Subject: [PATCH v7 0/2] fixes for sock memcg static branch disablement Posted by Glauber Costa on Fri, 25 May 2012 09:32:06 GMT

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Hi Andrew,

I believe this one addresses all of your previous comments.

Besides merging your patch, I tried to improve the comments so they would be more informative.

The first patch, I believe, is already merged at your tree. But I am including it here for completeness. I had no changes since last submission, so feel free to pick the second - or if there are still missing changes you'd like to see, point me to them.

Thanks

Glauber Costa (2): Always free struct memcg through schedule_work() decrement static keys on real destroy time
include/net/sock.h 22 +++++++++++++++++++++++++++++++++

1.7.7.6

Subject: [PATCH v7 1/2] Always free struct memcg through schedule_work() Posted by Glauber Costa on Fri, 25 May 2012 09:32:07 GMT View Forum Message <> Reply to Message

Right now we free struct memcg with kfree right after a rcu grace period, but defer it if we need to use vfree() to get rid of that memory area. We do that by need, because we need vfree to be called in a process context.

This patch unifies this behavior, by ensuring that even kfree will happen in a separate thread. The goal is to have a stable place to call the upcoming jump label destruction function outside the realm of the complicated and quite far-reaching cgroup lock (that can't be held when calling neither the cpu_hotplug.lock nor the jump_label_mutex)

Signed-off-by: Glauber Costa <glommer@parallels.com> Acked-by: Kamezawa Hiroyuki <kamezawa.hiroyu@jp.fujitsu.com>

```
CC: Tejun Heo <tj@kernel.org>
CC: Li Zefan < lizefan@huawei.com>
CC: Johannes Weiner <hannes@cmpxchg.org>
CC: Michal Hocko <mhocko@suse.cz>
CC: Andrew Morton <akpm@linux-foundation.org>
mm/memcontrol.c | 24 ++++++++++
1 files changed, 13 insertions(+), 11 deletions(-)
diff --git a/mm/memcontrol.c b/mm/memcontrol.c
index 932a734..0b4b4c8 100644
--- a/mm/memcontrol.c
+++ b/mm/memcontrol.c
@ @ -245,8 +245,8 @ @ struct mem_cgroup {
 struct rcu_head rcu_freeing;
- * But when using vfree(), that cannot be done at
 * interrupt time, so we must then gueue the work.
+ * We also need some space for a worker in deferred freeing.
+ * By the time we call it, rcu freeing is not longer in use.
 struct work_struct work_freeing;
 };
@ @ -4826,23 +4826,28 @ @ out_free:
}
- * Helpers for freeing a vzalloc()ed mem cgroup by RCU,
+ * Helpers for freeing a kmalloc()ed/vzalloc()ed mem_cgroup by RCU,
 * but in process context. The work freeing structure is overlaid
 * on the rcu freeing structure, which itself is overlaