.6.19-rc6-mm2/init/Kconfig
_____
--- 2.6.19-rc6-mm2.orig/init/Kconfig
+++ 2.6.19-rc6-mm2/init/Kconfig
@@ -222,6 +222,14 @@ config UTS NS
  vservers, to use uts namespaces to provide different
  uts info for different servers. If unsure, say N.
```

```
+config USER_NS
+ bool "User Namespaces"
+ default n
+ help
+ Support user namespaces. This allows containers, i.e.
+ vservers, to use user namespaces to provide different
+ user info for different servers. If unsure, say N.
config AUDIT
 bool "Auditing support"
 depends on NET
Index: 2.6.19-rc6-mm2/include/linux/user namespace.h
--- /dev/null
+++ 2.6.19-rc6-mm2/include/linux/user_namespace.h
@@ -0.0 +1.49 @@
+#ifndef LINUX USER NAMESPACE H
+#define LINUX USER NAMESPACE H
+#include linux/kref.h>
+#include linux/nsproxy.h>
+#define UIDHASH_BITS (CONFIG_BASE_SMALL ? 3 : 8)
+#define UIDHASH_SZ (1 << UIDHASH_BITS)
+struct user_namespace {
+ struct kref kref;
+ struct list head uidhash table[UIDHASH SZ];
+ struct user struct *root user;
+};
+extern struct user_namespace init_user_ns;
+#ifdef CONFIG_USER_NS
+static inline void get_user_ns(struct user_namespace *ns)
+ kref_get(&ns->kref);
+}
+extern int copy user ns(int flags, struct task struct *tsk);
+extern void free_user_ns(struct kref *kref);
+static inline void put_user_ns(struct user_namespace *ns)
+ kref_put(&ns->kref, free_user_ns);
+}
```

```
+#else
+static inline void get_user_ns(struct user_namespace *ns)
+{
+}
+
+static inline int copy_user_ns(int flags, struct task_struct *tsk)
+{
+ return 0;
+}
+static inline void put user ns(struct user namespace *ns)
+{
+}
+#endif
+#endif /* LINUX USER H */
Index: 2.6.19-rc6-mm2/kernel/user_namespace.c
--- /dev/null
+++ 2.6.19-rc6-mm2/kernel/user_namespace.c
@@ -0,0 +1,44 @@
+/*
+ * This program is free software; you can redistribute it and/or
+ * modify it under the terms of the GNU General Public License as
+ * published by the Free Software Foundation, version 2 of the
+ * License.
+ */
+#include linux/module.h>
+#include linux/version.h>
+#include linux/nsproxy.h>
+#include linux/user_namespace.h>
+struct user_namespace init_user_ns = {
+ .kref = {
+ .refcount = ATOMIC_INIT(2),
+ },
+ .root_user = &root_user,
+};
+EXPORT_SYMBOL_GPL(init_user_ns);
+#ifdef CONFIG_USER_NS
+int copy_user_ns(int flags, struct task_struct *tsk)
+{
+ struct user namespace *old ns = tsk->nsproxy->user ns;
```

```
+ int err = 0;
+ if (!old_ns)
+ return 0;
+ get_user_ns(old_ns);
+ return err;
+}
+void free user ns(struct kref *kref)
+{
+ struct user namespace *ns;
+ ns = container_of(kref, struct user_namespace, kref);
+ kfree(ns);
+}
+#endif /* CONFIG_USER_NS */
Index: 2.6.19-rc6-mm2/kernel/Makefile
--- 2.6.19-rc6-mm2.orig/kernel/Makefile
+++ 2.6.19-rc6-mm2/kernel/Makefile
@@ -4,7 +4,7 @@
        = sched.o fork.o exec_domain.o panic.o printk.o profile.o \
   exit.o itimer.o time.o softirq.o resource.o \
   sysctl.o capability.o ptrace.o timer.o user.o \
    sysctl.o capability.o ptrace.o timer.o user.o user namespace.o \
   signal.o sys.o kmod.o workqueue.o pid.o \
   rcupdate.o extable.o params.o posix-timers.o \
   kthread.o wait.o kfifo.o sys ni.o posix-cpu-timers.o mutex.o \
Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers
```

Subject: [patch -mm 03/17] namespace : export unshare of namespace and fs\_struct

Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:27:55 GMT

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From: Cedric Le Goater <clg@fr.ibm.com>

This is needed by the unshare\_ns syscall.

```
Signed-off-by: Cedric Le Goater <clg@fr.ibm.com>
include/linux/mnt_namespace.h |
kernel/fork.c
                         6 +++---
2 files changed, 7 insertions(+), 3 deletions(-)
Index: 2.6.19-rc6-mm2/include/linux/mnt namespace.h
--- 2.6.19-rc6-mm2.orig/include/linux/mnt namespace.h
+++ 2.6.19-rc6-mm2/include/linux/mnt namespace.h
@ @ -14,6 +14,10 @ @ struct mnt namespace {
int event:
};
+extern int unshare_fs(unsigned long unshare_flags, struct fs_struct **new_fsp);
+extern int unshare_mnt_ns(unsigned long unshare_flags,
    struct mnt namespace **new nsp.
    struct fs_struct *new_fs);
extern int copy_mnt_ns(int, struct task_struct *);
extern void put mnt ns(struct mnt namespace *ns);
extern struct mnt namespace *dup mnt ns(struct task struct *,
Index: 2.6.19-rc6-mm2/kernel/fork.c
--- 2.6.19-rc6-mm2.orig/kernel/fork.c
+++ 2.6.19-rc6-mm2/kernel/fork.c
@@ -1498,7 +1498,7 @@ static int unshare thread(unsigned long
/*
 * Unshare the filesystem structure if it is being shared
 */
-static int unshare_fs(unsigned long unshare_flags, struct fs_struct **new_fsp)
+extern int unshare fs(unsigned long unshare flags, struct fs struct **new fsp)
{
 struct fs_struct *fs = current->fs;
@@ -1515,7 +1515,7 @@ static int unshare_fs(unsigned long unsh
/*
  Unshare the mnt_namespace structure if it is being shared
-static int unshare_mnt_namespace(unsigned long unshare_flags,
+extern int unshare mnt ns(unsigned long unshare flags,
 struct mnt namespace **new nsp, struct fs struct *new fs)
 struct mnt_namespace *ns = current->nsproxy->mnt_ns;
@@ -1634,7 +1634,7 @@ asmlinkage long sys_unshare(unsigned lon
 goto bad unshare out:
 if ((err = unshare_fs(unshare_flags, &new_fs)))
 goto bad unshare cleanup thread;
- if ((err = unshare mnt namespace(unshare flags, &new ns, new fs)))
```

```
+ if ((err = unshare_mnt_ns(unshare_flags, &new_ns, new_fs)))
 goto bad unshare cleanup fs;
 if ((err = unshare_sighand(unshare_flags, &new_sigh)))
 goto bad_unshare_cleanup_ns;
Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers
Subject: [patch -mm 04/17] nsproxy: externalizes exit_task_namespaces
Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:27:56 GMT
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From: Cedric Le Goater <clg@fr.ibm.com>
this is required to remove a header dependency in sched.h which breaks
next patches.
Signed-off-by: Cedric Le Goater <clg@fr.ibm.com>
include/linux/nsproxy.h | 15 ++++-----
kernel/nsproxy.c
                   | 11 +++++++++
2 files changed, 15 insertions(+), 11 deletions(-)
Index: 2.6.19-rc6-mm2/include/linux/nsproxy.h
--- 2.6.19-rc6-mm2.orig/include/linux/nsproxy.h
+++ 2.6.19-rc6-mm2/include/linux/nsproxy.h
@ @ -2,7 +2,8 @ @
#define LINUX NSPROXY H
#include linux/spinlock.h>
-#include linux/sched.h>
+struct task_struct;
struct mnt namespace;
struct uts namespace;
@ @ -48,14 +49,6 @ @ static inline void put_nsproxy(struct ns
 }
}
-static inline void exit_task_namespaces(struct task_struct *p)
-{
```

```
- struct nsproxy *ns = p->nsproxy;
- if (ns) {
task_lock(p);
p->nsproxy = NULL;
task_unlock(p);
put_nsproxy(ns);
- }
-}
+extern void exit_task_namespaces(struct task_struct *p);
#endif
Index: 2.6.19-rc6-mm2/kernel/nsproxy.c
--- 2.6.19-rc6-mm2.orig/kernel/nsproxy.c
+++ 2.6.19-rc6-mm2/kernel/nsproxy.c
@@ -37,6 +37,17 @@ void get_task_namespaces(struct task_str
}
}
+void exit_task_namespaces(struct task_struct *p)
+ struct nsproxy *ns = p->nsproxy;
+ if (ns) {
+ task_lock(p);
+ p->nsproxy = NULL;
+ task_unlock(p);
+ put_nsproxy(ns);
+ }
+}
 * creates a copy of "orig" with refcount 1.
 * This does not grab references to the contained namespaces,
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Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers
```

Subject: [patch -mm 05/17] ipc namespace : externalizes unshare\_ipcs Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:27:57 GMT View Forum Message <> Reply to Message

From: Cedric Le Goater <clg@fr.ibm.com>

This patch externalizes unshare\_ipcs when CONFIG\_IPC\_NS is not

header dependencies on sched.h. Signed-off-by: Cedric Le Goater <clg@fr.ibm.com> include/linux/ipc.h | 2 +kernel/fork.c 1 2 +-2 files changed, 2 insertions(+), 2 deletions(-) Index: 2.6.19-rc6-mm2/include/linux/ipc.h --- 2.6.19-rc6-mm2.orig/include/linux/ipc.h +++ 2.6.19-rc6-mm2/include/linux/ipc.h @ @ -99,7 +99,6 @ @ extern struct ipc\_namespace init\_ipc\_ns; #ifdef CONFIG IPC NS extern void free ipc ns(struct kref \*kref); extern int copy ipcs(unsigned long flags, struct task struct \*tsk); -extern int unshare\_ipcs(unsigned long flags, struct ipc\_namespace \*\*ns); #else static inline int copy ipcs(unsigned long flags, struct task struct \*tsk) @ @ -107.6 +106.7 @ @ static inline int copy ipcs(unsigned lon #endif +extern int unshare ipcs(unsigned long flags, struct ipc namespace \*\*ns); static inline struct ipc namespace \*get ipc ns(struct ipc namespace \*ns) #ifdef CONFIG IPC NS Index: 2.6.19-rc6-mm2/kernel/fork.c \_\_\_\_\_ --- 2.6.19-rc6-mm2.orig/kernel/fork.c +++ 2.6.19-rc6-mm2/kernel/fork.c @@ -1591,7 +1591,7 @@ static int unshare semundo(unsigned long } #ifndef CONFIG IPC NS -static inline int unshare ipcs(unsigned long flags, struct ipc namespace \*\*ns) +extern int unshare ipcs(unsigned long flags, struct ipc namespace \*\*ns) if (flags & CLONE\_NEWIPC) return -EINVAL:

set. This is needed by the unshare ns syscall. Unfortunately, it could

not be handled with a dummy static inline in ipc.h because of

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Subject: [patch -mm 06/17] nsproxy: add extern to nsproxy functions Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:27:58 GMT View Forum Message <> Reply to Message

```
From: Cedric Le Goater <clq@fr.ibm.com>
Signed-off-by: Cedric Le Goater <clg@fr.ibm.com>
include/linux/nsproxy.h | 8 ++++----
1 file changed, 4 insertions(+), 4 deletions(-)
Index: 2.6.19-rc6-mm2/include/linux/nsproxy.h
--- 2.6.19-rc6-mm2.orig/include/linux/nsproxy.h
+++ 2.6.19-rc6-mm2/include/linux/nsproxy.h
@@ -37,10 +37,10 @@ struct nsproxy {
};
extern struct nsproxy init nsproxy;
-struct nsproxy *dup_namespaces(struct nsproxy *orig);
-int copy namespaces(int flags, struct task struct *tsk);
-void get_task_namespaces(struct task_struct *tsk);
-void free_nsproxy(struct nsproxy *ns);
+extern struct nsproxy *dup_namespaces(struct nsproxy *orig);
+extern int copy_namespaces(int flags, struct task_struct *tsk);
+extern void get task namespaces(struct task struct *tsk);
+extern void free_nsproxy(struct nsproxy *ns);
static inline void put_nsproxy(struct nsproxy *ns)
{
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```

Subject: [patch -mm 07/17] nsproxy: make put\_nsproxy an extern Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:27:59 GMT View Forum Message <> Reply to Message

From: Cedric Le Goater <clq@fr.ibm.com>

This routine will require modifications in the next patches and it won't be possible to keep it as an inline anymore.

```
Signed-off-by: Cedric Le Goater <clg@fr.ibm.com>
include/linux/nsproxy.h |
kernel/nsproxy.c
                     9++++++
2 files changed, 9 insertions(+), 9 deletions(-)
Index: 2.6.19-rc6-mm2/include/linux/nsproxy.h
--- 2.6.19-rc6-mm2.orig/include/linux/nsproxy.h
+++ 2.6.19-rc6-mm2/include/linux/nsproxy.h
@ @ -40,14 +40,7 @ @ extern struct nsproxy init_nsproxy;
extern struct nsproxy *dup_namespaces(struct nsproxy *orig);
extern int copy namespaces(int flags, struct task struct *tsk);
extern void get task namespaces(struct task struct *tsk);
-extern void free nsproxy(struct nsproxy *ns);
-static inline void put_nsproxy(struct nsproxy *ns)
- if (atomic_dec_and_test(&ns->count)) {
free_nsproxy(ns);
- }
-}
+extern void put_nsproxy(struct nsproxy *ns);
extern void exit task namespaces(struct task struct *p);
Index: 2.6.19-rc6-mm2/kernel/nsproxy.c
--- 2.6.19-rc6-mm2.orig/kernel/nsproxy.c
+++ 2.6.19-rc6-mm2/kernel/nsproxy.c
@@ -167,7 +167,7 @@ out ns:
 goto out:
-void free_nsproxy(struct nsproxy *ns)
+static void free nsproxy(struct nsproxy *ns)
 if (ns->mnt_ns)
 put_mnt_ns(ns->mnt_ns);
@ @ -183,3 +183,10 @ @ void free_nsproxy(struct nsproxy *ns)
 put_user_ns(ns->user_ns);
 kfree(ns);
```

```
+void put_nsproxy(struct nsproxy *ns)
+{
+ if (atomic_dec_and_test(&ns->count)) {
+ free_nsproxy(ns);
+ }
+}
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Subject: [patch -mm 08/17] nsproxy: add hashtable
Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:28:00 GMT
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From: Cedric Le Goater <clg@fr.ibm.com>
This patch adds a hashtable of nsproxy using the nsproxy as a key.
init nsproxy is hashed at init with key 0. This is considered to be
the 'host' nsproxy.
Signed-off-by: Cedric Le Goater <clg@fr.ibm.com>
include/linux/nsproxy.h | 3 +++
kernel/nsproxy.c
                    2 files changed, 46 insertions(+)
Index: 2.6.19-rc6-mm2/include/linux/nsproxy.h
--- 2.6.19-rc6-mm2.orig/include/linux/nsproxy.h
+++ 2.6.19-rc6-mm2/include/linux/nsproxy.h
@@ -2,6 +2,7 @@
#define _LINUX_NSPROXY_H
#include linux/spinlock.h>
+#include linux/list.h>
struct task struct;
@ @ -34,6 +35,8 @ @ struct nsproxy {
 struct pid namespace *pid ns;
 struct net_namespace *net_ns;
 struct user_namespace *user_ns;
```

+ struct hlist\_node ns\_hash\_node;

```
};
extern struct nsproxy init_nsproxy;
Index: 2.6.19-rc6-mm2/kernel/nsproxy.c
--- 2.6.19-rc6-mm2.orig/kernel/nsproxy.c
+++ 2.6.19-rc6-mm2/kernel/nsproxy.c
@ @ -22,6 +22,15 @ @
#include linux/pid namespace.h>
#include linux/net namespace.h>
+#define NS_HASH_BITS 3 /* this might need some configuration */
+#define NS_HASH_SIZE (1 << NS_HASH_BITS)
+#define NS_HASH_MASK (NS_HASH_SIZE - 1)
+#define ns_hashfn(id) (((id >> NS_HASH_BITS) + id) & NS_HASH_MASK)
+#define ns_hash_head(id) &ns_hash[ns_hashfn(id)]
+static struct hlist_head ns_hash[NS_HASH_SIZE];
+static DEFINE SPINLOCK(ns hash lock);
struct nsproxy init nsproxy = INIT NSPROXY(init nsproxy);
static inline void get_nsproxy(struct nsproxy *ns)
@ @ -190.3 +199.37 @ @ void put nsproxy(struct nsproxy *ns)
 free_nsproxy(ns);
 }
}
+/*
+ * This routine must be called with the ns_hash spin_locked
+static inline struct nsproxy *ns_hash_find(int id)
+{
+ struct hlist_node *elem;
+ struct nsproxy *ns;
+ hlist_for_each_entry(ns, elem, ns_hash_head(id), ns_hash_node) {
+ if (ns->id == id) {
+ get_nsproxy(ns);
+ return ns;
+ }
+ }
+ return NULL;
+}
+static int init nshash init(void)
+{
```

```
+ int i;
+ for (i = 0; i < NS_HASH_SIZE; ++i)
+ INIT_HLIST_HEAD(&ns_hash[i]);
+ spin_lock_irq(&ns_hash_lock);
+ hlist_add_head(&init_nsproxy.ns_hash_node, ns_hash_head(0));
+ spin_unlock_irq(&ns_hash_lock);
+ return 0;
+}
+module_init(nshash_init);
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Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers
Subject: [patch -mm 09/17] nsproxy: add namespace flags
Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:28:01 GMT
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From: Cedric Le Goater <clg@fr.ibm.com>
Signed-off-by: Cedric Le Goater <clg@fr.ibm.com>
include/linux/nsproxy.h | 11 ++++++++
1 file changed, 11 insertions(+)
Index: 2.6.19-rc6-mm2/include/linux/nsproxy.h
--- 2.6.19-rc6-mm2.orig/include/linux/nsproxy.h
+++ 2.6.19-rc6-mm2/include/linux/nsproxy.h
@ @ -14,6 +14,17 @ @ struct net_namespace;
struct user_namespace;
+ * namespaces flags
+ */
+#define NS MNT 0x00000001
+#define NS UTS 0x00000002
+#define NS_IPC 0x00000004
+#define NS_PID 0x00000008
+#define NS NET 0x00000010
+#define NS_USER 0x00000020
```

```
+#define NS_ALL (NS_MNT|NS_UTS|NS_IPC|NS_PID|NS_NET|NS_USER)
+/*
 * A structure to contain pointers to all per-process
 * namespaces - fs (mount), uts, network, sysvipc, etc.
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https://lists.osdl.org/mailman/listinfo/containers
Subject: [patch -mm 10/17] nsproxy: add unshare_ns and bind_ns syscalls
Posted by Cedric Le Goater on Tue. 05 Dec 2006 10:28:02 GMT
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From: Cedric Le Goater <clg@fr.ibm.com>
The following patch defines 2 new syscalls specific to nsproxy and
namespaces:
* unshare ns:
enables a process to unshare one or more namespaces. this
    duplicates the unshare syscall for the moment but we
expect to diverge when the number of namespaces increases
* bind ns:
allows a process to bind
1 - its nsproxy to some identifier
2 - to another nsproxy using an identifier or -pid
Here's a sample user space program to use them.
#include <stdio.h>
#include <stdlib.h>
#include <sched.h>
#include <unistd.h>
#include <string.h>
#include <errno.h>
#include <libgen.h>
#include linux/unistd.h>
#if __i386__
```

```
define NR unshare ns 324
# define __NR_bind_ns 325
#elif __ia64__
# define __NR_unshare_ns 1303
# define __NR_bind_ns 1304
#elif __powerpc__
   define __NR_unshare_ns 303
# define __NR_bind_ns 304
#elif s390x
# define NR_unshare_ns 313
# define NR bind ns 314
#elif x86 64
  define __NR_unshare_ns 284
# define __NR_bind_ns 285
#else
   error "Architecture not supported"
#endif
static inline syscall 1 (int, unshare ns, unsigned long, flags)
static inline syscall2 (int, bind ns, int, id, unsigned long, flags)
#define NS MNT 0x00000001
#define NS UTS 0x00000002
#define NS_IPC 0x00000004
#define NS PID 0x00000008
#define NS NET 0x00000010
#define NS_USER 0x00000020
static const char* procname;
static void usage(const char *name)
printf("usage: %s [-h] [-I id] [-muiUnp] [command [arg ...]]\n", name);
printf("\n");
printf(" -h this message\n");
printf("\n");
printf(" -I <id> bind process to nsproxy <id>\n");
printf(" -m mount namespace\n");
printf(" -u utsname namespace\n");
printf(" -i ipc namespace\n");
printf(" -U user namespace\n");
printf(" -n net namespace\n");
printf(" -p pid namespace\n");
printf("\n");
printf("(C) Copyright IBM Corp. 2006\n");
printf("\n");
exit(1);
}
```

```
int main(int argc, char *argv[])
int c;
unsigned long flags = 0;
int id = -1;
procname = basename(argv[0]);
while ((c = getopt(argc, argv, "+muiUnphl:")) != EOF) {
 switch (c) {
 case 'I': if (optarg)
  id = atoi(optarg); break;
 case 'm': flags |= NS_MNT; break;
 case 'u': flags |= NS_UTS; break;
 case 'i': flags |= NS IPC; break;
 case 'U': flags |= NS_USER; break;
 case 'n': flags |= NS_NET; break;
 case 'p': flags |= NS_PID; break;
 case 'h':
 default:
 usage(procname);
};
argv = &argv[optind];
argc = argc - optind;
if (!strcmp(procname, "unsharens")) {
 if (unshare ns(flags) == -1) {
 perror("unshare_ns");
 return 1;
 }
}
if (bind_ns(id, flags) == -1) {
 perror("bind_ns");
 return 1;
if (argc) {
 execve(argv[0], argv, __environ);
 perror("execve");
 return 1;
return 0;
```

```
}
Signed-off-by: Cedric Le Goater <clg@fr.ibm.com>
arch/i386/kernel/syscall_table.S |
arch/ia64/kernel/entry.S
                               2
arch/s390/kernel/compat_wrapper.S |
arch/s390/kernel/syscalls.S
arch/x86 64/ia32/ia32entry.S
                                  2
include/asm-i386/unistd.h
include/asm-ia64/unistd.h
include/asm-powerpc/systbl.h
include/asm-powerpc/unistd.h
                                   4
include/asm-s390/unistd.h
include/asm-x86 64/unistd.h
include/linux/syscalls.h
                               3
kernel/nsproxy.c
                            kernel/sys ni.c
                            4
14 files changed, 366 insertions(+), 9 deletions(-)
Index: 2.6.19-rc6-mm2/include/linux/syscalls.h
--- 2.6.19-rc6-mm2.orig/include/linux/syscalls.h
+++ 2.6.19-rc6-mm2/include/linux/syscalls.h
@@ -605,6 +605,9 @@ asmlinkage long sys_set_robust_list(stru
     size t len):
asmlinkage long sys getcpu(unsigned user *cpu, unsigned user *node, struct getcpu cache
user *cache);
+asmlinkage long sys unshare ns(unsigned long unshare flags);
+asmlinkage long sys_bind_ns(int id, unsigned long unshare_flags);
int kernel_execve(const char *filename, char *const argv[], char *const envp[]);
asmlinkage long sys_kevent_get_events(int ctl_fd, unsigned int min, unsigned int max,
Index: 2.6.19-rc6-mm2/kernel/nsproxy.c
--- 2.6.19-rc6-mm2.orig/kernel/nsproxy.c
+++ 2.6.19-rc6-mm2/kernel/nsproxy.c
@@ -22,7 +22,11 @@
#include linux/pid_namespace.h>
#include linux/net_namespace.h>
-#define NS_HASH_BITS 3 /* this might need some configuration */
+ * nsproxies are stored in a hash but a rbtree might be more
+ * appropriate.
```

```
+ */
+#define NS HASH BITS 3
#define NS_HASH_SIZE (1 << NS_HASH_BITS)
#define NS_HASH_MASK (NS_HASH_SIZE - 1)
#define ns_hashfn(id) (((id >> NS_HASH_BITS) + id) & NS_HASH_MASK)
@ @ -193,11 +197,26 @ @ static void free_nsproxy(struct nsproxy
 kfree(ns):
}
+/*
   put_nsproxy() is similar to free_uid() in kernel/user.c
+ * the lock can be taken from a tasklet context (task getting freed by
+ * RCU) which requires to be irg safe.
+ */
void put_nsproxy(struct nsproxy *ns)
- if (atomic_dec_and_test(&ns->count)) {
free nsproxy(ns);
- }
+ unsigned long flags;
+ local_irq_save(flags);
+ if (atomic_dec_and_lock(&ns->count, &ns_hash_lock)) {
+ BUG_ON(!ns->id);
+ if (ns->id != -1)
+ hlist_del(&ns->ns_hash_node);
+ spin unlock irgrestore(&ns hash lock, flags);
+ free nsproxy(ns);
+ } else {
+ local irg restore(flags);
+ }
}
@ @ -218,6 +237,304 @ @ static inline struct nsproxy *ns hash fi
 return NULL:
}
+static int bind ns(int id, struct nsproxy *ns)
+{
+ struct nsproxy *prev;
+ int ret = 0:
+
+ if (id < 0)
+ return -EINVAL;
+ spin lock irg(&ns hash lock);
```

```
+ prev = ns_hash_find(id);
+ if (!prev) {
+ ns->id = id;
+ hlist_add_head(&ns->ns_hash_node, ns_hash_head(ns->id));
+ }
+ spin_unlock_irq(&ns_hash_lock);
+ if (prev) {
+ ret = -EBUSY;
+ put_nsproxy(prev);
+ }
+ return ret;
+}
+static int switch_ns(int id, unsigned long flags)
+{
+ int err = 0;
+ struct nsproxy *ns = NULL, *old_ns = NULL, *new_ns = NULL;
+ if (flags & ~NS_ALL)
+ return -EINVAL;
+ /* Let 0 be a default value ? */
+ if (!flags)
+ flags = NS_ALL;
+ if (id < 0) {
+ struct task_struct *p;
+ err = -ESRCH;
+ read_lock(&tasklist_lock);
+ p = find_task_by_pid(-id);
+ if (p) {
+ task_lock(p);
+ get_nsproxy(p->nsproxy);
+ ns = p->nsproxy;
+ task_unlock(p);
+ }
+ read_unlock(&tasklist_lock);
+ } else {
+ err = -ENOENT;
+ spin_lock_irq(&ns_hash_lock);
+ ns = ns_hash_find(id);
+ spin_unlock_irq(&ns_hash_lock);
+ }
+ if (!ns)
+ goto out;
```

```
+ new_ns = ns;
+ /*
+ * clone current nsproxy and populate it with the namespaces
+ * chosen by flags.
+ */
+ if (flags != NS_ALL) {
+ new_ns = dup_namespaces(current->nsproxy);
+ if (!new ns) {
+ err = -ENOMEM;
+ goto out_ns;
+ }
+ if (flags & NS_MNT) {
+ put_mnt_ns(new_ns->mnt_ns);
+ get_mnt_ns(ns->mnt_ns);
+ new_ns->mnt_ns = ns->mnt_ns;
+ }
+ if (flags & NS_UTS) {
+ put uts ns(new ns->uts ns);
+ get_uts_ns(ns->uts_ns);
+ new_ns->uts_ns = ns->uts_ns;
+ }
+ if (flags & NS_IPC) {
+ put_ipc_ns(new_ns->ipc_ns);
+ new_ns->ipc_ns = get_ipc_ns(ns->ipc_ns);
+ }
+ out ns:
+ put_nsproxy(ns);
+ }
+ task_lock(current);
+ if (new_ns) {
+ old_ns = current->nsproxy;
+ current->nsproxy = new_ns;
+ }
+ task_unlock(current);
+ if (old_ns)
+ put_nsproxy(old_ns);
+ err = 0;
+out:
+ return err;
+}
```

```
+
+
+/*
   bind_ns - bind the nsproxy of a task to an id or bind a task to a
         identified nsproxy
   @id: nsproxy identifier if positive or pid if negative
   @flags: identifies the namespaces to bind to
+ * bind ns serves 2 purposes.
+ * The first is to bind the nsproxy of the current task to the
+ * identifier @id. If the identifier is already used, -EBUSY is
+ * returned. If the nsproxy is already bound, -EACCES is returned.
+ * flags is not used in that case.
+ * The second use is to bind the current task to a subset of
+ * namespaces of an identified nsproxy. If positive, @id is considered
+ * being an asproxy identifier previously used to bind the asproxy to
+ * @id. If negative, @id is the pid of a task which is another way to
+ * identify a nsproxy. Switching nsproxy is restricted to tasks within
+ * nsproxy 0, the default nsproxy. If unknown, -ENOENT is returned.
+ * @flags is used to bind the task to the selected namespaces.
   Both uses may return -EINVAL for invalid arguments and -EPERM for
+ * insufficient privileges.
+ * Returns 0 on success.
+asmlinkage long sys_bind_ns(int id, unsigned long flags)
+ struct nsproxy *ns = current->nsproxy;
+ int ret = 0:
+ /*
+ * ns is being changed by switch ns(), protect it
+ get_nsproxy(ns);
+
+ * protect ns->id
+ */
+ spin_lock(&ns->nslock);
+ switch (ns->id) {
+ case -1:
+ * only an unbound nsproxy can be bound to an id.
  */
```

```
+ ret = bind_ns(id, ns);
+ break:
+ case 0:
+ if (!capable(CAP_SYS_ADMIN)) {
+ ret = -EPERM;
  goto unlock;
+ }
+
+ /*
  * only nsproxy 0 can switch nsproxy. if target id is
 * 0, this is a nop.
+ */
+ if (id)
+ ret = switch_ns(id, flags);
+ break;
+ default:
+ /*
 * current nsproxy is already bound. forbid any
+ * switch.
+ */
+ ret = -EACCES;
+ }
+unlock:
+ spin_unlock(&ns->nslock);
+ put_nsproxy(ns);
+ return ret;
+}
+
+/*
+ * sys_unshare_ns - unshare one or more of namespaces which were
             originally shared using clone.
+ * @unshare_ns_flags: identifies the namespaces to unshare
+ * this is for the moment a pale copy of unshare but we expect to
+ * diverge when the number of namespaces increases, the number of
+ * available clone flags is also very low. This is one way to get
+ * some air.
+ *
+asmlinkage long sys_unshare_ns(unsigned long unshare_ns_flags)
+{
+ int err = 0;
+ struct nsproxy *new_nsproxy = NULL, *old_nsproxy = NULL;
+ struct fs struct *fs, *new fs = NULL;
+ struct mnt namespace *mnt, *new mnt = NULL;
```

```
+ struct uts_namespace *uts, *new_uts = NULL;
+ struct ipc_namespace *ipc, *new_ipc = NULL;
+ unsigned long unshare_flags = 0;
+
+ /* Return -EINVAL for all unsupported flags */
+ err = -EINVAL;
+ if (unshare_ns_flags & ~NS_ALL)
 goto bad_unshare_ns_out;
+ /*
  * compatibility with unshare()/clone() : convert ns flags to
  * clone flags
 */
+ if (unshare_ns_flags & NS_MNT)
+ unshare_flags |= CLONE_NEWNS|CLONE_FS;
+ if (unshare_ns_flags & NS_UTS)
+ unshare flags |= CLONE NEWUTS;
+ if (unshare_ns_flags & NS_IPC)
+ unshare_flags |= CLONE_NEWIPC;
+ if ((err = unshare_fs(unshare_flags, &new_fs)))
+ goto bad unshare ns out;
+ if ((err = unshare_mnt_ns(unshare_flags, &new_mnt, new_fs)))
+ goto bad_unshare_ns_cleanup_fs;
+ if ((err = unshare_utsname(unshare_flags, &new_uts)))
+ goto bad_unshare_ns_cleanup_mnt;
+ if ((err = unshare_ipcs(unshare_flags, &new_ipc)))
+ goto bad_unshare_ns_cleanup_uts;
+ if (new_mnt || new_uts || new_ipc) {
+ old nsproxy = current->nsproxy;
+ new_nsproxy = dup_namespaces(old_nsproxy);
+ if (!new_nsproxy) {
  err = -ENOMEM;
   goto bad_unshare_ns_cleanup_ipc;
+
  }
+ }
 if (new_fs || new_mnt || new_uts || new_ipc) {
 task lock(current);
+
  if (new_nsproxy) {
   current->nsproxy = new_nsproxy;
   new_nsproxy = old_nsproxy;
+
  if (new fs) {
```

```
+ fs = current->fs;
   current->fs = new fs;
+ new_fs = fs;
+ }
+
+ if (new_mnt) {
   mnt = current->nsproxy->mnt_ns;
   current->nsproxy->mnt_ns = new_mnt;
+
   new_mnt = mnt;
  }
+
+ if (new_uts) {
   uts = current->nsproxy->uts_ns;
   current->nsproxy->uts_ns = new_uts;
   new_uts = uts;
+
+
  }
+ if (new_ipc) {
  ipc = current->nsproxy->ipc_ns;
   current->nsproxy->ipc_ns = new_ipc;
   new_ipc = ipc;
+
+
  task_unlock(current);
+ }
+ if (new_nsproxy)
+ put_nsproxy(new_nsproxy);
+bad_unshare_ns_cleanup_ipc:
+ if (new_ipc)
+ put_ipc_ns(new_ipc);
+bad_unshare_ns_cleanup_uts:
+ if (new_uts)
+ put_uts_ns(new_uts);
+bad_unshare_ns_cleanup_mnt:
+ if (new_mnt)
+ put_mnt_ns(new_mnt);
+
+bad_unshare_ns_cleanup_fs:
+ if (new_fs)
+ put_fs_struct(new_fs);
+bad_unshare_ns_out:
+ return err;
+}
```

```
static int init nshash init(void)
int i;
Index: 2.6.19-rc6-mm2/kernel/sys_ni.c
_____
--- 2.6.19-rc6-mm2.orig/kernel/sys ni.c
+++ 2.6.19-rc6-mm2/kernel/sys_ni.c
@ @ -146,3 +146,7 @ @ cond_syscall(compat_sys_migrate_pages);
cond_syscall(sys_bdflush);
cond_syscall(sys_ioprio_set);
cond syscall(sys ioprio get);
+/* user resources syscalls */
+cond_syscall(sys_unshare_ns);
+cond_syscall(sys_bind_ns);
Index: 2.6.19-rc6-mm2/arch/i386/kernel/syscall_table.S
--- 2.6.19-rc6-mm2.orig/arch/i386/kernel/syscall table.S
+++ 2.6.19-rc6-mm2/arch/i386/kernel/syscall_table.S
@ @ -323,3 +323,5 @ @ ENTRY(sys call table)
 .long sys kevent ctl
 .long sys_kevent_wait
 .long sys_kevent_ring_init
+ .long sys_unshare_ns
+ .long sys_bind_ns /* 325 */
Index: 2.6.19-rc6-mm2/arch/ia64/kernel/entry.S
--- 2.6.19-rc6-mm2.orig/arch/ia64/kernel/entry.S
+++ 2.6.19-rc6-mm2/arch/ia64/kernel/entry.S
@@ -1610,5 +1610,7 @@ sys call table:
 data8 sys_sync_file_range // 1300
 data8 sys_tee
 data8 sys_vmsplice
+ data8 sys_unshare_ns
+ data8 sys_bind_ns
 .org sys_call_table + 8*NR_syscalls // guard against failures to increase NR_syscalls
Index: 2.6.19-rc6-mm2/arch/s390/kernel/compat wrapper.S
--- 2.6.19-rc6-mm2.orig/arch/s390/kernel/compat wrapper.S
+++ 2.6.19-rc6-mm2/arch/s390/kernel/compat_wrapper.S
@ @ -1665,3 +1665,14 @ @ sys_getcpu_wrapper:
 llgtr %r3,%r3 # unsigned *
 llgtr %r4,%r4 # struct getcpu_cache *
 jg sys_getcpu
+ .globl sys unshare ns wrapper
```

```
+sys unshare ns wrapper:
+ llgfr %r2,%r2
                        # unsigned long
+ jg
      sys_unshare_ns
+ .qlobl sys_bind_ns_wrapper
+sys_bind_ns_wrapper:
+ lgfr %r2,%r2
                        # int
+ llgfr %r3,%r3
                       # unsigned long
      sys bind ns
+ iq
Index: 2.6.19-rc6-mm2/arch/s390/kernel/syscalls.S
--- 2.6.19-rc6-mm2.orig/arch/s390/kernel/syscalls.S
+++ 2.6.19-rc6-mm2/arch/s390/kernel/syscalls.S
@@ -321,3 +321,5 @@ SYSCALL(sys_vmsplice,sys_vmsplice,compat
                /* 310 sys_move_pages */
NI SYSCALL
SYSCALL(sys_getcpu,sys_getcpu_wrapper)
SYSCALL(sys epoll pwait,sys epoll pwait,sys ni syscall)
+SYSCALL(sys unshare ns,sys unshare ns,sys unshare ns wrapper)
+SYSCALL(sys bind ns,sys bind ns,sys bind ns wrapper)
Index: 2.6.19-rc6-mm2/arch/x86 64/ia32/ia32entry.S
--- 2.6.19-rc6-mm2.orig/arch/x86 64/ia32/ia32entry.S
+++ 2.6.19-rc6-mm2/arch/x86_64/ia32/ia32entry.S
@ @ -722,4 +722,6 @ @ ia32_sys_call_table:
 .quad sys_kevent_ctl /* 320 */
 .quad sys kevent wait
 .quad sys_kevent_ring_init
+ .quad sys unshare ns
+ .quad sys bind ns
ia32 syscall end:
Index: 2.6.19-rc6-mm2/include/asm-i386/unistd.h
--- 2.6.19-rc6-mm2.orig/include/asm-i386/unistd.h
+++ 2.6.19-rc6-mm2/include/asm-i386/unistd.h
@@ -329,10 +329,12 @@
#define NR kevent ctl 321
#define __NR_kevent_wait 322
#define NR kevent ring init 323
+#define NR unshare ns 324
+#define NR bind ns 325
#ifdef KERNEL
-#define NR_syscalls 324
+#define NR syscalls 326
#define ARCH WANT IPC PARSE VERSION
#define ARCH WANT OLD READDIR
```

```
_____
--- 2.6.19-rc6-mm2.orig/include/asm-ia64/unistd.h
+++ 2.6.19-rc6-mm2/include/asm-ia64/unistd.h
@@ -291,11 +291,13 @@
#define __NR_sync_file_range 1300
#define NR tee 1301
#define __NR_vmsplice 1302
+#define NR unshare ns 1303
+#define NR bind ns 1304
#ifdef KERNEL
-#define NR_syscalls 279 /* length of syscall table */
+#define NR_syscalls 281 /* length of syscall table */
#define ARCH WANT SYS RT SIGACTION
Index: 2.6.19-rc6-mm2/include/asm-powerpc/systbl.h
--- 2.6.19-rc6-mm2.orig/include/asm-powerpc/systbl.h
+++ 2.6.19-rc6-mm2/include/asm-powerpc/systbl.h
@ @ -305,3 +305,5 @ @ SYSCALL_SPU(faccessat)
COMPAT_SYS_SPU(get_robust_list)
COMPAT SYS SPU(set robust list)
COMPAT_SYS(move_pages)
+SYSCALL SPU(unshare ns)
+SYSCALL SPU(bind ns)
Index: 2.6.19-rc6-mm2/include/asm-powerpc/unistd.h
--- 2.6.19-rc6-mm2.orig/include/asm-powerpc/unistd.h
+++ 2.6.19-rc6-mm2/include/asm-powerpc/unistd.h
@ @ -324,10 +324,12 @ @
#define __NR_get_robust_list 299
#define NR set robust list 300
#define __NR_move_pages 301
+#define NR unshare ns 302
+#define NR bind ns 303
#ifdef KERNEL
-#define __NR_syscalls 302
+#define __NR_syscalls 304
#define __NR__exit __NR_exit
#define NR syscalls NR syscalls
Index: 2.6.19-rc6-mm2/include/asm-s390/unistd.h
```

```
--- 2.6.19-rc6-mm2.orig/include/asm-s390/unistd.h
+++ 2.6.19-rc6-mm2/include/asm-s390/unistd.h
@@ -250.8 +250.10 @@
/* Number 310 is reserved for new sys move pages */
#define __NR_getcpu 311
#define NR epoll pwait 312
+#define __NR_unshare ns 313
+#define NR bind ns 314
-#define NR_syscalls 313
+#define NR syscalls 315
 * There are some system calls that are not present on 64 bit, some
Index: 2.6.19-rc6-mm2/include/asm-x86 64/unistd.h
--- 2.6.19-rc6-mm2.orig/include/asm-x86 64/unistd.h
+++ 2.6.19-rc6-mm2/include/asm-x86 64/unistd.h
@@ -627,8 +627,12 @@ SYSCALL( NR kevent ctl, sys kevent ct
  SYSCALL( NR kevent wait, sys kevent wait)
#define NR kevent ring init 283
  SYSCALL(__NR_kevent_ring_init, sys_kevent_ring_init)
+#define __NR_unshare_ns 284
+__SYSCALL(__NR_unshare_ns, sys_unshare_ns)
+#define NR bind ns 285
+__SYSCALL(__NR_bind_ns, sys_bind_ns)
-#define NR syscall max NR kevent ring init
+#define __NR_syscall_max __NR_bind_ns
#ifndef NO STUBS
#define __ARCH_WANT_OLD_READDIR
Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers
```

Subject: [patch -mm 11/17] user namespace: add user\_namespace ptr to vfsmount Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:28:03 GMT View Forum Message <> Reply to Message

From: Serge E. Hallyn <serue@us.ibm.com>

Add user\_namespace ptr to vfsmount, and define a helper to compare it

```
to the task's user ns.
Signed-off-by: Serge E. Hallyn <serue@us.ibm.com>
fs/namespace.c
                      4 ++++
include/linux/mount.h | 2 ++
include/linux/sched.h | 12 +++++++++
3 files changed, 18 insertions(+)
Index: 2.6.19-rc6-mm2/fs/namespace.c
--- 2.6.19-rc6-mm2.orig/fs/namespace.c
+++ 2.6.19-rc6-mm2/fs/namespace.c
@@ -25,6 +25,7 @@
#include linux/security.h>
#include linux/mount.h>
#include linux/ramfs.h>
+#include linux/user namespace.h>
#include <asm/uaccess.h>
#include <asm/unistd.h>
#include "pnode.h"
@@ -56,6 +57,8 @@ struct vfsmount *alloc vfsmnt(const char
 struct vfsmount *mnt = kmem_cache_alloc(mnt_cache, GFP_KERNEL);
 if (mnt) {
 memset(mnt, 0, sizeof(struct vfsmount));
+ mnt->mnt_user_ns = current->nsproxy->user_ns;
+ get_user_ns(mnt->mnt_user_ns);
 atomic set(&mnt->mnt count, 1);
 INIT LIST HEAD(&mnt->mnt hash);
 INIT_LIST_HEAD(&mnt->mnt_child);
@ @ -88.6 +91.7 @ @ EXPORT SYMBOL(simple set mnt):
void free_vfsmnt(struct vfsmount *mnt)
+ put_user_ns(mnt->mnt_user_ns);
 kfree(mnt->mnt devname);
 kmem_cache_free(mnt_cache, mnt);
}
Index: 2.6.19-rc6-mm2/include/linux/mount.h
--- 2.6.19-rc6-mm2.orig/include/linux/mount.h
+++ 2.6.19-rc6-mm2/include/linux/mount.h
@ @ -21,6 +21,7 @ @ struct super_block;
struct vfsmount;
struct dentry;
struct mnt_namespace;
+struct user_namespace;
```

```
#define MNT NOSUID 0x01
#define MNT NODEV 0x02
@ @ -53,6 +54,7 @ @ struct vfsmount {
 struct list_head mnt_slave; /* slave list entry */
 struct vfsmount *mnt_master; /* slave is on master->mnt_slave_list */
 struct mnt_namespace *mnt_ns; /* containing namespace */
+ struct user namespace *mnt user ns: /* namespace for uid interpretation */
int mnt_pinned;
};
Index: 2.6.19-rc6-mm2/include/linux/sched.h
--- 2.6.19-rc6-mm2.orig/include/linux/sched.h
+++ 2.6.19-rc6-mm2/include/linux/sched.h
@@ -83,6 +83,8 @@ struct sched_param {
#include linux/timer.h>
#include linux/hrtimer.h>
#include linux/task io accounting.h>
+#include linux/nsproxy.h>
+#include linux/mount.h>
#include <asm/processor.h>
@ @ -1589,6 +1591,16 @ @ extern int cond_resched(void);
extern int cond_resched_lock(spinlock_t * lock);
extern int cond_resched_softirg(void);
+static inline int task_mnt_same_uid(struct task_struct *tsk,
    struct vfsmount *mnt)
+{
+ if (tsk->nsproxy == init_task.nsproxy)
+ return 1:
+ if (mnt->mnt_user_ns == tsk->nsproxy->user_ns)
+ return 1;
+ return 0;
+}
 * Does a critical section need to be broken due to another
 * task waiting?:
Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers
```

Page 41 of 124 ---- Generated from OpenVZ Forum

Subject: [patch -mm 12/17] user namespace: hook permission Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:28:04 GMT

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From: Serge Hallyn <serue@us.ibm.com> Hook permission to check vfsmnt->user ns against current. Signed-off-by: Serge E. Hallyn <serue@us.ibm.com> fs/namei.c | 4 ++++ 1 file changed, 4 insertions(+) Index: 2.6.19-rc6-mm2/fs/namei.c --- 2.6.19-rc6-mm2.orig/fs/namei.c +++ 2.6.19-rc6-mm2/fs/namei.c @ @ -246,6 +246,8 @ @ int permission(struct inode \*inode, int return -EACCES: } + if (nd && !task\_mnt\_same\_uid(current, nd->mnt)) + return -EACCES; \* MAY EXEC on regular files requires special handling: We override @@ -433,6 +435,8 @@ static int exec\_permission\_lite(struct i umode t mode = inode->i mode; + if (!task mnt same uid(current, nd->mnt)) + return -EACCES; if (inode->i op && inode->i op->permission) return -EAGAIN; Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: [patch -mm 13/17] user namespace: implement shared mounts Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:28:05 GMT View Forum Message <> Reply to Message

From: Serge E. Hallyn <serue@us.ibm.com>

Implement shared and private userns mounts. Private userns mounts are ones which no process in another user namespace can use.

Here is a sample mount program which only does a bind mount of arg1 onto arg2, but making the destination a private userns mount.

```
int main(int argc, char *argv[])
int type;
if(argc != 3) {
fprintf(stderr, "usage: %s src dest", argv[0]);
 return 1;
}
fprintf(stdout, "%s %s %s\n", argv[0], argv[1], argv[2]);
type = MS_PRIV_USERNS | MS_BIND;
setfsuid(getuid());
if(mount(argv[1], argv[2], "none", type, "") == -1) {
 perror("mount");
 return 1;
return 0;
Signed-off-by: Serge E. Hallyn <serue@us.ibm.com>
fs/namespace.c
                    | 16 +++++++++
fs/pnode.h
                     1+
include/linux/fs.h | 1+
include/linux/mount.h | 1 +
include/linux/sched.h | 2 ++
5 files changed, 17 insertions(+), 4 deletions(-)
Index: 2.6.19-rc6-mm2/fs/namespace.c
--- 2.6.19-rc6-mm2.orig/fs/namespace.c
+++ 2.6.19-rc6-mm2/fs/namespace.c
@@ -259,6 +259,8 @@ static struct vfsmount *clone mnt(struct
 }
 if (flag & CL_MAKE_SHARED)
  set_mnt_shared(mnt);
+ if (flag & CL_PRIV_USERNS)
+ mnt->mnt_flags |= MNT_PRIV_USERNS;
 /* stick the duplicate mount on the same expiry list
```

```
* as the original if that was on one */
@ @ -370,6 +372,7 @ @ static int show vfsmnt(struct seg file *
 { MNT_NOSUID, ",nosuid" },
 { MNT_NODEV, ",nodev" },
 { MNT_NOEXEC, ",noexec" },
+ { MNT_PRIV_USERNS, ",privuserns" },
 { MNT NOATIME, ",noatime" },
 { MNT_NODIRATIME, ",nodiratime" },
 { 0, NULL }
@ @ -901,11 +904,14 @ @ static int do change type(struct nameida
 * do loopback mount.
 */
-static int do_loopback(struct nameidata *nd, char *old_name, int recurse)
+static int do_loopback(struct nameidata *nd, char *old_name, int recurse,
     int uidns_share)
 struct nameidata old nd;
 struct vfsmount *mnt = NULL:
 int err = mount is safe(nd);
+ int flag = (uidns_share ? CL_PRIV_USERNS : 0);
 if (err)
 return err:
 if (!old_name || !*old_name)
@@ -924,9 +930,9 @@ static int do_loopback(struct nameidata
 err = -ENOMEM;
 if (recurse)
- mnt = copy_tree(old_nd.mnt, old_nd.dentry, 0);
+ mnt = copy tree(old nd.mnt, old nd.dentry, flag);
 else
mnt = clone_mnt(old_nd.mnt, old_nd.dentry, 0);
+ mnt = clone_mnt(old_nd.mnt, old_nd.dentry, flag);
 if (!mnt)
 goto out;
@ @ -1409,6 +1415,8 @ @ long do mount(char *dev name, char *dir
 mnt flags |= MNT NOATIME;
 if (flags & MS_NODIRATIME)
 mnt_flags |= MNT_NODIRATIME;
+ if (flags & MS_PRIV_USERNS)
+ mnt_flags |= MNT_PRIV_USERNS;
 flags &= ~(MS_NOSUID | MS_NOEXEC | MS_NODEV | MS_ACTIVE |
   MS_NOATIME | MS_NODIRATIME);
@ @ -1426,7 +1434,7 @ @ long do mount(char *dev name, char *dir
 retval = do remount(&nd, flags & ~MS REMOUNT, mnt flags,
```

```
data page);
 else if (flags & MS BIND)
retval = do_loopback(&nd, dev_name, flags & MS_REC);
+ retval = do loopback(&nd, dev name, flags & MS REC, flags & MS PRIV USERNS);
else if (flags & (MS_SHARED | MS_PRIVATE | MS_SLAVE | MS_UNBINDABLE))
 retval = do_change_type(&nd, flags);
 else if (flags & MS MOVE)
Index: 2.6.19-rc6-mm2/fs/pnode.h
______
--- 2.6.19-rc6-mm2.orig/fs/pnode.h
+++ 2.6.19-rc6-mm2/fs/pnode.h
@@ -22.6 +22.7 @@
#define CL_COPY_ALL 0x04
#define CL MAKE SHARED 0x08
#define CL PROPAGATION 0x10
+#define CL_PRIV_USERNS 0x20
static inline void set_mnt_shared(struct vfsmount *mnt)
{
Index: 2.6.19-rc6-mm2/include/linux/fs.h
--- 2.6.19-rc6-mm2.orig/include/linux/fs.h
+++ 2.6.19-rc6-mm2/include/linux/fs.h
@ @ -120,6 +120,7 @ @ extern int dir notify enable:
#define MS_PRIVATE (1<<18) /* change to private */
#define MS SLAVE (1<<19) /* change to slave */
#define MS_SHARED (1<<20) /* change to shared */
+#define MS_PRIV_USERNS (1<<21) /* compare user namespaces for permission */
#define MS ACTIVE (1<<30)
#define MS_NOUSER (1<<31)
Index: 2.6.19-rc6-mm2/include/linux/mount.h
--- 2.6.19-rc6-mm2.orig/include/linux/mount.h
+++ 2.6.19-rc6-mm2/include/linux/mount.h
@ @ -34.6 +34.7 @ @ struct user namespace:
#define MNT_SHARED 0x1000 /* if the vfsmount is a shared mount */
#define MNT_UNBINDABLE 0x2000 /* if the vfsmount is a unbindable mount */
#define MNT PNODE MASK 0x3000 /* propogation flag mask */
+#define MNT_PRIV_USERNS 0x4000 /* compare user namespaces for permission */
struct vfsmount {
struct list head mnt hash;
Index: 2.6.19-rc6-mm2/include/linux/sched.h
______
--- 2.6.19-rc6-mm2.orig/include/linux/sched.h
+++ 2.6.19-rc6-mm2/include/linux/sched.h
@ @ -1596,6 +1596,8 @ @ static inline int task mnt same uid(stru
```

```
if (tsk->nsproxy == init_task.nsproxy)
 return 1;
+ if (!(mnt->mnt_flags & MNT_PRIV_USERNS))
+ return 1;
 if (mnt->mnt_user_ns == tsk->nsproxy->user_ns)
 return 1;
 return 0;
Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers
Subject: [patch -mm 14/17] user namespace: maintain user ns for priv_userns
mounts to vfsmount
Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:28:06 GMT
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From: Serge E. Hallyn <serue@us.ibm.com>
If a vfsmount is marked priv_userns, then a bind mount should
maintain the old->user ns. Not doing so presents a simple way
around the priv_userns security mechanism:)
Signed-off-by: Serge E. Hallyn <serue@us.ibm.com>
fs/namespace.c |
1 file changed, 2 insertions(+)
Index: 2.6.19-rc6-mm2/fs/namespace.c
--- 2.6.19-rc6-mm2.orig/fs/namespace.c
+++ 2.6.19-rc6-mm2/fs/namespace.c
@@ -240,6 +240,8 @@ static struct vfsmount *clone mnt(struct
 if (mnt) {
 mnt->mnt_flags = old->mnt_flags;
+ if (mnt->mnt flags & MNT PRIV USERNS)
+ mnt->mnt user ns = old->mnt user ns;
 atomic_inc(&sb->s_active);
 mnt->mnt sb = sb;
 mnt->mnt_root = dget(root);
```

```
Subject: [patch -mm 15/17] pid namespace: add unshare Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:28:07 GMT View Forum Message <> Reply to Message
```

```
From: Cedric Le Goater <clg@fr.ibm.com>
Signed-off-by: Cedric Le Goater <clq@fr.ibm.com>
include/linux/pid_namespace.h |
kernel/nsproxy.c
                      25 +++++++++++++++++++++
kernel/pid.c
                    3 files changed, 68 insertions(+), 3 deletions(-)
Index: 2.6.19-rc6-mm2/include/linux/pid_namespace.h
_____
--- 2.6.19-rc6-mm2.orig/include/linux/pid_namespace.h
+++ 2.6.19-rc6-mm2/include/linux/pid namespace.h
@@ -29,6 +29,8 @@ static inline void get_pid_ns(struct pid
kref_get(&ns->kref);
}
+extern int unshare_pid_ns(unsigned long unshare_ns_flags,
   struct pid_namespace **new_pid);
extern int copy_pid_ns(int flags, struct task_struct *tsk);
extern void free pid ns(struct kref *kref);
Index: 2.6.19-rc6-mm2/kernel/pid.c
--- 2.6.19-rc6-mm2.orig/kernel/pid.c
+++ 2.6.19-rc6-mm2/kernel/pid.c
@ @ -361,6 +361,50 @ @ struct pid *find_ge_pid(int nr)
EXPORT_SYMBOL_GPL(find_get_pid);
+static struct pid_namespace *clone_pid_ns(struct pid_namespace *old_ns)
+ struct pid namespace *ns;
+ int i;
+ ns = kmalloc(sizeof(struct pid_namespace), GFP_KERNEL);
+ if (!ns)
+ return ns;
```

```
+ kref_init(&ns->kref);
+ atomic_set(&ns->pidmap[0].nr_free, BITS_PER_PAGE - 1);
+ ns->pidmap[0].page = kzalloc(PAGE_SIZE, GFP_KERNEL);
+ if (!ns->pidmap[0].page) {
+ kfree(ns);
+ return NULL;
+ }
+ set bit(0, ns->pidmap[0].page);
+ for (i = 1; i < PIDMAP_ENTRIES; i++) {
+ atomic_set(&ns->pidmap[i].nr_free, BITS_PER_PAGE);
+ ns->pidmap[i].page = NULL;
+ }
+ ns->last_pid = 0;
+ ns->child reaper = current;
+ return ns;
+}
+int unshare_pid_ns(unsigned long unshare_ns_flags,
    struct pid namespace **new pid)
+{
+ if (unshare_ns_flags & NS_PID) {
+ if (!capable(CAP_SYS_ADMIN))
 return -EPERM;
+ *new pid = clone pid ns(current->nsproxy->pid ns);
+ if (!*new pid)
+ return -ENOMEM;
+ }
+ return 0;
+}
int copy_pid_ns(int flags, struct task_struct *tsk)
 struct pid namespace *old ns = tsk->nsproxy->pid ns;
Index: 2.6.19-rc6-mm2/kernel/nsproxy.c
______
--- 2.6.19-rc6-mm2.orig/kernel/nsproxy.c
+++ 2.6.19-rc6-mm2/kernel/nsproxy.c
@ @ -324,6 +324,11 @ @ static int switch_ns(int id, unsigned lo
  put_ipc_ns(new_ns->ipc_ns);
  new_ns->ipc_ns = get_ipc_ns(ns->ipc_ns);
+ if (flags & NS PID) {
+ get pid ns(ns->pid ns);
```

```
+ put pid ns(new ns->pid ns);
+ new ns->pid ns = ns->pid ns;
+ }
 out_ns:
 put_nsproxy(ns);
@@ -440.6 +445.7 @@ asmlinkage long sys_unshare_ns(unsigned
 struct mnt_namespace *mnt, *new_mnt = NULL;
 struct uts namespace *uts, *new uts = NULL;
 struct ipc namespace *ipc, *new ipc = NULL;
+ struct pid_namespace *pid, *new_pid = NULL;
 unsigned long unshare flags = 0;
 /* Return -EINVAL for all unsupported flags */
@ @ -467,16 +473,19 @ @ asmlinkage long sys_unshare_ns(unsigned
 if ((err = unshare_ipcs(unshare_flags, &new_ipc)))
  goto bad unshare ns cleanup uts:
- if (new mnt || new uts || new ipc) {
+ if ((err = unshare_pid_ns(unshare_ns_flags, &new_pid)))
+ goto bad unshare ns cleanup ipc;
+ if (new_mnt || new_uts || new_ipc || new_pid) {
  old_nsproxy = current->nsproxy;
  new_nsproxy = dup_namespaces(old_nsproxy);
  if (!new nsproxy) {
  err = -ENOMEM;
  goto bad_unshare_ns_cleanup_ipc;
 goto bad unshare ns cleanup pid;
  }
 }
- if (new_fs || new_mnt || new_uts || new_ipc) {
+ if (new_fs || new_mnt || new_uts || new_ipc || new_pid) {
  task_lock(current);
@@ -509,12 +518,22 @@ asmlinkage long sys unshare ns(unsigned
  new ipc = ipc;
  }
+ if (new pid) {
+ pid = current->nsproxy->pid_ns;
+ current->nsproxy->pid_ns = new_pid;
+ new pid = pid;
+ }
  task unlock(current);
```

```
}
 if (new_nsproxy)
  put_nsproxy(new_nsproxy);
+bad_unshare_ns_cleanup_pid:
+ if (new_pid)
+ put_pid_ns(new_pid);
bad unshare ns cleanup ipc:
 if (new_ipc)
  put_ipc_ns(new_ipc);
Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers
Subject: [patch -mm 16/17] net namespace: add unshare
Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:28:08 GMT
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From: Cedric Le Goater <clq@fr.ibm.com>
Signed-off-by: Cedric Le Goater <clg@fr.ibm.com>
include/linux/net_namespace.h | 13 ++++++++++
kernel/nsproxy.c
                      3 files changed, 70 insertions(+), 3 deletions(-)
Index: 2.6.19-rc6-mm2/include/linux/net namespace.h
--- 2.6.19-rc6-mm2.orig/include/linux/net namespace.h
+++ 2.6.19-rc6-mm2/include/linux/net_namespace.h
@@ -3,6 +3,7 @@
#include linux/kref.h>
#include linux/nsproxv.h>
+#include linux/errno.h>
struct net_namespace {
 struct kref kref;
@ @ -19,6 +20,9 @ @ static inline void get_net_ns(struct net
 kref_get(&ns->kref);
```

```
+extern int unshare net ns(unsigned long unshare flags,
    struct net_namespace **new_net);
extern int copy_net_ns(int flags, struct task_struct *tsk);
extern void free net ns(struct kref *kref);
@ @ -36,6 +40,15 @ @ static inline void get_net_ns(struct net
}
+static inline int unshare net ns(unsigned long unshare flags,
    struct net_namespace **new_net)
+{
+ if (unshare_flags & NS_NET)
+ return -EINVAL;
+ return 0;
+}
static inline int copy_net_ns(int flags, struct task_struct *tsk)
return 0;
Index: 2.6.19-rc6-mm2/kernel/nsproxy.c
_____
--- 2.6.19-rc6-mm2.orig/kernel/nsproxy.c
+++ 2.6.19-rc6-mm2/kernel/nsproxy.c
@ @ -329,6 +329,12 @ @ static int switch ns(int id, unsigned lo
  put pid ns(new ns->pid ns);
  new_ns->pid_ns = ns->pid_ns;
+ if (flags & NS NET) {
+ get_net_ns(ns->net_ns);
+ put_net_ns(new_ns->net_ns);
+ new_ns->net_ns = ns->net_ns;
+ }
 out ns:
 put_nsproxy(ns);
@@ -446,6 +452,7 @@ asmlinkage long sys unshare ns(unsigned
 struct uts_namespace *uts, *new_uts = NULL;
 struct ipc_namespace *ipc, *new_ipc = NULL;
 struct pid_namespace *pid, *new_pid = NULL;
+ struct net_namespace *net, *new_net = NULL;
 unsigned long unshare_flags = 0;
 /* Return -EINVAL for all unsupported flags */
```

```
@ @ -475,17 +482,19 @ @ asmlinkage long sys_unshare_ns(unsigned
 if ((err = unshare_pid_ns(unshare_ns_flags, &new_pid)))
 goto bad_unshare_ns_cleanup_ipc;
+ if ((err = unshare_net_ns(unshare_ns_flags, &new_net)))
+ goto bad_unshare_ns_cleanup_pid;
- if (new_mnt || new_uts || new_ipc || new_pid) {
+ if (new mnt || new uts || new ipc || new pid || new net) {
  old nsproxy = current->nsproxy;
  new_nsproxy = dup_namespaces(old_nsproxy);
  if (!new nsproxy) {
  err = -ENOMEM;
  goto bad_unshare_ns_cleanup_pid;
 goto bad_unshare_ns_cleanup_net;
  }
 }
- if (new_fs || new_mnt || new_uts || new_ipc || new_pid) {
+ if (new_fs || new_mnt || new_uts || new_ipc || new_pid || new_net) {
  task lock(current);
@ @ -524,12 +533,22 @ @ asmlinkage long sys_unshare_ns(unsigned
  new_pid = pid;
 }
+ if (new net) {
+ net = current->nsproxy->net ns;
+ current->nsproxy->net_ns = new_net;
+ new net = net;
+ }
  task_unlock(current);
 if (new_nsproxy)
  put_nsproxy(new_nsproxy);
+bad_unshare_ns_cleanup_net:
+ if (new net)
+ put_net_ns(new_net);
bad_unshare_ns_cleanup_pid:
if (new_pid)
 put_pid_ns(new_pid);
Index: 2.6.19-rc6-mm2/net/core/net namespace.c
```

```
--- 2.6.19-rc6-mm2.orig/net/core/net_namespace.c
+++ 2.6.19-rc6-mm2/net/core/net namespace.c
@ @ -18,6 +18,41 @ @ struct net_namespace init_net_ns = {
#ifdef CONFIG_NET_NS
+/*
+ * Clone a new ns copying an original net ns, setting refcount to 1
+ * @old ns: namespace to clone
+ * Return NULL on error (failure to kmalloc), new ns otherwise
+static struct net_namespace *clone_net_ns(struct net_namespace *old_ns)
+ struct net_namespace *ns;
+ ns = kmalloc(sizeof(struct net_namespace), GFP_KERNEL);
+ if (!ns)
+ return NULL;
+ kref_init(&ns->kref);
+ return ns:
+}
+
+/*
+ * unshare the current process' net namespace.
+ */
+int unshare_net_ns(unsigned long unshare_flags,
    struct net_namespace **new_net)
+{
+ if (unshare_flags & NS_NET) {
+ if (!capable(CAP_SYS_ADMIN))
+ return -EPERM;
+ *new_net = clone_net_ns(current->nsproxy->net_ns);
+ if (!*new_net)
+ return -ENOMEM;
+ }
+ return 0;
+}
int copy_net_ns(int flags, struct task_struct *tsk)
 struct net_namespace *old_ns = tsk->nsproxy->net_ns;
```

Containers mailing list

Subject: [patch -mm 17/17] user namespace: add unshare Posted by Cedric Le Goater on Tue, 05 Dec 2006 10:28:09 GMT View Forum Message <> Reply to Message

From: Cedric Le Goater <clq@fr.ibm.com> Signed-off-by: Cedric Le Goater <clg@fr.ibm.com> include/linux/user namespace.h | 11 ++++++ kernel/nsproxy.c | 26 ++++++++++++++ kernel/user namespace.c 3 files changed, 91 insertions(+), 3 deletions(-) Index: 2.6.19-rc6-mm2/kernel/nsproxy.c \_\_\_\_\_ --- 2.6.19-rc6-mm2.orig/kernel/nsproxy.c +++ 2.6.19-rc6-mm2/kernel/nsproxy.c @ @ -335,6 +335,12 @ @ static int switch ns(int id, unsigned lo new ns->net ns = ns->net ns; } + if (flags & NS\_USER) { + get\_user\_ns(ns->user\_ns); + put\_user\_ns(new\_ns->user\_ns); + new\_ns->user\_ns = ns->user\_ns; + } out ns: put\_nsproxy(ns); @@ -453,6 +459,7 @@ asmlinkage long sys unshare ns(unsigned struct ipc\_namespace \*ipc, \*new\_ipc = NULL; struct pid\_namespace \*pid, \*new\_pid = NULL; struct net\_namespace \*net, \*new\_net = NULL; + struct user\_namespace \*user, \*new\_user = NULL; unsigned long unshare\_flags = 0; /\* Return -EINVAL for all unsupported flags \*/ @ @ -484,17 +491,20 @ @ asmlinkage long sys\_unshare\_ns(unsigned goto bad unshare ns cleanup ipc; if ((err = unshare\_net\_ns(unshare\_ns\_flags, &new\_net))) goto bad\_unshare\_ns\_cleanup\_pid; + if ((err = unshare\_user\_ns(unshare\_ns\_flags, &new\_user)))

+ goto bad\_unshare\_ns\_cleanup\_net;

```
- if (new mnt || new uts || new ipc || new pid || new net) {
+ if (new_mnt || new_uts || new_ipc || new_pid || new_net || new_user) {
  old nsproxy = current->nsproxy;
  new_nsproxy = dup_namespaces(old_nsproxy);
  if (!new_nsproxy) {
  err = -ENOMEM;
  goto bad_unshare_ns_cleanup_net;
  goto bad unshare ns cleanup user;
  }
- if (new_fs || new_mnt || new_uts || new_ipc || new_pid || new_net) {
+ if (new_fs || new_mnt || new_uts || new_ipc || new_pid || new_net ||
    new_user) {
  task_lock(current);
@ @ -539,12 +549,22 @ @ asmlinkage long sys unshare ns(unsigned
  new net = net;
+ if (new_user) {
+ user = current->nsproxy->user_ns;
+ current->nsproxy->user_ns = new_user;
 new user = user;
+ }
  task unlock(current);
 if (new_nsproxy)
  put_nsproxy(new_nsproxy);
+bad_unshare_ns_cleanup_user:
+ if (new user)
+ put_user_ns(new_user);
+
bad_unshare_ns_cleanup_net:
if (new net)
 put net ns(new net);
Index: 2.6.19-rc6-mm2/include/linux/user_namespace.h
--- 2.6.19-rc6-mm2.orig/include/linux/user_namespace.h
+++ 2.6.19-rc6-mm2/include/linux/user namespace.h
@ @ -22,6 +22,8 @ @ static inline void get_user_ns(struct us
kref get(&ns->kref);
}
```

```
+extern int unshare_user_ns(unsigned long unshare_flags,
    struct user_namespace **new_user);
extern int copy_user_ns(int flags, struct task_struct *tsk);
extern void free_user_ns(struct kref *kref);
@ @ -36,6 +38,15 @ @ static inline void get_user_ns(struct us
}
+static inline int unshare_user_ns(unsigned long unshare_flags,
    struct user_namespace **new_user)
+{
+ if (unshare_flags & NS_USER)
+ return -EINVAL;
+ return 0;
+}
static inline int copy_user_ns(int flags, struct task_struct *tsk)
 return 0;
Index: 2.6.19-rc6-mm2/kernel/user_namespace.c
--- 2.6.19-rc6-mm2.orig/kernel/user_namespace.c
+++ 2.6.19-rc6-mm2/kernel/user_namespace.c
@ @ -21,6 +21,63 @ @ EXPORT_SYMBOL_GPL(init_user_ns);
#ifdef CONFIG USER NS
+/*
+ * Clone a new ns copying an original user ns, setting refcount to 1
+ * @old_ns: namespace to clone
+ * Return NULL on error (failure to kmalloc), new ns otherwise
+ */
+static struct user_namespace *clone_user_ns(struct user_namespace *old_ns)
+ struct user namespace *ns;
+ struct user_struct *new_user;
+ int n:
+
+ ns = kmalloc(sizeof(struct user_namespace), GFP_KERNEL);
+ if (!ns)
+ return NULL;
+ kref_init(&ns->kref);
+ for(n = 0; n < UIDHASH SZ; ++n)
```

```
+ INIT_LIST_HEAD(ns->uidhash_table + n);
+ /* Insert new root user. */
+ ns->root_user = alloc_uid(ns, 0);
+ if (!ns->root_user) {
+ kfree(ns);
+ return NULL;
+ }
+ /* Reset current->user with a new one */
+ new_user = alloc_uid(ns, current->uid);
+ if (!new user) {
+ free_uid(ns->root_user);
+ kfree(ns);
+ return NULL;
+ }
+ switch_uid(new_user);
+ return ns;
+}
+
+/*
+ * unshare the current process' user namespace.
+int unshare_user_ns(unsigned long unshare_flags,
    struct user_namespace **new_user)
+{
+ if (unshare flags & NS USER) {
+ if (!capable(CAP SYS ADMIN))
+ return -EPERM;
+ *new_user = clone_user_ns(current->nsproxy->user_ns);
+ if (!*new_user)
+ return -ENOMEM;
+ }
+ return 0;
+}
int copy_user_ns(int flags, struct task_struct *tsk)
 struct user_namespace *old_ns = tsk->nsproxy->user_ns;
Containers mailing list
Containers@lists.osdl.org
```

https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 11/17] user namespace: add user\_namespace ptr to vfsmount
Posted by serue on Tue, 05 Dec 2006 18:27:44 GMT

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Quoting clg@fr.ibm.com (clg@fr.ibm.com):
> From: Serge E. Hallyn <serue@us.ibm.com>
> Add user\_namespace ptr to vfsmount, and define a helper to compare it

Eric.

> to the task's user ns.

Patches 11-14 are my next iteration of the patches to tag vfsmounts with a user\_namespace ptr. Would these make the user namespaces acceptable to you?

These still leave signals unaddressed, but one can argue that signals are implicitly addressed through pidspaces. If we decide that's insufficient, we can add user\_ns checks to signal permission checks.

thanks, -serge > Signed-off-by: Serge E. Hallyn <serue@us.ibm.com> > fs/namespace.c 4 ++++ > include/linux/mount.h | 2++ > include/linux/sched.h | 12 +++++++++ > 3 files changed, 18 insertions(+) > > Index: 2.6.19-rc6-mm2/fs/namespace.c > --- 2.6.19-rc6-mm2.orig/fs/namespace.c > +++ 2.6.19-rc6-mm2/fs/namespace.c > @ @ -25,6 +25,7 @ @ > #include ux/security.h> > #include ux/mount.h> > #include ux/ramfs.h> > +#include nux/user\_namespace.h> > #include <asm/uaccess.h> > #include <asm/unistd.h> > #include "pnode.h" > @ @ -56,6 +57,8 @ @ struct vfsmount \*alloc vfsmnt(const char > struct vfsmount \*mnt = kmem\_cache\_alloc(mnt\_cache, GFP\_KERNEL); > if (mnt) { memset(mnt, 0, sizeof(struct vfsmount)); > + mnt->mnt\_user\_ns = current->nsproxy->user\_ns;

```
> + get_user_ns(mnt->mnt_user_ns);
 atomic set(&mnt->mnt count, 1);
   INIT_LIST_HEAD(&mnt->mnt_hash);
   INIT_LIST_HEAD(&mnt->mnt_child);
> @ @ -88,6 +91,7 @ @ EXPORT_SYMBOL(simple_set_mnt);
> void free_vfsmnt(struct vfsmount *mnt)
> {
> + put user ns(mnt->mnt user ns);
> kfree(mnt->mnt devname);
 kmem_cache_free(mnt_cache, mnt);
> Index: 2.6.19-rc6-mm2/include/linux/mount.h
> --- 2.6.19-rc6-mm2.orig/include/linux/mount.h
> +++ 2.6.19-rc6-mm2/include/linux/mount.h
> @ @ -21.6 +21.7 @ @ struct super block;
> struct vfsmount;
> struct dentry;
> struct mnt_namespace;
> +struct user_namespace;
> #define MNT NOSUID 0x01
> #define MNT NODEV 0x02
> @ @ -53,6 +54,7 @ @ struct vfsmount {
> struct list_head mnt_slave; /* slave list entry */
> struct vfsmount *mnt_master; /* slave is on master->mnt_slave_list */
> struct mnt namespace *mnt ns; /* containing namespace */
> + struct user namespace *mnt user ns; /* namespace for uid interpretation */
> int mnt_pinned;
> }:
>
> Index: 2.6.19-rc6-mm2/include/linux/sched.h
> --- 2.6.19-rc6-mm2.orig/include/linux/sched.h
> +++ 2.6.19-rc6-mm2/include/linux/sched.h
> @ @ -83.6 +83.8 @ @ struct sched param {
> #include ux/timer.h>
> #include ux/hrtimer.h>
> #include ux/task io accounting.h>
> +#include ux/nsproxy.h>
> +#include linux/mount.h>
> #include <asm/processor.h>
>
> @ @ -1589,6 +1591,16 @ @ extern int cond_resched(void);
> extern int cond resched lock(spinlock t * lock);
> extern int cond resched softirg(void);
```

```
> +static inline int task mnt same uid(struct task struct *tsk,
      struct vfsmount *mnt)
> +{
> + if (tsk->nsproxy == init_task.nsproxy)
> + return 1;
> + if (mnt->mnt user ns == tsk->nsproxy->user ns)
> + return 1;
> + return 0;
> +}
> +
  * Does a critical section need to be broken due to another
  * task waiting?:
> --
Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers
Subject: Re: [patch -mm 05/17] ipc namespace : externalizes unshare_ipcs
Posted by Herbert Poetzl on Tue, 05 Dec 2006 23:32:11 GMT
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On Tue, Dec 05, 2006 at 11:27:57AM +0100, clg@fr.ibm.com wrote:
> From: Cedric Le Goater <clg@fr.ibm.com>
> This patch externalizes unshare ipcs when CONFIG IPC NS is not
> set. This is needed by the unshare ns syscall. Unfortunately, it could
> not be handled with a dummy static inline in ipc.h because of
> header dependencies on sched.h.
hmm, maybe sched.h should be further broken down ...
best.
Herbert
> Signed-off-by: Cedric Le Goater <clg@fr.ibm.com>
> include/linux/ipc.h |
                       2 +-
> kernel/fork.c
                  1 2 +-
> 2 files changed, 2 insertions(+), 2 deletions(-)
> Index: 2.6.19-rc6-mm2/include/linux/ipc.h
```

>

```
> --- 2.6.19-rc6-mm2.orig/include/linux/ipc.h
> +++ 2.6.19-rc6-mm2/include/linux/ipc.h
> @ @ -99,7 +99,6 @ @ extern struct ipc_namespace init_ipc_ns;
> #ifdef CONFIG IPC NS
> extern void free_ipc_ns(struct kref *kref);
> extern int copy_ipcs(unsigned long flags, struct task_struct *tsk);
> -extern int unshare_ipcs(unsigned long flags, struct ipc_namespace **ns);
> #else
> static inline int copy ipcs(unsigned long flags, struct task struct *tsk)
> @ @ -107,6 +106,7 @ @ static inline int copy ipcs(unsigned lon
> #endif
>
> +extern int unshare_ipcs(unsigned long flags, struct ipc_namespace **ns);
> static inline struct ipc_namespace *get_ipc_ns(struct ipc_namespace *ns)
> #ifdef CONFIG IPC NS
> Index: 2.6.19-rc6-mm2/kernel/fork.c
> --- 2.6.19-rc6-mm2.orig/kernel/fork.c
> +++ 2.6.19-rc6-mm2/kernel/fork.c
> @ @ -1591,7 +1591,7 @ @ static int unshare_semundo(unsigned long
> }
>
> #ifndef CONFIG_IPC_NS
> -static inline int unshare_ipcs(unsigned long flags, struct ipc_namespace **ns)
> +extern int unshare ipcs(unsigned long flags, struct ipc namespace **ns)
> {
> if (flags & CLONE_NEWIPC)
  return -EINVAL:
>
> Containers mailing list
> Containers@lists.osdl.org
> https://lists.osdl.org/mailman/listinfo/containers
Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers
```

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Herbert Poetzl on Tue, 05 Dec 2006 23:36:03 GMT

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On Tue, Dec 05, 2006 at 11:28:00AM +0100, clg@fr.ibm.com wrote:

```
> From: Cedric Le Goater <clg@fr.ibm.com>
> This patch adds a hashtable of nsproxy using the nsproxy as a key.
> init_nsproxy is hashed at init with key 0. This is considered to be
> the 'host' nsproxy.
hmm? how is that going to work?
I mean, a simple clone() call (on the host) will
change the nsproxy and spawn new ones, but that
doesn't mean that the process has left the 'host'
it just means that some namespace is not shared
with the 'other' processes ...
best.
Herbert
> Signed-off-by: Cedric Le Goater <clg@fr.ibm.com>
> include/linux/nsproxy.h | 3 +++
> kernel/nsproxy.c
                     > 2 files changed, 46 insertions(+)
>
> Index: 2.6.19-rc6-mm2/include/linux/nsproxy.h
> --- 2.6.19-rc6-mm2.orig/include/linux/nsproxy.h
> +++ 2.6.19-rc6-mm2/include/linux/nsproxy.h
> @ @ -2,6 +2,7 @ @
> #define _LINUX_NSPROXY_H
> #include ux/spinlock.h>
> +#include ux/list.h>
> struct task_struct;
>
> @ @ -34.6 +35.8 @ @ struct nsproxy {
> struct pid_namespace *pid_ns;
> struct net namespace *net ns;
  struct user_namespace *user_ns;
> + struct hlist node ns hash node;
> extern struct nsproxy init_nsproxy;
>
> Index: 2.6.19-rc6-mm2/kernel/nsproxy.c
> --- 2.6.19-rc6-mm2.orig/kernel/nsproxy.c
> +++ 2.6.19-rc6-mm2/kernel/nsproxy.c
```

```
> @ @ -22,6 +22,15 @ @
> #include linux/pid namespace.h>
> #include <linux/net_namespace.h>
> +#define NS_HASH_BITS 3 /* this might need some configuration */
> +#define NS_HASH_SIZE (1 << NS_HASH_BITS)
> +#define NS HASH MASK (NS HASH SIZE - 1)
> +#define ns_hashfn(id) (((id >> NS_HASH_BITS) + id) & NS_HASH_MASK)
> +#define ns hash head(id) &ns hash[ns hashfn(id)]
> +
> +static struct hlist_head ns_hash[NS_HASH_SIZE];
> +static DEFINE SPINLOCK(ns hash lock);
> +
> struct nsproxy init_nsproxy = INIT_NSPROXY(init_nsproxy);
> static inline void get_nsproxy(struct nsproxy *ns)
> @ @ -190.3 +199.37 @ @ void put nsproxy(struct nsproxy *ns)
  free_nsproxy(ns);
  }
>
> }
> +
> +/*
> + * This routine must be called with the ns_hash spin_locked
> +static inline struct nsproxy *ns_hash_find(int id)
> +{
> + struct hlist_node *elem;
> + struct nsproxy *ns;
> +
> + hlist_for_each_entry(ns, elem, ns_hash_head(id), ns_hash_node) {
> + if (ns->id == id) {
> + get_nsproxy(ns);
> + return ns;
> + }
> + }
> + return NULL;
> +}
> +static int init nshash init(void)
> +{
> + int i;
> + for (i = 0; i < NS_HASH_SIZE; ++i)
> + INIT_HLIST_HEAD(&ns_hash[i]);
> + spin lock irg(&ns hash lock);
> + hlist add head(&init nsproxy.ns hash node, ns hash head(0));
```

Subject: Re: [patch -mm 09/17] nsproxy: add namespace flags Posted by Herbert Poetzl on Tue, 05 Dec 2006 23:41:58 GMT

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```
On Tue, Dec 05, 2006 at 11:28:01AM +0100, clg@fr.ibm.com wrote:
> From: Cedric Le Goater <clg@fr.ibm.com>
> Signed-off-by: Cedric Le Goater <clg@fr.ibm.com>
> ---
> include/linux/nsproxy.h | 11 ++++++++
 1 file changed, 11 insertions(+)
>
> Index: 2.6.19-rc6-mm2/include/linux/nsproxy.h
> --- 2.6.19-rc6-mm2.orig/include/linux/nsproxy.h
> +++ 2.6.19-rc6-mm2/include/linux/nsproxy.h
> @ @ -14,6 +14,17 @ @ struct net namespace;
> struct user_namespace;
>
> /*
> + * namespaces flags
> +#define NS MNT 0x00000001
> +#define NS UTS 0x00000002
> +#define NS IPC 0x00000004
> +#define NS PID 0x00000008
> +#define NS_NET 0x00000010
> +#define NS_USER 0x00000020
> +#define NS_ALL (NS_MNT|NS_UTS|NS_IPC|NS_PID|NS_NET|NS_USER)
```

hmm, why \_another\_ set of flags to refer to the namespaces? is the clone()/unshare() set of flags not sufficient for that? if so, shouldn't we switch (or even better change? the unshare() too) to a new set of syscalls?

note: even if they are just for internal purpose and will never get exposed to userspace, we should think twice before we create just another set of flags, and if we do so, please let us change them all, including certain clone flags (and add a single compatibility wrapper for the 'old' syscalls)

best, Herbert

> +/\*

> \* A structure to contain pointers to all per-process

> \* namespaces - fs (mount), uts, network, sysvipc, etc.

> \*

>

> -

>\_\_\_\_\_

> Containers mailing list

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Containers mailing list
Containers@lists.osdl.org
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Subject: Re: [patch -mm 05/17] ipc namespace: externalizes unshare\_ipcs Posted by Cedric Le Goater on Thu, 07 Dec 2006 09:00:16 GMT View Forum Message <> Reply to Message

## Herbert Poetzl wrote:

> On Tue, Dec 05, 2006 at 11:27:57AM +0100, clg@fr.ibm.com wrote:

>> From: Cedric Le Goater <clg@fr.ibm.com>

>>

>> This patch externalizes unshare\_ipcs when CONFIG\_IPC\_NS is not

>> set. This is needed by the unshare\_ns syscall. Unfortunately, it could

>> not be handled with a dummy static inline in ipc.h because of

>> header dependencies on sched.h.

>

> hmm, maybe sched.h should be further broken down ...

that will require a few thousands patches:) but yes, sched.h is just so full of everything that it breaks the kernel build real often.

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Cedric Le Goater on Thu, 07 Dec 2006 09:06:03 GMT View Forum Message <> Reply to Message

- >> This patch adds a hashtable of nsproxy using the nsproxy as a key.
- >> init\_nsproxy is hashed at init with key 0. This is considered to be
- >> the 'host' nsproxy.

> hmm? how is that going to work?

>

- > I mean, a simple clone() call (on the host) will
- > change the nsproxy and spawn new ones, but that
- > doesn't mean that the process has left the 'host'
- > it just means that some namespace is not shared
- > with the 'other' processes ...

right. the new nsproxy is built with the new and old namespaces and the id can be identify it.

nsproxy 0 is init\_nsproxy which aggregates all the initial namespaces.

C.

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 09/17] nsproxy: add namespace flags Posted by Cedric Le Goater on Thu, 07 Dec 2006 09:29:28 GMT

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```
>> /*
```

>> + \* namespaces flags

>> + \*/

- >> +#define NS MNT 0x00000001
- >> +#define NS UTS 0x00000002
- >> +#define NS\_IPC 0x00000004

```
>> +#define NS PID 0x00000008
>> +#define NS NET 0x00000010
>> +#define NS_USER 0x00000020
>> +#define NS_ALL (NS_MNT|NS_UTS|NS_IPC|NS_PID|NS_NET|NS_USER)
> hmm, why _another_ set of flags to refer to the
> namespaces?
well, because namespaces are a new kind in the kernel
> is the clone()/unshare() set of flags not sufficient
> for that?
because we are reaching the limits of the CLONE_ flags.
> if so, shouldn't we switch (or even better change?
> the unshare() too) to a new set of syscalls?
unshare ns() is a new syscall and we don't really need a
clone anyway. nop?
we could make the clone flags and namespace flags compatible
with:
#define NS_MNT_CLONE_NEWNS
#define NS UTS CLONE NEWUTS
#define NS_IPC CLONE_NEWIPC
that shouldn't be a big issue but we could also
remove/deprecate:
#define CLONE_NEWNS 0x00020000 /* New namespace group? */
#define CLONE_NEWUTS 0x04000000 /* New utsname group? */
#define CLONE_NEWIPC 0x08000000 /* New ipcs */
to make sure namespaces can not be unshared using
unshare()
> note: even if they are just for internal purpose
> and will never get exposed to userspace,
they will. unshare_ns() is a syscall.
> we should think twice before we create just another
> set of flags, and if we do so, please let us change
> them all, including certain clone flags (and add a
```

> single compatibility wrapper for the 'old' syscalls)

compatible with the new ones.

something like this:

int sys\_unshare(int unshare\_flags)
{
 int unshare\_ns\_flags;

 unshare\_ns\_flags = convert\_flags(unshare\_flags);

 return sys\_unshare\_ns(unshare\_ns\_flags);
}

C.

Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers

so you would keep the unshare as is but change the set of flags its using, making sure the old ones are still

Subject: Re: [patch -mm 10/17] nsproxy: add unshare\_ns and bind\_ns syscalls Posted by ebiederm on Fri, 08 Dec 2006 19:26:49 GMT

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## NAK

Don't use global identifiers. Use pids. i.e. struct pid \* for your identifiers. Is there is a reason pids are unsuitable?

I'm also worried about the security implications of switching namespaces on a process. That is something that needs to be looked at very closely.

These two changes certainly don't belong in a single patch, and they certainly use a bit more explanation. syscalls are not something to add lightly. Because they must be supported forever.

Eric

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by ebiederm on Fri, 08 Dec 2006 19:30:31 GMT View Forum Message <> Reply to Message

clg@fr.ibm.com writes:

- > From: Cedric Le Goater <clg@fr.ibm.com>
- >
- > This patch adds a hashtable of nsproxy using the nsproxy as a key.
- > init nsproxy is hashed at init with key 0. This is considered to be
- > the 'host' nsproxy.

NAK. Which namespace do these ids live in?

It sounds like you are setting up to make the 'host' nsproxy special and have special rules. That also sounds wrong.

Even letting the concept of nsproxy escape to user space sounds wrong. nsproxy is an internal space optimization. It's not struct container and I don't think we want it to become that.

Eric

\_\_\_\_\_

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

## Subject: Re: [patch -mm 09/17] nsproxy: add namespace flags Posted by ebiederm on Fri, 08 Dec 2006 19:40:33 GMT

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Cedric Le Goater <clg@fr.ibm.com> writes:

```
>>> /*
>>> + * namespaces flags
>>> + */
>>> +#define NS MNT 0x00000001
>>> +#define NS_UTS 0x00000002
>>> +#define NS IPC 0x00000004
>>> +#define NS PID 0x00000008
>>> +#define NS NET 0x00000010
>>> +#define NS USER 0x00000020
>>> +#define NS ALL (NS MNT|NS UTS|NS IPC|NS PID|NS NET|NS USER)
>> hmm, why _another_ set of flags to refer to the
>> namespaces?
> well, because namespaces are a new kind in the kernel
Gratuitous incompatibility.
>> is the clone()/unshare() set of flags not sufficient
>> for that?
> because we are reaching the limits of the CLONE_ flags.
Not really. There are at least 8 bits that clone cannot use
but that unshare can.
>> if so, shouldn't we switch (or even better change?
>> the unshare() too) to a new set of syscalls?
> unshare ns() is a new syscall and we don't really need a
> clone anyway. nop?
Huh? Clone should be the primary. There are certain namespaces
that it are very hard to unshare, without creating a new process.
> we could make the clone flags and namespace flags compatible
> with:
> #define NS MNT CLONE NEWNS
> #define NS UTS CLONE NEWUTS
> #define NS_IPC CLONE_NEWIPC
> that shouldn't be a big issue but we could also
```

```
> remove/deprecate:
> #define CLONE_NEWNS 0x00020000 /* New namespace group? */
> #define CLONE_NEWUTS 0x04000000 /* New utsname group? */
> #define CLONE_NEWIPC 0x08000000 /* New ipcs */
> to make sure namespaces can not be unshared using
> unshare()
Nope. We got them now they can not be removed. It's part of the ABI.
And since we are not removing the functionality it makes not sense
at all to just change the name of the flags.
>> we should think twice before we create just another
>> set of flags, and if we do so, please let us change
>> them all, including certain clone flags (and add a
>> single compatibility wrapper for the 'old' syscalls)
> so you would keep the unshare as is but change the set
> of flags its using, making sure the old ones are still
> compatible with the new ones.
>
> something like this:
> int sys_unshare(int unshare_flags)
> {
> int unshare_ns_flags;
> unshare ns flags = convert flags(unshare flags);
> return sys unshare ns(unshare ns flags);
> }
>?
If necessary.
Eric
Containers mailing list
Containers@lists.osdl.org
```

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by serue on Fri, 08 Dec 2006 19:53:06 GMT

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Quoting Eric W. Biederman (ebiederm@xmission.com): > clg@fr.ibm.com writes: > > From: Cedric Le Goater <clg@fr.ibm.com> >> This patch adds a hashtable of nsproxy using the nsproxy as a key. >> init nsproxy is hashed at init with key 0. This is considered to be > > the 'host' nsproxy. > NAK. Which namespace do these ids live in? > It sounds like you are setting up to make the 'host' nsproxy special > and have special rules. That also sounds wrong. > Even letting the concept of nsproxy escape to user space sounds wrong. > nsproxy is an internal space optimization. It's not struct container > and I don't think we want it to become that. > Eric

So would you advocate referring to containers just by the pid of a process containing the nsproxy, and letting userspace maintain a mapping of id's to containers through container create/enter commands? Or is there some other way you were thinking of doing this?

thanks. -serge

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 04/17] nsproxy: externalizes exit\_task\_namespaces Posted by ebiederm on Fri, 08 Dec 2006 20:16:22 GMT

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clg@fr.ibm.com writes:

- > From: Cedric Le Goater <clg@fr.ibm.com>
- > this is required to remove a header dependency in sched.h which breaks > next patches.

This just doesn't feel right.

Why with unshare working now are you needing to rework everything?

Page 72 of 124 ---- Generated from OpenVZ Forum

Containers mailing list
Containers@lists.osdl.org

https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by ebiederm on Fri, 08 Dec 2006 20:57:38 GMT View Forum Message <> Reply to Message

"Serge E. Hallyn" <serue@us.ibm.com> writes:

- > Quoting Eric W. Biederman (ebiederm@xmission.com):
- >> clg@fr.ibm.com writes:

>>

>> > From: Cedric Le Goater <clg@fr.ibm.com>

>> >

- >> > This patch adds a hashtable of nsproxy using the nsproxy as a key.
- >> > init\_nsproxy is hashed at init with key 0. This is considered to be
- >> > the 'host' nsproxy.

>>

>> NAK. Which namespace do these ids live in?

>>

- >> It sounds like you are setting up to make the 'host' nsproxy special
- >> and have special rules. That also sounds wrong.

>>

- >> Even letting the concept of nsproxy escape to user space sounds wrong.
- >> nsproxy is an internal space optimization. It's not struct container
- >> and I don't think we want it to become that.

>>

>> Eric

>

- > So would you advocate referring to containers just by the pid of a
- > process containing the nsproxy, and letting userspace maintain a mapping
- > of id's to containers through container create/enter commands? Or is
- > there some other way you were thinking of doing this?

There are two possible ways.

- 1) Just use a process using the namespace. This is easiest to implement.
- 2) Have a struct pid reference in the namespace itself, and probably an extra pointer in struct pid to find it.

This is the most stable, because fork/exit won't affect which pid you need to use.

Beyond that yes it seems to make sense to let user space maintain any mapping

of containers to ids.

Eric

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Herbert Poetzl on Sat, 09 Dec 2006 04:11:10 GMT View Forum Message <> Reply to Message

On Fri, Dec 08, 2006 at 01:57:38PM -0700, Eric W. Biederman wrote: > "Serge E. Hallyn" <serue@us.ibm.com> writes: > > Quoting Eric W. Biederman (ebiederm@xmission.com): >>> clg@fr.ibm.com writes: >>> From: Cedric Le Goater <clg@fr.ibm.com> > >> > >>> This patch adds a hashtable of nsproxy using the nsproxy as a key. >>> > init nsproxy is hashed at init with key 0. This is considered to be >>> the 'host' nsproxy. > >> NAK. Which namespace do these ids live in? well, I gave a similar answer in another email, so I fully agree with the NAK here ... >>> It sounds like you are setting up to make the 'host' nsproxy >>> special and have special rules. That also sounds wrong. > >> >>> Even letting the concept of nsproxy escape to user space sounds >>> wrong, nsproxy is an internal space optimization. It's not struct >>> container and I don't think we want it to become that. > >> > >> Eric > > > > So would you advocate referring to containers just by the pid of > > a process containing the nsproxy, and letting userspace maintain > > a mapping of id's to containers through container create/enter >> commands? Or is there some other way you were thinking of doing > > this? > There are two possible ways.

- > 1) Just use a process using the namespace.
- > This is easiest to implement.

- > 2) Have a struct pid reference in the namespace itself,
- > and probably an extra pointer in struct pid to find it.
- > This is the most stable, because fork/exit won't affect
- > which pid you need to use.

while I agree that nsproxy is definitely the wrong point to tie a 'context' too, as it can contain a mixture of spaces from inside and outside a context, and it would require to forbid doing things like clone() with the space flags, both inside and outside a 'container' to allow to use them for actual vps applications, I think that we have to have some kind of handle to tie specific sets of namespaces too

that 'can' be an nsproxy or something different, but I'm absolutely unhappy with tying it to a process, as I already mentioned several times, that lightweight 'containers' do not use/have an init process, and no single process might survive the entire life span of that 'container' ...

- > Beyond that yes it seems to make sense to let user space
- > maintain any mapping of containers to ids.

I agree with that, but we need something to move around between the various spaces ...

for example, Linux-VServer ties the namespaces to the context structure (atm) which allows userspace to set and enter specific spaces of a guest context (I assume OpenVZ does similar)

HTC, Herbert

- > Eric
- >
- > Containers mailing list
- > Containers@lists.osdl.org
- > https://lists.osdl.org/mailman/listinfo/containers

Containers mailing list
Containers@lists.osdl.org

https://lists.osdl.org/mailman/listinfo/containers

```
On Fri, Dec 08, 2006 at 12:26:49PM -0700, Eric W. Biederman wrote:
> clg@fr.ibm.com writes:
> > From: Cedric Le Goater <clg@fr.ibm.com>
>> The following patch defines 2 new syscalls specific to nsproxy and
> > namespaces :
> > * unshare_ns:
>> enables a process to unshare one or more namespaces. this
        duplicates the unshare syscall for the moment but we
>> expect to diverge when the number of namespaces increases
> Are we out of clone flags yet? If not this is premature.
no, but a different nevertheless related question:
does anybody, except for 'us' use the unshare() syscall?
because if not, then why not simply extend that one
to 64bit and be done, we probably won't need a clone64()
but if we find we do (at some point) adding that with
the new flags would be trivial ...
OTOH, we could also just add an unshare64() too
anyway, we _will_ run out of flags in the near future
> > * bind ns:
> >
>> allows a process to bind
>> 1 - its nsproxy to some identifier
>> 2 - to another nsproxy using an identifier or -pid
>
> NAK
> Don't use global identifiers. Use pids. i.e. struct pid * for your
> identifiers. Is there is a reason pids are unsuitable?
yes, see my reply in the other mail
> I'm also worried about the security implications of switching
> namespaces on a process.
> That is something that needs to be looked at very closely.
```

Linux-VServer currently uses a capability to prevent changing between namespaces (a very generic one) but it probably makes sense to add something like that in general ... btw, did I mention that the capability flags are running out too?

- > These two changes certainly don't belong in a single patch,
- > and they certainly use a bit more explanation.
- > syscalls are not something to add lightly.

well, and they will take ages to get into mainline for all archs, or has that changed sine we reserved sys\_vserver()?

> Because they must be supported forever.

I'm not sure about that, most archs 'reuse' syscalls when there is no user left ...

best, Herbert

> Eric

>\_\_\_\_\_

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Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 10/17] nsproxy: add unshare\_ns and bind\_ns syscalls Posted by ebiederm on Sat, 09 Dec 2006 07:40:03 GMT View Forum Message <> Reply to Message

Herbert Poetzl <a href="mailto:herbert@13thfloor.at">herbert@13thfloor.at</a> writes:

- > On Fri, Dec 08, 2006 at 12:26:49PM -0700, Eric W. Biederman wrote:
- >> clg@fr.ibm.com writes:

>>

>> > From: Cedric Le Goater <clg@fr.ibm.com>

>> >

- >> > The following patch defines 2 new syscalls specific to nsproxy and
- >> > namespaces:

>> >

>> > \* unshare\_ns :

```
>> >
>> > enables a process to unshare one or more namespaces. this
>> > duplicates the unshare syscall for the moment but we
>> > expect to diverge when the number of namespaces increases
>>
>> Are we out of clone flags yet? If not this is premature.
> > no, but a different nevertheless related question:
> does anybody, except for 'us' use the unshare() syscall?
```

The pam\_namespace module if I have looked at things properly. I believe that is what it was added to support.

```
because if not, then why not simply extend that one
to 64bit and be done, we probably won't need a clone64()
but if we find we do (at some point) adding that with
the new flags would be trivial ...
OTOH, we could also just add an unshare64() too
anyway, we _will_ run out of flags in the near future
```

Agreed. Please let's cross that bridge when we come to it.

>> I'm also worried about the security implications of switching
>> namespaces on a process.
>> That is something that needs to be looked at very closely.
>
Linux-VServer currently uses a capability to prevent
> changing between namespaces (a very generic one) but
> it probably makes sense to add something like that
> in general ... btw, did I mention that the capability
> flags are running out too?

I think they have run out. Not that sys\_capability needs a revision but it appears the format of the data does which is likely just as bad.

```
>> These two changes certainly don't belong in a single patch,
>> and they certainly use a bit more explanation.
>> syscalls are not something to add lightly.
>
> well, and they will take ages to get into mainline
> for all archs, or has that changed sine we reserved
> sys_vserver()?
```

I think it is likely a little better. I'm not certain what your definition of ages is.

>> Because they must be supported forever.

>

- > I'm not sure about that, most archs 'reuse' syscalls
- > when there is no user left ...

I haven't seen that on i386. Except for experimental syscalls I have a hard time believing we have any syscalls that have had all of their users disappear.

Reusing syscall numbers is in a lot of ways completely irresponsible once you start supporting a binary interface. Even if you remove the syscall because there are no users or it makes absolutely no sense any more.

Fric

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by ebiederm on Sat, 09 Dec 2006 07:54:26 GMT View Forum Message <> Reply to Message

Herbert Poetzl <a href="mailto:herbert@13thfloor.at">herbert@13thfloor.at</a> writes:

- >> There are two possible ways.
- >> 1) Just use a process using the namespace.
- >> This is easiest to implement.

>

- >> 2) Have a struct pid reference in the namespace itself,
- >> and probably an extra pointer in struct pid to find it.
- >> This is the most stable, because fork/exit won't affect
- >> which pid you need to use.

>

- > that 'can' be an nsproxy or something different, but
- > I'm absolutely unhappy with tying it to a process,
- > as I already mentioned several times, that lightweight
- > 'containers' do not use/have an init process, and no
- > single process might survive the entire life span of
- > that 'container' ...

Herbert think of a session id. That is a pid that is tied to something besides a single process.

It is easy and recursion safe to tie a pid to a namespace

or anything else that make sense, as I suggested above.

The pid namespace feels like the right place for this kind of activity.

- >> Beyond that yes it seems to make sense to let user space
- >> maintain any mapping of containers to ids.

>

- > I agree with that, but we need something to move
- > around between the various spaces ...

If you have CAP\_SYS\_PTRACE or you have a child process in a container you can create another with ptrace.

Now I don't mind optimizing that case, with something like the proposed bind\_ns syscall. But we need to be darn certain why it is safe, and does not change the security model that we currently have.

I have not seen that discussion yet, and until I do I have serious concerns. That discussion needs to be on lkml as well. Why did Al Viro think this was a bad idea when it was proposed for the mount namespace?

This is where you are on the edge of some very weird interface interactions. Without suid programs it would be completely safe for anyone to unshare their mount namespace. With suid programs allowed an unprivileged unshare mount namespace unshare is next to impossible.

- > for example, Linux-VServer ties the namespaces to
- > the context structure (atm) which allows userspace
- > to set and enter specific spaces of a guest context
- > (I assume OpenVZ does similar)

Yep, and we certainly need to find a way to fulfill this usage requirement.

Eric

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Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 10/17] nsproxy: add unshare\_ns and bind\_ns syscalls Posted by Herbert Poetzl on Sat, 09 Dec 2006 08:22:02 GMT

```
On Sat, Dec 09, 2006 at 12:40:03AM -0700, Eric W. Biederman wrote:
> Herbert Poetzl <herbert@13thfloor.at> writes:
> > On Fri, Dec 08, 2006 at 12:26:49PM -0700, Eric W. Biederman wrote:
>>> clg@fr.ibm.com writes:
>>> From: Cedric Le Goater <clg@fr.ibm.com>
>>> The following patch defines 2 new syscalls specific to nsproxy and
>>> > namespaces:
> >> >
>>> * unshare ns:
>>> enables a process to unshare one or more namespaces. this
           duplicates the unshare syscall for the moment but we
>>> expect to diverge when the number of namespaces increases
> >>
>>> Are we out of clone flags yet? If not this is premature.
> > no, but a different nevertheless related question:
> > does anybody, except for 'us' use the unshare() syscall?
> The pam_namespace module if I have looked at things properly.
> I believe that is what it was added to support.
hmm, too bad then ...
> > because if not, then why not simply extend that one
>> to 64bit and be done, we probably won't need a clone64()
> > but if we find we do (at some point) adding that with
> > the new flags would be trivial ...
> >
> > OTOH, we could also just add an unshare64() too
> > anyway, we _will_ run out of flags in the near future
> Agreed. Please let's cross that bridge when we come to it.
well, personally I'd prefer to 'know' that the
interfaces we introduce now to handle the 'new'
namespaces will still work in a few months (or
maybe years?) and not require _another_ change
let me give an example here:
```

Page 81 of 124 ---- Generated from OpenVZ Forum

we now change our userspace tools to support new clone() flags, we can do that with a little support from the kernel quite easily without changing the tools with every kernel release and more important, without breaking backwards compatibility in the Linux-VServer ABI (and API)

if we decide to switch to a completely different API in, lets say six month from now, 'we' have two options:

- add \_another\_ compatibility layer to handle the 'old' (i.e. now introduced) ABI
- accept that older tools will fail and/or produce strange results because 'our' ABI isn't supported anymore ...

so naturally, I'm not very excited to introduce and/or utilize and interface which we all \_know\_ is not able to satisfy the upcoming demands ...

nevertheless, that isn't considered a major issue here, so feel free to ignore my personal feelings

- >>> I'm also worried about the security implications of switching
- >>> namespaces on a process.
- >>> That is something that needs to be looked at very closely.

> >

- > > Linux-VServer currently uses a capability to prevent
- > > changing between namespaces (a very generic one) but
- > > it probably makes sense to add something like that
- > > in general ... btw, did I mention that the capability
- > > flags are running out too?

>

- > I think they have run out. Not that sys\_capability needs a
- > revision but it appears the format of the data does which
- > is likely just as bad.

gladly not in 2.6.19, as we absolutely need to add one capability (CAP\_CONTEXT) in Linux-VServer (the one allowing to utilize sys\_vserver)

but once that is taken by mainline, we have to add another workaround ...

sidenote: I made several suggestions to extend the capability system (in number of caps and functionality) of which none was even considered

>>> These two changes certainly don't belong in a single patch,

> >> and they certainly use a bit more explanation. >>> syscalls are not something to add lightly. > > well, and they will take ages to get into mainline > > for all archs, or has that changed sine we reserved > > sys\_vserver()? > > I think it is likely a little better. I'm not certain what > your definition of ages is. took slightly more than a year, IIRC ... >>> Because they must be supported forever. > > > > I'm not sure about that, most archs 'reuse' syscalls > > when there is no user left ... > I haven't seen that on i386. Except for experimental > syscalls I have a hard time believing we have any syscalls > that have had all of their users disappear. okay, so be it ... > Reusing syscall numbers is in a lot of ways completely > irresponsible once you start supporting a binary interface. > Even if you remove the syscall because there are no users > or it makes absolutely no sense any more. okay, no problem with keeping them around ... best, Herbert > Eric Containers mailing list

Subject: Re: [patch -mm 10/17] nsproxy: add unshare\_ns and bind\_ns syscalls Posted by Cedric Le Goater on Mon, 11 Dec 2006 15:21:05 GMT View Forum Message <> Reply to Message

Eric W. Biederman wrote:

Containers@lists.osdl.org

> clg@fr.ibm.com writes:

\_

>> From: Cedric Le Goater <clg@fr.ibm.com>

https://lists.osdl.org/mailman/listinfo/containers

>>
>> The following patch defines 2 new syscalls specific to nsproxy and
>> namespaces :
>>
>> * unshare_ns :
>> anables a present to unabare one or mare names as this
<ul><li>&gt;&gt; enables a process to unshare one or more namespaces. this</li><li>&gt;&gt; duplicates the unshare syscall for the moment but we</li></ul>
<ul><li>outpilcates the unshare syscall for the moment but we</li><li>expect to diverge when the number of namespaces increases</li></ul>
> expect to diverge when the humber of hamespaces increases
> Are we out of clone flags yet? If not this is premature.
>
>> * bind_ns :
_ >>
>> allows a process to bind
>> 1 - its nsproxy to some identifier
>> 2 - to another nsproxy using an identifier or -pid
>
> NAK
> Doubt was alabal identificant. Her mide it a structurid * for your
> Don't use global identifiers. Use pids. i.e. struct pid * for your
> identifiers. Is there is a reason pids are unsuitable?
(1) gives a little more freedom to the sysadmin managing its
(2) uses pids. do you also nak it?
(,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,
do you always have access to pid?
> I'm also worried about the security implications of switching namespaces
> on a process. That is something that needs to be looked at very closely.
this is required by at least 3 products I know of.
> These two changes certainly don't belong in a single patch, and they
> certainly use a bit more explanation. syscalls are not something to
> add lightly. Because they must be supported forever.
r and lightly. Boodado they much be supported forever.
agree.
C.
Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers

## Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Cedric Le Goater on Mon, 11 Dec 2006 15:23:29 GMT

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Eric W. Biederman wrote:

> clg@fr.ibm.com writes:

>

>> From: Cedric Le Goater <clg@fr.ibm.com>

>>

- >> This patch adds a hashtable of nsproxy using the nsproxy as a key.
- >> init\_nsproxy is hashed at init with key 0. This is considered to be
- >> the 'host' nsproxy.

>

> NAK. Which namespace do these ids live in?

>

- > It sounds like you are setting up to make the 'host' nsproxy special
- > and have special rules.

exactly.

> That also sounds wrong.

sounds very nice to me and a few others.

- > Even letting the concept of nsproxy escape to user space sounds wrong.
- > nsproxy is an internal space optimization. It's not struct container
- > and I don't think we want it to become that.

i don't agree here. we need that, so does openvz, vserver, people working on resource management.

C.

\_\_\_\_\_

Containers mailing list
Containers@lists.osdl.org

https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 04/17] nsproxy: externalizes exit\_task\_namespaces Posted by Cedric Le Goater on Mon, 11 Dec 2006 15:26:57 GMT

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Eric W. Biederman wrote:

- > clg@fr.ibm.com writes:
- >
- >> From: Cedric Le Goater <clg@fr.ibm.com>
- >>
- >> this is required to remove a header dependency in sched.h which breaks

```
>> next patches.
> This just doesn't feel right.
> Why with unshare working now are you needing to rework everything?
This is not everything, this is just 3 lines compared to millions.
The issue is not here anyway.
C.
Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers
Subject: Re: [patch -mm 09/17] nsproxy: add namespace flags
Posted by Cedric Le Goater on Mon, 11 Dec 2006 15:27:16 GMT
View Forum Message <> Reply to Message
Eric W. Biederman wrote:
> Cedric Le Goater <clq@fr.ibm.com> writes:
>
>>> /*
>>> + * namespaces flags
>>>> + */
>>> +#define NS MNT 0x00000001
>>> +#define NS_UTS 0x00000002
>>> +#define NS IPC 0x00000004
>>> +#define NS PID 0x00000008
>>> +#define NS NET 0x00000010
>>> +#define NS USER 0x00000020
>>> +#define NS_ALL (NS_MNT|NS_UTS|NS_IPC|NS_PID|NS_NET|NS_USER)
>>> hmm, why another set of flags to refer to the
>>> namespaces?
>> well, because namespaces are a new kind in the kernel
```

?
>>> is the clone()/unshare() set of flags not sufficient
>>> for that?
>> because we are reaching the limits of the CLONE\_ flags.
>
> Not really. There are at least 8 bits that clone cannot use
> but that unshare can.

> Gratuitous incompatibility.

```
please, could you list them?
>>> if so, shouldn't we switch (or even better change?
>>> the unshare() too) to a new set of syscalls?
>> unshare_ns() is a new syscall and we don't really need a
>> clone anyway. nop?
> Huh? Clone should be the primary. There are certain namespaces
> that it are very hard to unshare, without creating a new process.
You just said above that clone had less available flags than
unshare ...
anyway, could you elaborate a bit more? I have the opposite
feeling and you gave me that impression also a few month ago.
No problem for me, i just want a way to use this stuff without
>>> we should think twice before we create just another
>>> set of flags, and if we do so, please let us change
>>> them all, including certain clone flags (and add a
>>> single compatibility wrapper for the 'old' syscalls)
>> so you would keep the unshare as is but change the set
>> of flags its using, making sure the old ones are still
>> compatible with the new ones.
>> something like this:
>>
>> int sys unshare(int unshare flags)
>> {
>> int unshare_ns_flags;
>> unshare_ns_flags = convert_flags(unshare_flags);
>>
>> return sys_unshare_ns(unshare_ns_flags);
>> }
>>
>> ?
> If necessary.
ok good. will check it out.
C.
```

Containers mailing list

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by serue on Mon, 11 Dec 2006 15:29:35 GMT View Forum Message <> Reply to Message

```
Quoting Eric W. Biederman (ebiederm@xmission.com):
> Herbert Poetzl <herbert@13thfloor.at> writes:
>>> There are two possible ways.
>>> 1) Just use a process using the namespace.
>>> This is easiest to implement.
> >
>>> 2) Have a struct pid reference in the namespace itself.
>>> and probably an extra pointer in struct pid to find it.
       This is the most stable, because fork/exit won't affect
> >>
       which pid you need to use.
> >>
> > that 'can' be an nsproxy or something different, but
> > I'm absolutely unhappy with tying it to a process,
> > as I already mentioned several times, that lightweight
> > 'containers' do not use/have an init process, and no
> > single process might survive the entire life span of
> > that 'container' ...
>
> Herbert think of a session id. That is a pid that is
> tied to something besides a single process.
>
> It is easy and recursion safe to tie a pid to a namespace
```

Recursion safe, but limiting in that you can only descend one pid namespace at a time. That limitation aside, providing task notifiers for all unshares, plus a syscall to jump into all namespaces belonging to a process known to you by a particular pid, could be a good approach. Now you can have a userspace daemon keeping namespace id's tied to processes, giving you the ability to say

> or anything else that make sense, as I suggested above.

```
ns_exec -a -I ns12 my_prog
(unshare all namespaces and run my_prog in
a container known as 'ns12')
ns_enter -I ns12 /bin/ps
(jump into ns12 and run /bin/ps)
```

The likely requirement to run a namespace tracking daemon in

each pid namespace that wants such functionality could become a resource hog, but that may be just a theoretical problem, since you'll only need that if you want to play with namespaces, meaning that for it to be a problem you'd have to have lots of virtual servers each maintaining namespaces to either do process migration or spawn more virtual servers (which each maintain namespaces to...)

> The pid namespace feels like the right place for this kind > of activity. >
> >> Beyond that yes it seems to make sense to let user space > >> maintain any mapping of containers to ids. >>
> > I agree with that, but we need something to move > > around between the various spaces ... >
> If you have CAP\_SYS\_PTRACE or you have a child process > in a container you can create another with ptrace. >
Now I don't mind optimizing that case, with something like

Sigh, and that's going to have to be a discussion per namespace.

> the proposed bind\_ns syscall. But we need to be darn certain> why it is safe, and does not change the security model that

- > I have not seen that discussion yet, and until I do I have
- > serious concerns. That discussion needs to be on lkml as
- > well. Why did Al Viro think this was a bad idea when it
- > was proposed for the mount namespace?

> This is where you are on the edge of some very weird interface

- > interactions. Without suid programs it would be completely safe
- > for anyone to unshare their mount namespace. With suid programs
- > allowed an unprivileged unshare mount namespace unshare is next to
- > impossible.

>

> we currently have.

- > > for example, Linux-VServer ties the namespaces to
- > > the context structure (atm) which allows userspace
- > > to set and enter specific spaces of a guest context
- > > (I assume OpenVZ does similar)

> Yep, and we certainly need to find a way to fulfill this usage

> requirement.

Containers mailing list
Containers@lists.osdl.org

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by serue on Mon, 11 Dec 2006 15:56:22 GMT

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Quoting Serge E. Hallyn (serue@us.ibm.com):

- > Quoting Eric W. Biederman (ebiederm@xmission.com):
- > > Herbert Poetzl <herbert@13thfloor.at> writes:
- >>> Beyond that yes it seems to make sense to let user space
- >>> maintain any mapping of containers to ids.
- >>>
- >>> I agree with that, but we need something to move
- >> around between the various spaces ...
- > >
- > > If you have CAP\_SYS\_PTRACE or you have a child process
- > > in a container you can create another with ptrace.
- > >
- > > Now I don't mind optimizing that case, with something like
- >> the proposed bind\_ns syscall. But we need to be darn certain
- >> why it is safe, and does not change the security model that
- > > we currently have.

>

> Sigh, and that's going to have to be a discussion per namespace.

Well, assuming that we're using pids as identifiers, that means we can only enter decendent namespaces, which means 'we' must have created them. So anything we could do by entering the ns, we could have done by creating it as well, right?

-serge

\_\_\_\_\_

Containers mailing list Containers@lists.osdl.org

https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Cedric Le Goater on Mon, 11 Dec 2006 16:09:33 GMT View Forum Message <> Reply to Message

## Herbert Poetzl wrote:

- > On Fri, Dec 08, 2006 at 01:57:38PM -0700, Eric W. Biederman wrote:
- >> "Serge E. Hallyn" <serue@us.ibm.com> writes:

>>

>>> Quoting Eric W. Biederman (ebiederm@xmission.com):

```
>>>> clg@fr.ibm.com writes:
>>>>
>>>> From: Cedric Le Goater <clg@fr.ibm.com>
>>>>
>>>> This patch adds a hashtable of nsproxy using the nsproxy as a key.
>>>> init_nsproxy is hashed at init with key 0. This is considered to be
>>>> the 'host' nsproxy.
>>>> NAK. Which namespace do these ids live in?
> well, I gave a similar answer in another email,
> so I fully agree with the NAK here ...
hmm, I wasn't that clear to me. OK, let's dig :)
>>>> It sounds like you are setting up to make the 'host' nsproxy
>>> special and have special rules. That also sounds wrong.
>>>>
>>>> Even letting the concept of nsproxy escape to user space sounds
>>> wrong, nsproxy is an internal space optimization. It's not struct
>>> container and I don't think we want it to become that.
>>>>
>>>> Eric
>>> So would you advocate referring to containers just by the pid of
>>> a process containing the nsproxy, and letting userspace maintain
>>> a mapping of id's to containers through container create/enter
>>> commands? Or is there some other way you were thinking of doing
>>> this?
>> There are two possible ways.
>> 1) Just use a process using the namespace.
    This is easiest to implement.
>
>> 2) Have a struct pid reference in the namespace itself,
    and probably an extra pointer in struct pid to find it.
     This is the most stable, because fork/exit won't affect
>>
    which pid you need to use.
>>
> while I agree that nsproxy is definitely the wrong
> point to tie a 'context' too, as it can contain a
> mixture of spaces from inside and outside a context,
> and it would require to forbid doing things like
> clone() with the space flags, both inside and outside
> a 'container' to allow to use them for actual vps
> applications, I think that we have to have some kind
> of handle to tie specific sets of namespaces too
```

this is nsproxy ...

- > that 'can' be an nsproxy or something different, but
- > I'm absolutely unhappy with tying it to a process,

hmm, what do you mean? nsproxy survives the death of any process. It's not tied to any process in particular. One process creates it with an unshare but that's all.

the ->nsproxy in task\_struct is a way to find it.

- > as I already mentioned several times, that lightweight
- > 'containers' do not use/have an init process, and no
- > single process might survive the entire life span of
- > that 'container' ...

I think there is a misunderstanding here. a 'container' or 'nsproxy' or what ever is a set of namespaces which are not tied to a process.

you can do that today on 2.6.19 with utsname.

- >> Beyond that yes it seems to make sense to let user space
- >> maintain any mapping of containers to ids.

>

- > I agree with that, but we need something to move
- > around between the various spaces ...

the bind\_ns syscall lets the user specify the mapping. this is not done by the kernel.

I had to introduce some rules, like giving more capabilities to some processes, but that can be changed. For the moment, they have to live in "init\_proxy".

- > for example, Linux-VServer ties the namespaces to
- > the context structure (atm) which allows userspace
- > to set and enter specific spaces of a quest context
- > (I assume OpenVZ does similar)

What's the big difference with nsproxy?

C.

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

# Subject: Re: [patch -mm 10/17] nsproxy: add unshare\_ns and bind\_ns syscalls Posted by Cedric Le Goater on Mon, 11 Dec 2006 17:05:16 GMT

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Herbert Poetzl wrote: > On Fri, Dec 08, 2006 at 12:26:49PM -0700, Eric W. Biederman wrote: >> clg@fr.ibm.com writes: >>
>>> From: Cedric Le Goater <clg@fr.ibm.com></clg@fr.ibm.com>
>>> The following patch defines 2 new syscalls specific to nsproxy and >>> namespaces :
>>> * unshare_ns :
>>> enables a process to unshare one or more namespaces. this >>> duplicates the unshare syscall for the moment but we >>> expect to diverge when the number of namespaces increases >> Are we out of clone flags yet? If not this is premature.
<ul><li>no, but a different nevertheless related question:</li><li>does anybody, except for 'us' use the unshare() syscall?</li></ul>
<ul> <li>because if not, then why not simply extend that one</li> <li>to 64bit and be done, we probably won't need a clone64()</li> <li>but if we find we do (at some point) adding that with</li> <li>the new flags would be trivial</li> </ul>
> OTOH, we could also just add an unshare64() too >
> anyway, we _will_ run out of flags in the near future
yes. that's probably the way to go. I'll rework unshare_ns() in a unshare64(). it will give some air to the 32bits clone() and unshare() and will let us use the >32bits flags for namespaces.
thanks,
C.
Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by ebiederm on Mon, 11 Dec 2006 19:35:12 GMT View Forum Message <> Reply to Message

- "Serge E. Hallyn" <serue@us.ibm.com> writes:
- > Quoting Serge E. Hallyn (serue@us.ibm.com):
- >> Quoting Eric W. Biederman (ebiederm@xmission.com):
- >> > Herbert Poetzl < herbert@13thfloor.at> writes:
- >> > > Beyond that yes it seems to make sense to let user space
- >> >> maintain any mapping of containers to ids.
- >> > >
- >> > I agree with that, but we need something to move
- >> > around between the various spaces ...
- >> >
- >> > If you have CAP\_SYS\_PTRACE or you have a child process
- >> > in a container you can create another with ptrace.
- >> >
- >> > Now I don't mind optimizing that case, with something like
- >> > the proposed bind\_ns syscall. But we need to be darn certain
- >> > why it is safe, and does not change the security model that
- >> > we currently have.

>>

- >> Sigh, and that's going to have to be a discussion per namespace.
- > Well, assuming that we're using pids as identifiers, that means
- > we can only enter decendent namespaces, which means 'we' must
- > have created them. So anything we could do by entering the ns,
- > we could have done by creating it as well, right?

It isn't strict descendents who we can see. i.e. init can create the thing, and we could have just logged into the network but init and us still share the same pid namespace.

But yes it would be we can only enter descendent namespaces, for some definition of enter.

There are two issues.

- We may have a namespace we want to create and then remove the ability for the sysadmin to fiddle with, so it can play with encrypted data or something like that safely. Not quite unix but it is certainly worth considering.
- 2) When we only partially enter a namespace it is very easy for additional properties to enter that namespace. For example we enter the pid namespace and the mount namespace, but keep our current working directory in the previous namespace. Then a process in the restricted namespace can get out by cd into /proc/<?>/cwd.

If someones permissions to various objects does not depend on the namespace they are in quite possibly this is a non-issue. If we actually depend on the isolation to keep things secure enter is a setup for a first rate escape.

Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 09/17] nsproxy: add namespace flags Posted by ebiederm on Mon, 11 Dec 2006 20:02:15 GMT View Forum Message <> Reply to Message

Cedric Le Goater <clg@fr.ibm.com> writes:

```
> Eric W. Biederman wrote:
>> Cedric Le Goater <clg@fr.ibm.com> writes:
>>>> /*
>>>> + * namespaces flags
>>>> + */
>>>> +#define NS_MNT 0x00000001
>>>> +#define NS UTS 0x00000002
>>>> +#define NS IPC 0x00000004
>>>> +#define NS PID 0x00000008
>>>> +#define NS NET 0x00000010
>>>> +#define NS_USER 0x00000020
>>>> +#define NS_ALL (NS_MNT|NS_UTS|NS_IPC|NS_PID|NS_NET|NS_USER)
>>>> hmm, why _another_ set of flags to refer to the
>>>> namespaces?
>>> well, because namespaces are a new kind in the kernel
>> Gratuitous incompatibility.
>?
Changing the numbers for no good reason. We can easily keep the existing numbers.
>>> is the clone()/unshare() set of flags not sufficient
>>>> for that?
>>> because we are reaching the limits of the CLONE_ flags.
>> Not really. There are at least 8 bits that clone cannot use
>> but that unshare can.
> please, could you list them?
```

CSIGNAL. There are a several others as well that will never mean anything in an unshare context:

I believe the 10 flags below are also nonsense from an unshare perspective. CLONE\_PTRACE 0x00002000 CLONE VFORK 0x00004000 CLONE\_PARENT 0x00008000 CLONE\_SETTLS 0x00080000 CLONE PARENT SETTID 0x00100000 CLONE CHILD CLEARTID 0x00200000 CLONE DETACHED 0x00400000 CLONE UNTRACED 0x00800000 CLONE CHILD SETTID 0x01000000 CLONE STOPPED 0x02000000 >>> if so, shouldn't we switch (or even better change? >>>> the unshare() too) to a new set of syscalls? >>> unshare\_ns() is a new syscall and we don't really need a >>> clone anyway, nop? >> Huh? Clone should be the primary. There are certain namespaces >> that it are very hard to unshare, without creating a new process. > You just said above that clone had less available flags than > unshare ...

I did. It is just easier to support clone than unshare.

- > anyway, could you elaborate a bit more? I have the opposite
- > feeling and you gave me that impression also a few month ago.

> No problem for me, i just want a way to use this stuff without

My feeling is basically that there are some things that we can do much more cleanly at process creation time.

>From an implementation standpoint unshare is fairly nasty because it keeps things from being invariants across the lifetime of a process. Which means it contains more races and is harder to support. That's why we can't unshare we can share with clone right now.

### Eric

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable

## Posted by serue on Mon, 11 Dec 2006 20:03:01 GMT

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```
Quoting Eric W. Biederman (ebiederm@xmission.com):
> "Serge E. Hallyn" <serue@us.ibm.com> writes:
> > Quoting Serge E. Hallyn (serue@us.ibm.com):
> >> Quoting Eric W. Biederman (ebiederm@xmission.com):
>>> Herbert Poetzl <a href="mailto:herbert@13thfloor.at">herbert@13thfloor.at</a> writes:
>>> > Beyond that yes it seems to make sense to let user space
>>> >> maintain any mapping of containers to ids.
>>>>>
>>> > I agree with that, but we need something to move
>>> > around between the various spaces ...
> >> >
>>> If you have CAP_SYS_PTRACE or you have a child process
>>> > in a container you can create another with ptrace.
> >> >
>>> > Now I don't mind optimizing that case, with something like
>>> the proposed bind_ns syscall. But we need to be darn certain
>>> why it is safe, and does not change the security model that
>>> > we currently have.
>>> Sigh, and that's going to have to be a discussion per namespace.
> >
>> Well, assuming that we're using pids as identifiers, that means
>> we can only enter decendent namespaces, which means 'we' must
> > have created them. So anything we could do by entering the ns,
> > we could have done by creating it as well, right?
> It isn't strict descendents who we can see. i.e. init can create
> the thing, and we could have just logged into the network but init
> and us still share the same pid namespace.
>
> But yes it would be we can only enter descendent namespaces, for
> some definition of enter.
> There are two issues.
> 1) We may have a namespace we want to create and then remove the ability
> for the sysadmin to fiddle with, so it can play with encrypted data or
   something like that safely. Not quite unix but it is certainly worth
   considering.
```

Yeah, that occurred to me, but it doesn't seem like we can possibly make sufficient guarantees to the client to make this worthwhile.

I'd love to be wrong about that, but if nothing else we can't prove to the client that they're running on an unhacked host. So the host admin will always have to be trusted.

- > 2) When we only partially enter a namespace it is very easy for additional
- > properties to enter that namespace. For example we enter the pid
- > namespace and the mount namespace, but keep our current working directory
- > in the previous namespace. Then a process in the restricted namespace
- > can get out by cd into /proc/<?>/cwd.

Yup, entering existing namespaces should be all-or-nothing.

- > If someones permissions to various objects does not depend on the namespace
- > they are in quite possibly this is a non-issue. If we actually depend on
- > the isolation to keep things secure enter is a setup for a first rate escape.

I don't believe the isolation can be effective between two namespaces where one is an ancestor of another. It can be so long as one isn't the ancestor of another, but then we're not allowing either to enter the other namespace. So it's not a problem.

The bind\_ns() proposed by Cedric is stricter, only allowing nsid 0 to switch namespaces. So it may be overly restrictive, and does introduce a new global namespace, but it is safe.

-serge

>

\_\_\_\_\_

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by ebiederm on Mon, 11 Dec 2006 20:34:40 GMT View Forum Message <> Reply to Message

"Serge E. Hallyn" <serue@us.ibm.com> writes:

- > Quoting Eric W. Biederman (ebiederm@xmission.com):
- > Yeah, that occurred to me, but it doesn't seem like we can possibly make
- > sufficient guarantees to the client to make this worthwhile.
- > I'd love to be wrong about that, but if nothing else we can't prove to
- > the client that they're running on an unhacked host. So the host admin
- > will always have to be trusted.

To some extent that is true. Although all security models we have currently fall down if you hack the kernel, or run your kernel in a hacked virtual environment. It would be nice if under normal conditions you could mount an encrypted filesystem only in a container

and not have concerns of those files escaping.

Which would probably be a matter of having a separate uid\_ns and not allowing process outside of your container to have any permissions in that filesystem.

- >> 2) When we only partially enter a namespace it is very easy for additional
- >> properties to enter that namespace. For example we enter the pid
- >> namespace and the mount namespace, but keep our current working directory
- >> in the previous namespace. Then a process in the restricted namespace
- >> can get out by cd into /proc/<?>/cwd.

>

> Yup, entering existing namespaces should be all-or-nothing.

A truly all-or-nothing has the problem that there is no external input into the container, and a very controlled external input to the existing container is what this is about.

- >> If someones permissions to various objects does not depend on the namespace
- >> they are in quite possibly this is a non-issue. If we actually depend on
- >> the isolation to keep things secure enter is a setup for a first rate escape.

\_

- > I don't believe the isolation can be effective between two namespaces
- > where one is an ancestor of another. It can be so long as one isn't
- > the ancestor of another, but then we're not allowing either to enter
- > the other namespace. So it's not a problem.

#### Reasonable.

- > The bind\_ns() proposed by Cedric is stricter, only allowing nsid 0 to
- > switch namespaces. So it may be overly restrictive, and does introduce
- > a new global namespace, but it is safe.

I will look a little more. There are a lot patches out there that need review. What disturbs a little is that with ptrace we have an existing mechanism that can do everything we want enter or bind\_ns to be able to do.

I actually have code that will let me fork a process in a new namespace today with out needing bind\_ns. What is more I don't even have to be root to use it.

I would very much prefer to see us optimizing our debugging and control interfaces so they are efficient then see us implement something completely new that is problem domain specific.

Eric	
Containers mailing list	

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by serue on Mon, 11 Dec 2006 22:01:15 GMT

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Quoting Eric W. Biederman (ebiederm@xmission.com):

- > "Serge E. Hallyn" <serue@us.ibm.com> writes:
- > > Quoting Eric W. Biederman (ebiederm@xmission.com):
- > >
- >> Yeah, that occurred to me, but it doesn't seem like we can possibly make
- > > sufficient guarantees to the client to make this worthwhile.
- > >
- >> I'd love to be wrong about that, but if nothing else we can't prove to
- > > the client that they're running on an unhacked host. So the host admin
- > > will always have to be trusted.

>

- > To some extent that is true. Although all security models we have
- > currently fall down if you hack the kernel, or run your kernel
- > in a hacked virtual environment. It would be nice if under normal
- > conditions you could mount an encrypted filesystem only in a container
- > and not have concerns of those files escaping.

Hmm, well perhaps I'm being overly pessimistic - IBM research did have a demo based on TPM of remote attestation, which may be usable for ensuring that you're connecting to a service on your virtual machine on a certain (unhacked) kernel on particular hardware, in which case what you're talking about may be possible - given a stringent initial environment (i.e. not the 'gimme \$20/month for a hosted partition in arizona' environment).

Given that, perhaps having a virtual machine with access to encrypted storage - safe from the host machine admins - may not be unattainable after all. And given that, it would be worth designing the ns\_enter() system call so that a parent cannot enter some child namespace.

- > Which would probably be a matter of having a separate uid\_ns and not
- > allowing process outside of your container to have any permissions in
- > that filesystem.

Yup. Or even just a separate uid\_ns and an ecryptfs partition, so that the host can back up the encrypted data incrementally (per file, i.e. not just the whole dmcrypted loop file).

-serge

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by serue on Mon, 11 Dec 2006 22:18:34 GMT

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Quoting Eric W. Biederman (ebiederm@xmission.com):

- > I actually have code that will let me fork a process in a new namespace today
- > with out needing bind\_ns. What is more I don't even have to be root
- > to use it.

Can you elaborate? The user namespace patches don't enforce ptrace yet, so you could unshare as root, become uid 500, then as uid 500 in the original namespace ptrace the process in the new namespace. Is that what you're doing? If (when) ptrace enforces the uid namespace, will that stop what you're doing?

-serge

\_\_\_\_\_

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Dave Hansen on Mon, 11 Dec 2006 22:53:08 GMT View Forum Message <> Reply to Message

On Mon, 2006-12-11 at 16:23 +0100, Cedric Le Goater wrote:

- > > Even letting the concept of nsproxy escape to user space sounds wrong.
- > > nsproxy is an internal space optimization. It's not struct container
- > > and I don't think we want it to become that.

>

- > i don't agree here. we need that, so does openvz, vserver, people working
- > on resource management.

I think what those projects need is \_some\_ way to group tasks. I'm not sure they actually need nsproxies.

Two tasks in the same container could very well have different nsproxies. The nsproxy defines how the pid namespace, and pid<->task mappings happen for a given task. The init process for a container is special and might actually appear in more than one pid namespace, while

its children might only appear in one. That means that this init process's nsproxy can and should actually be different from its children's. This is despite the fact that they are in the same container.

If we really need this 'container' grouping, it can easily be something pointed to \_by\_ the nsproxy, but it shouldn't \_be\_ the nsproxy.

-- Dave

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by ebiederm on Tue, 12 Dec 2006 03:28:40 GMT View Forum Message <> Reply to Message

"Serge E. Hallyn" <serue@us.ibm.com> writes:

- > Quoting Eric W. Biederman (ebiederm@xmission.com):
- >> I actually have code that will let me fork a process in a new namespace today
- >> with out needing bind\_ns. What is more I don't even have to be root
- >> to use it.

>

- > Can you elaborate? The user namespace patches don't enforce ptrace
- > yet, so you could unshare as root, become uid 500, then as uid 500
- > in the original namespace ptrace the process in the new namespace.
- > Is that what you're doing? If (when) ptrace enforces the uid namespace,
- > will that stop what you're doing?

sys\_ptrace is allowed in 2 situations.

- The user and group identities are the same.
- The calling process has CAP\_SYS\_PTRACE capability.

So currently if the uid namespace enforces the user and group checks that will prevent the first case, and is very desirable. But it won't stop someone with CAP\_SYS\_PTRACE. Which given the normal case seems reasonable.

Getting to the point where you can't trace what a process is doing would probably require some additional interprocess firewalling from something like selinux.

Eric			

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Cedric Le Goater on Tue, 12 Dec 2006 07:09:45 GMT View Forum Message <> Reply to Message

## Dave Hansen wrote:

- > On Mon, 2006-12-11 at 16:23 +0100, Cedric Le Goater wrote:
- >>> Even letting the concept of nsproxy escape to user space sounds wrong.
- >>> nsproxy is an internal space optimization. It's not struct container
- >>> and I don't think we want it to become that.
- >> i don't agree here. we need that, so does openvz, vserver, people working
- >> on resource management.

>

- > I think what those projects need is \_some\_ way to group tasks. I'm not
- > sure they actually need nsproxies.

not only tasks. ipc, fs, etc.

- > Two tasks in the same container could very well have different
- > nsproxies. The nsproxy defines how the pid namespace, and pid<->task
- > mappings happen for a given task.

not only. there are other namespaces in nsproxy.

- > The init process for a container is
- > special and might actually appear in more than one pid namespace, while
- > its children might only appear in one. That means that this init
- > process's nsproxy can and should actually be different from its
- > children's. This is despite the fact that they are in the same
- > container.

>

- > If we really need this 'container' grouping, it can easily be something
- > pointed to \_by\_ the nsproxy, but it shouldn't \_be\_ the nsproxy.

ok so let's add a container object, containing a nsproxy and add another indirection ...

C.

Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Cedric Le Goater on Tue, 12 Dec 2006 07:11:33 GMT

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## Serge E. Hallyn wrote:

- > Quoting Serge E. Hallyn (serue@us.ibm.com):
- >> Quoting Eric W. Biederman (ebiederm@xmission.com):
- >>> Herbert Poetzl <herbert@13thfloor.at> writes:
- >>>> Beyond that yes it seems to make sense to let user space
- >>>> maintain any mapping of containers to ids.
- >>>> I agree with that, but we need something to move
- >>> around between the various spaces ...
- >>> If you have CAP\_SYS\_PTRACE or you have a child process
- >>> in a container you can create another with ptrace.

>>>

- >>> Now I don't mind optimizing that case, with something like
- >>> the proposed bind\_ns syscall. But we need to be darn certain
- >>> why it is safe, and does not change the security model that
- >>> we currently have.
- >> Sigh, and that's going to have to be a discussion per namespace.

>

> Well, assuming that we're using pids as identifiers, that means

we can't because a process could die while the namespace is still referenced by an other subsystem. We need some kind of id.

- > we can only enter decendent namespaces, which means 'we' must
- > have created them. So anything we could do by entering the ns,
- > we could have done by creating it as well, right?

\_\_\_\_\_

Containers mailing list Containers@lists.osdl.org

https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by ebiederm on Tue, 12 Dec 2006 07:52:14 GMT View Forum Message <> Reply to Message

Cedric Le Goater <clg@fr.ibm.com> writes:

- > Serge E. Hallyn wrote:
- >> Well, assuming that we're using pids as identifiers, that means
- > we can't because a process could die while the namespace is still
- > referenced by an other subsystem. We need some kind of id.

Think of a session think of a process group heck think of threads a pid is not tied to one task struct. It is absolutely not a problem for a namespace to do get\_pid(...) when it is initialized and put\_pid(...) just before it is freed.

All of the mechanisms for using pids for something like this are already in place.

What we don't have is a fast pid to namespace transfer. But that is just an extra pointer in struct pid. Really that is a trivial patch. Giving every namespace a pid pointer in struct pid takes a little more space then I would like but it is not a big deal.

Eric

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by ebiederm on Tue, 12 Dec 2006 08:37:48 GMT

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Cedric Le Goater <clg@fr.ibm.com> writes:

- > Dave Hansen wrote:
- >> On Mon, 2006-12-11 at 16:23 +0100, Cedric Le Goater wrote:
- >>>> Even letting the concept of nsproxy escape to user space sounds wrong.
- >>> nsproxy is an internal space optimization. It's not struct container
- >>> and I don't think we want it to become that.
- >>> i don't agree here. we need that, so does openvz, vserver, people working
- >>> on resource management.

>>

>> I think what those projects need is \_some\_ way to group tasks. I'm not

>> sure they actually need nsproxies.

\_

> not only tasks. ipc, fs, etc.

What is the important aspect that you need to group. What concept are you trying to convey?

How do you describe a container in which someone is using the pam\_namespace module? So different tasks in the container have a different mount namespace?

- >> Two tasks in the same container could very well have different
- >> nsproxies. The nsproxy defines how the pid namespace, and pid<->task

>> mappings happen for a given task. > not only. there are other namespaces in nsproxy. The point is that there is not a one to one mapping between containers and nsproxies. There are likely to be more nsproxies than containers. >> The init process for a container is >> special and might actually appear in more than one pid namespace, while >> its children might only appear in one. That means that this init >> process's nsproxy can and should actually be different from its >> children's. This is despite the fact that they are in the same >> container. >> >> If we really need this 'container' grouping, it can easily be something >> pointed to \_by\_ the nsproxy, but it shouldn't \_be\_ the nsproxy. > ok so let's add a container object, containing a nsproxy and add > another indirection ... Well that isn't what Dave suggested, and I don't think it will give you what you want.

Eric

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by dev on Tue, 12 Dec 2006 08:43:38 GMT View Forum Message <> Reply to Message

>>>Even letting the concept of nsproxy escape to user space sounds wrong.
>>>nsproxy is an internal space optimization. It's not struct container
>>>and I don't think we want it to become that.
>>
>>i don't agree here. we need that, so does openvz, vserver, people working
>>on resource management.
>
> I think what those projects need is \_some\_ way to group tasks. I'm not
> sure they actually need nsproxies.
>
> Two tasks in the same container could very well have different
> nsproxies.
what is container then from your POV?

- > The nsproxy defines how the pid namespace, and pid<->task
- > mappings happen for a given task. The init process for a container is
- > special and might actually appear in more than one pid namespace, while
- > its children might only appear in one. That means that this init
- > process's nsproxy can and should actually be different from its
- > children's. This is despite the fact that they are in the same
- > container.

nsproxy has references to all namespaces, not just pid namespace.

Thus it is a container "view" effectively.

If container is something different, then please define it.

- > If we really need this 'container' grouping, it can easily be something
- > pointed to \_by\_ the nsproxy, but it shouldn't \_be\_ the nsproxy.

You can add another indirection if really want it so much...

But is it required?

We created nsproxy which adds another level of indirection, but from performance POV it is questinable. I can say that we had a nice experience, when adding a single dereference in TCP code resulted in ~0.5% performance degradation.

Thanks, Kirill

\_\_\_\_\_

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by ebiederm on Tue, 12 Dec 2006 08:57:48 GMT

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Kirill Korotaev <dev@sw.ru> writes:

>>

- >> I think what those projects need is \_some\_ way to group tasks. I'm not
- >> sure they actually need nsproxies.

>>

- >> Two tasks in the same container could very well have different
- >> nsproxies.
- > what is container then from your POV?

A nested instance of user space. User space may unshare things such as the mount namespace so it can give users the ability to control their own mounts and the like.

>> The nsproxy defines how the pid namespace, and pid<->task

- >> mappings happen for a given task. The init process for a container is
- >> special and might actually appear in more than one pid namespace, while
- >> its children might only appear in one. That means that this init
- >> process's nsproxy can and should actually be different from its
- >> children's. This is despite the fact that they are in the same
- >> container.
- > nsproxy has references to all namespaces, not just pid namespace.
- > Thus it is a container "view" effectively.
- > If container is something different, then please define it.

nsproxy has exactly one instance of all namespaces. A container in the general case can hold other containers, and near containers (like processes with separate mount namespaces). As well as processes.

So nsproxy currently captures the common case for containers but not the general case.

- >> If we really need this 'container' grouping, it can easily be something
- >> pointed to \_by\_ the nsproxy, but it shouldn't \_be\_ the nsproxy.
- > You can add another indirection if really want it so much...
- > But is it required?
- > We created nsproxy which adds another level of indirection, but from performance
- > POV
- > it is questinable. I can say that we had a nice experience, when adding
- > a single dereference in TCP code resulted in ~0.5% performance degradation.

I totally agree with that, nsproxy is something we need to watch from a performance point of view. nsproxy is primarily a space optimization to keep from bloating task struct, and possibly a fork time optimization. At least at the point we added it no one could measure overhead from using it.

That is one of the reasons I don't want nsproxy to become explicit and be exported to user space. So if it is a performance problem we can change the implementation without affecting users.

Eric

Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by serue on Tue, 12 Dec 2006 15:29:12 GMT

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Quoting Eric W. Biederman (ebiederm@xmission.com): > "Serge E. Hallyn" <serue@us.ibm.com> writes: > > Quoting Eric W. Biederman (ebiederm@xmission.com): >>> I actually have code that will let me fork a process in a new namespace today >>> with out needing bind\_ns. What is more I don't even have to be root > >> to use it. > > >> Can you elaborate? The user namespace patches don't enforce ptrace >> yet, so you could unshare as root, become uid 500, then as uid 500 >> in the original namespace ptrace the process in the new namespace. >> Is that what you're doing? If (when) ptrace enforces the uid namespace, > > will that stop what you're doing? > > sys\_ptrace is allowed in 2 situations. > - The user and group identities are the same. > - The calling process has CAP SYS PTRACE capability. > So currently if the uid namespace enforces the user and group checks > that will prevent the first case, and is very desirable. But it won't > stop someone with CAP\_SYS\_PTRACE. Which given the normal case seems > reasonable.

Yes, I was forgetting that intra-container ptrace is generally inhibited by lack of a handle to processes in the other container. So:

- in checkpoint/restart usage, the normal CAP\_SYS\_PTRACE semantics is fine
- . inside a vserver, the normal CAP\_SYS\_PTRACE is fine
- . in general, a process inside one vserver cannot reference a process in another vserver, so we don't need to worry about ptrace permissions at all
- . however, if we want to (as per emails yesterday) provide some bit of enforcement of limits from parent namespaces to child namespaces - where a pid is in fact available for at least the init process (and, depending on our final implementation, perhaps all processes) - then we need something more.

As you say, selinux permissions would be one way to obtain this.

- > Getting to the point where you can't trace what a process is doing
- > would probably require some additional interprocess firewalling
- > from something like selinux.

Yup.

thanks, -serge

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by serue on Tue, 12 Dec 2006 15:45:07 GMT View Forum Message <> Reply to Message

Quoting Cedric Le Goater (clg@fr.ibm.com):

- > Dave Hansen wrote:
- > > On Mon, 2006-12-11 at 16:23 +0100, Cedric Le Goater wrote:
- >>>> Even letting the concept of nsproxy escape to user space sounds wrong.
- >>> nsproxy is an internal space optimization. It's not struct container
- >>>> and I don't think we want it to become that.
- >>> i don't agree here. we need that, so does openvz, vserver, people working
- >>> on resource management.

> >

- >> I think what those projects need is some way to group tasks. I'm not
- > > sure they actually need nsproxies.

>

> not only tasks. ipc, fs, etc.

>

- > > Two tasks in the same container could very well have different
- > > nsproxies. The nsproxy defines how the pid namespace, and pid<->task
- > > mappings happen for a given task.

>

> not only. there are other namespaces in nsproxy.

Right, and as Eric has pointed out, you may well want to use one id to refer to several nsproxies - for instance if you are using unshare to provide per-user private mount namespaces using pam\_namespace.so (that's mostly for LSPP systems right now, but I do this on my laptop too). All my accounts are in the same 'container', but have different mount namespaces, hence different nsproxies.

- > > The init process for a container is
- > > special and might actually appear in more than one pid namespace, while
- > > its children might only appear in one. That means that this init
- > > process's nsproxy can and should actually be different from its
- >> children's. This is despite the fact that they are in the same
- > > container.

> >

- >> If we really need this 'container' grouping, it can easily be something
- > > pointed to \_by\_ the nsproxy, but it shouldn't \_be\_ the nsproxy.

> ok so let's add a container object, containing a nsproxy and add > another indirection
No thanks.
Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers
Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Cedric Le Goater on Tue, 12 Dec 2006 18:29:08 GMT View Forum Message <> Reply to Message
>>> If someones permissions to various objects does not depend on the namespace >>> they are in quite possibly this is a non-issue. If we actually depend on >>> the isolation to keep things secure enter is a setup for a first rate escape. >> I don't believe the isolation can be effective between two namespaces >> where one is an ancestor of another. It can be so long as one isn't >> the ancestor of another, but then we're not allowing either to enter >> the other namespace. So it's not a problem.
> Reasonable.
>> The bind_ns() proposed by Cedric is stricter, only allowing nsid 0 to >> switch namespaces. So it may be overly restrictive, and does introduce >> a new global namespace, but it is safe.
<ul> <li>I will look a little more. There are a lot patches out there that need</li> <li>review. What disturbs a little is that with ptrace we have an existing</li> <li>mechanism that can do everything we want enter or bind_ns to be able to do.</li> </ul>
Eric, you have this habit of flooding us with email whenever a patchset is sent on this topic. It is a bad habit. Please take some time to look at it before. There is work behind it and it tries to address some issues.
This patchset has been sent on container@ as a proposal for -mm. I'll try to make a summary of how we can improve next one to move forward.
I still need to read all your emails :)
thanks,
C.
Containers mailing list Containers@lists.osdl.org

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Herbert Poetzl on Tue, 12 Dec 2006 23:22:22 GMT View Forum Message <> Reply to Message

```
On Mon, Dec 11, 2006 at 04:01:15PM -0600, Serge E. Hallyn wrote:
> Quoting Eric W. Biederman (ebiederm@xmission.com):
> > "Serge E. Hallyn" <serue@us.ibm.com> writes:
>> Quoting Eric W. Biederman (ebiederm@xmission.com):
>> Yeah, that occurred to me, but it doesn't seem like we can possibly make
>> sufficient guarantees to the client to make this worthwhile.
>> I'd love to be wrong about that, but if nothing else we can't prove to
>>> the client that they're running on an unhacked host. So the host admin
>> will always have to be trusted.
>> To some extent that is true. Although all security models we have
> > currently fall down if you hack the kernel, or run your kernel
>> in a hacked virtual environment. It would be nice if under normal
>> conditions you could mount an encrypted filesystem only in a container
> > and not have concerns of those files escaping.
> Hmm, well perhaps I'm being overly pessimistic - IBM research did have a
> demo based on TPM of remote attestation, which may be usable for
> ensuring that you're connecting to a service on your virtual machine on
> a certain (unhacked) kernel on particular hardware, in which case what
> you're talking about may be possible - given a stringent initial
> environment (i.e. not the 'gimme $20/month for a hosted partition in
> arizona' environment).
```

interesting, how would you \_ensure\_ from inside such an environment, that nobody tampered with the kernel you are running on?

- > Given that, perhaps having a virtual machine with access to encrypted
- > storage safe from the host machine admins may not be unattainable
- > after all. And given that, it would be worth designing the ns enter()
- > system call so that a parent cannot enter some child namespace.

we currently call this Context Privacy, and it is partially implemented, but of course, it does only work if the kernel is known good

> > Which would probably be a matter of having a separate uid\_ns and not

- > > allowing process outside of your container to have any permissions in
- > > that filesystem.

>

- > Yup. Or even just a separate uid\_ns and an ecryptfs partition, so
- > that the host can back up the encrypted data incrementally (per file,
- > i.e. not just the whole dmcrypted loop file).

it's simple to avoid access to certain 'tagged' devices and/or filesystems, it's hard to handle kernel modifications or even simple things like reading the kernel memory ...

best, Herbert

> -serge

Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Herbert Poetzl on Wed, 13 Dec 2006 04:55:58 GMT View Forum Message <> Reply to Message

On Tue, Dec 12, 2006 at 11:43:38AM +0300, Kirill Korotaev wrote:

- >>>>Even letting the concept of nsproxy escape to user space sounds wrong.
- >>>nsproxy is an internal space optimization. It's not struct container
- >>>and I don't think we want it to become that.
- > >>i don't agree here. we need that, so does openvz, vserver, people
- >>>working on resource management.
- > >
- > >
- >> I think what those projects need is \_some\_ way to group tasks. I'm
- > > not sure they actually need nsproxies.
- > >
- >> Two tasks in the same container could very well have different
- > > nsproxies.

and typically, they will ...

> what is container then from your POV?

from my PoV, a container is something keeping processes \_inside\_ which basically requires the following elements:

- isolation from other containers
- virtualization of unique elements
- limitation on resources
- policy on all interfaces

the current spaces mostly address the isolation and to some degree, the virtualization, which is a good thing, but the container also requires the resource limitation and the policy, to handle interfaces to the outside (should not be new to you, actually:)

so the container (may it be represented by a structure or not), may reference an nsproxy (as we do in the 2.6.19 versions of Linux-VServer) but an nsproxy is not the proper element to define a container ..

we also want to be able to have sub spaces inside a container, as long as they do not interfere or overcome the limitations and policy

- > > The nsproxy defines how the pid namespace, and pid<->task
- > > mappings happen for a given task. The init process for a container is
- > > special and might actually appear in more than one pid namespace, while
- >> its children might only appear in one. That means that this init
- > > process's nsproxy can and should actually be different from its
- >> children's. This is despite the fact that they are in the same
- > > container.
- > nsproxy has references to all namespaces, not just pid namespace.
- > Thus it is a container "view" effectively.

it is a view into the world of one or more processes, but not necessarily the view of all processes inside a container:)

> If container is something different, then please define it.

see above ...

- >> If we really need this 'container' grouping, it can easily be something
- > > pointed to \_by\_ the nsproxy, but it shouldn't \_be\_ the nsproxy.
- > You can add another indirection if really want it so much...
- > But is it required?
- > We created nsproxy which adds another level of indirection, but from

> performance POV it is questinable.

I'm not very happy with the nsproxy abstraction, as I think it would be better handled per task, and I still have no real world test results what overhead the nsproxy indirection causes

- > I can say that we had a nice experience, when adding a single
- > dereference in TCP code resulted in ~0.5% performance degradation.

yes, that is what I fear is happening right now with the nsproxy ... but I think we need to test that, and if it makes sense, switch to task direct spaces (as we had before), just more of them ...

best, Herbert

- > Thanks,
- > Kirill

>

>\_\_\_\_\_

- > Containers mailing list
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- > https://lists.osdl.org/mailman/listinfo/containers

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Cedric Le Goater on Wed, 13 Dec 2006 14:44:06 GMT View Forum Message <> Reply to Message

Eric W. Biederman wrote:

> Cedric Le Goater <clg@fr.ibm.com> writes:

>

>> Serge E. Hallyn wrote:

- >>> Well, assuming that we're using pids as identifiers, that means
- >> we can't because a process could die while the namespace is still
- >> referenced by an other subsystem. We need some kind of id.

>

- > Think of a session think of a process group heck think of threads
- > a pid is not tied to one task struct. It is absolutely not a problem
- > for a namespace to do get\_pid(...) when it is initialized and put\_pid(...)
- > just before it is freed.

>

> All of the mechanisms for using pids for something like this are already > in place.

- > What we don't have is a fast pid to namespace transfer. But that is just
- > an extra pointer in struct pid. Really that is a trivial patch.
- > Giving every namespace a pid pointer in struct pid takes a little more
- > space then I would like but it is not a big deal.

I'm not sure I understand how you want to do this.

Let me try: you would add a 'struct pid pid' field to all namespaces and assign that 'pid' field with the struct pid of the task creating the namespace?

C.

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Cedric Le Goater on Wed, 13 Dec 2006 15:00:35 GMT View Forum Message <> Reply to Message

Serge E. Hallyn wrote:

- > Quoting Cedric Le Goater (clg@fr.ibm.com):
- >> Dave Hansen wrote:
- >>> On Mon, 2006-12-11 at 16:23 +0100, Cedric Le Goater wrote:
- >>>> Even letting the concept of nsproxy escape to user space sounds wrong.
- >>>> nsproxy is an internal space optimization. It's not struct container
- >>>> and I don't think we want it to become that.
- >>> i don't agree here. we need that, so does openvz, vserver, people working
- >>> on resource management.
- >>> I think what those projects need is some way to group tasks. I'm not
- >>> sure they actually need nsproxies.
- >> not only tasks. ipc, fs, etc.

- >>> Two tasks in the same container could very well have different
- >>> nsproxies. The nsproxy defines how the pid namespace, and pid<->task
- >>> mappings happen for a given task.
- >> not only. there are other namespaces in nsproxy.

- > Right, and as Eric has pointed out, you may well want to use one id to
- > refer to several nsproxies for instance if you are using unshare
- > to provide per-user private mount namespaces using pam\_namespace.so
- > (that's mostly for LSPP systems right now, but I do this on my laptop
- > too). All my accounts are in the same 'container', but have different

> mount namespaces, hence different nsproxies.

I think we have definition issue here: what is a 'container'?

I don't see any issue with the above scenario. unsharing mount namespace results in the creation of a new nsproxy which will require a new identifier in order to find this new mount namespace.

so yes, different mount namespaces, hence different nsproxies, hence different ids if you want to find that new mount namespace.

- >>> The init process for a container is
- >>> special and might actually appear in more than one pid namespace, while
- >>> its children might only appear in one. That means that this init
- >>> process's nsproxy can and should actually be different from its
- >>> children's. This is despite the fact that they are in the same
- >>> container.

>>>

- >>> If we really need this 'container' grouping, it can easily be something
- >>> pointed to \_by\_ the nsproxy, but it shouldn't \_be\_ the nsproxy.
- >> ok so let's add a container object, containing a nsproxy and add
- >> another indirection ...

>

> No thanks.

exactly.

C.

Containers mailing list

Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Cedric Le Goater on Wed, 13 Dec 2006 15:02:56 GMT View Forum Message <> Reply to Message

Eric W. Biederman wrote:

> Cedric Le Goater <clg@fr.ibm.com> writes:

>

>> Dave Hansen wrote:

>>> On Mon, 2006-12-11 at 16:23 +0100, Cedric Le Goater wrote:

>>>> Even letting the concept of nsproxy escape to user space sounds wrong.

>>>> nsproxy is an internal space optimization. It's not struct container

>>>> and I don't think we want it to become that.

>>>> i don't agree here. we need that, so does openvz, vserver, people working >>>> on resource management.
>>> I think what those projects need is \_some\_ way to group tasks. I'm not >>> sure they actually need nsproxies.

>> not only tasks. ipc, fs, etc.

>

- > What is the important aspect that you need to group. What concept
- > are you trying to convey?

>

- > How do you describe a container in which someone is using the
- > pam\_namespace module? So different tasks in the container have
- > a different mount namespace?

let's define a container first. I'm not sure for what you are using that term.

>>> Two tasks in the same container could very well have different

- >>> nsproxies. The nsproxy defines how the pid namespace, and pid<->task
- >>> mappings happen for a given task.
- >> not only. there are other namespaces in nsproxy.

>

- > The point is that there is not a one to one mapping between containers
- > and nsproxies. There are likely to be more nsproxies than containers.

again. please explain the difference that you see between a container and a nsproxy. I don't get it and i might be missing something important doing this short cut.

C.

Containers mailing list Containers@lists.osdl.org

https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Cedric Le Goater on Wed, 13 Dec 2006 15:17:39 GMT View Forum Message <> Reply to Message

## Herbert Poetzl wrote:

> On Tue, Dec 12, 2006 at 11:43:38AM +0300, Kirill Korotaev wrote:

>>>> Even letting the concept of nsproxy escape to user space sounds wrong.

>>>> nsproxy is an internal space optimization. It's not struct container

>>>> and I don't think we want it to become that.

>

>>> i don't agree here. we need that, so does openvz, vserver, people

>>> working on resource management.

>>>

```
>>> I think what those projects need is _some_ way to group tasks. I'm
>>> not sure they actually need nsproxies.
>>>
>>> Two tasks in the same container could very well have different
>>> nsproxies.
> and typically, they will ...
that means we are missing a container object then, a vps, a vcontext, a
vsomething, nop?
>> what is container then from your POV?
>
> from my PoV, a container is something keeping
> processes _inside_ which basically requires
> the following elements:
> - isolation from other containers
> - virtualization of unique elements
> - limitation on resources
> - policy on all interfaces
> the current spaces mostly address the isolation
> and to some degree, the virtualization, which
> is a good thing, but the container also requires
> the resource limitation and the policy, to handle
> interfaces to the outside (should not be new to
> you, actually:)
>
> so the container (may it be represented by a
> structure or not), may reference an nsproxy
> (as we do in the 2.6.19 versions of Linux-VServer)
> but an nsproxy is not the proper element to
> define a container ..
agree. it's not complete.
should we address that by introducing a new object?
could that be done on per-product basis? I mean like
in a driver model.
> we also want to be able to have sub spaces inside
> a container, as long as they do not interfere or
> overcome the limitations and policy
>>> The nsproxy defines how the pid namespace, and pid<->task
>>> mappings happen for a given task. The init process for a container is
>>> special and might actually appear in more than one pid namespace, while
```

```
>>> its children might only appear in one. That means that this init
>>> process's nsproxy can and should actually be different from its
>>> children's. This is despite the fact that they are in the same
>>> container.
>> nsproxy has references to all namespaces, not just pid namespace.
>> Thus it is a container "view" effectively.
> it is a view into the world of one or more processes,
> but not necessarily the view of all processes inside
> a container :)
>> If container is something different, then please define it.
> see above ...
>>> If we really need this 'container' grouping, it can easily be something
>>> pointed to _by_ the nsproxy, but it shouldn't _be_ the nsproxy.
>> You can add another indirection if really want it so much...
>> But is it required?
>> We created nsproxy which adds another level of indirection, but from
>> performance POV it is questinable.
> I'm not very happy with the nsproxy abstraction,
> as I think it would be better handled per task,
> and I still have no real world test results what
> overhead the nsproxy indirection causes
>
>> I can say that we had a nice experience, when adding a single
>> dereference in TCP code resulted in ~0.5% performance degradation.
>
> yes, that is what I fear is happening right now
> with the nsproxy ... but I think we need to test
> that, and if it makes sense, switch to task direct
> spaces (as we had before), just more of them ...
getting some figures would be nice and we might also be able
to improve the current naproxy model.
C.
Containers mailing list
Containers@lists.osdl.org
https://lists.osdl.org/mailman/listinfo/containers
```

## Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by ebiederm on Wed, 13 Dec 2006 18:53:56 GMT

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Cedric Le Goater <clg@fr.ibm.com> writes:

> Eric W. Biederman wrote:

>>

- >> What we don't have is a fast pid to namespace transfer. But that is just
- >> an extra pointer in struct pid. Really that is a trivial patch.
- >> Giving every namespace a pid pointer in struct pid takes a little more
- >> space then I would like but it is not a big deal.

>

> I'm not sure I understand how you want to do this.

>

- > Let me try: you would add a 'struct pid pid' field to all namespaces and
- > assign that 'pid' field with the struct pid of the task creating the
- > namespace ?

Yes a struct pid \*pid field, that we did the proper reference counting on.

As for which pid to assign, that is a little trickier. The struct pid of the task creating the namespace is the obvious choice and that will always work for clone. For unshare that would only work if we added the restriction you can't unshare if someone already has used that pid for that kind of namespace.

Eric

Containers mailing list

Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by Cedric Le Goater on Thu, 14 Dec 2006 13:17:47 GMT View Forum Message <> Reply to Message

- >> Let me try: you would add a 'struct pid pid' field to all namespaces and
- >> assign that 'pid' field with the struct pid of the task creating the
- >> namespace?

>

- > Yes a struct pid \*pid field, that we did the proper reference counting
- > on.

sure.

- > As for which pid to assign, that is a little trickier. The struct pid
- > of the task creating the namespace is the obvious choice and that will
- > always work for clone. For unshare that would only work if we added
- > the restriction you can't unshare if someone already has used that pid
- > for that kind of namespace.

clone will also be an issue if more than one namespace is unshared. Do you use the same 'struct pid\*' for each namespace? hmm, it feels also wrong.

having an id field in the namespace and using a bind\_ns like syscall to let the user assign whatever id he wants to, doesn't seem to be such a bad idea.

C.

Containers mailing list

Containers@lists.osdl.org

https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by ebiederm on Thu, 14 Dec 2006 21:08:58 GMT View Forum Message <> Reply to Message

Cedric Le Goater <clg@fr.ibm.com> writes:

- >>> Let me try: you would add a 'struct pid pid' field to all namespaces and
- >>> assign that 'pid' field with the struct pid of the task creating the
- >>> namespace?

>>

- >> Yes a struct pid \*pid field, that we did the proper reference counting
- >> on.

>

> sure.

>

- >> As for which pid to assign, that is a little trickier. The struct pid
- >> of the task creating the namespace is the obvious choice and that will
- >> always work for clone. For unshare that would only work if we added
- >> the restriction you can't unshare if someone already has used that pid
- >> for that kind of namespace.

>

- > clone will also be an issue if more than one namespace is unshared. Do
- > you use the same 'struct pid\*' for each namespace ? hmm, it feels also
- > wrong.

>

- > having an id field in the namespace and using a bind\_ns like syscall
- > to let the user assign whatever id he wants to, doesn't seem to be

> such a bad idea.

I think I would probably have suggested simply taking the next available id in that case in practice. Partly it depends on exactly what we are trying to do with these.

But I do agree getting that last little details right so some corner case doesn't feel wrong is hard. That is why I try and put off this kind of things until as much is known about how we are going to use it as possible. So we can make a good decision and solve practical problems, and not theoretical ones.

## Eric

Containers mailing list Containers@lists.osdl.org https://lists.osdl.org/mailman/listinfo/containers

Subject: Re: [patch -mm 08/17] nsproxy: add hashtable Posted by serue on Wed, 20 Dec 2006 06:12:03 GMT

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Quoting Herbert Poetzl (herbert@13thfloor.at):

- > On Mon, Dec 11, 2006 at 04:01:15PM -0600, Serge E. Hallyn wrote:
- > > Quoting Eric W. Biederman (ebiederm@xmission.com):
- >> "Serge E. Hallyn" <serue@us.ibm.com> writes:
- > > >
- >>> Quoting Eric W. Biederman (ebiederm@xmission.com):
- >>>>
- >>> Yeah, that occurred to me, but it doesn't seem like we can possibly make
- >>> sufficient guarantees to the client to make this worthwhile.
- >>>>
- >>> I'd love to be wrong about that, but if nothing else we can't prove to
- >>> the client that they're running on an unhacked host. So the host admin
- >>> will always have to be trusted.
- >>>
- >> To some extent that is true. Although all security models we have
- >>> currently fall down if you hack the kernel, or run your kernel
- >>> in a hacked virtual environment. It would be nice if under normal
- >> conditions you could mount an encrypted filesystem only in a container
- > > and not have concerns of those files escaping.
- > >
- >> Hmm, well perhaps I'm being overly pessimistic IBM research did have a
- >> demo based on TPM of remote attestation, which may be usable for
- > > ensuring that you're connecting to a service on your virtual machine on
- > > a certain (unhacked) kernel on particular hardware, in which case what
- > > you're talking about may be possible given a stringent initial

- > > environment (i.e. not the 'gimme \$20/month for a hosted partition in
- > > arizona' environment).

\_

- > interesting, how would you \_ensure\_ from inside
- > such an environment, that nobody tampered with
- > the kernel you are running on?

Sorry, took awhile to find the best reference, but I guess this would be it:

http://domino.research.ibm.com/library/cyberdig.nsf/1e4115aea78b6e7c85256b360066f0d4/459e9a2b1f668aee85256f330067589f

Another description is

http://domino.research.ibm.com/comm/research\_people.nsf/pages/sailer.ima.html

I guess it was in 2004 that they did a demo of remote attestion at the RSA conference, as described in the third to last paragraph in http://cio.co.nz/cio.nsf/0/03943645293DB008CC256E47005D8EA2?OpenDocument and in

http://domino.research.ibm.com/comm/pr.nsf/pages/news.20040218\_linux.html

-serge

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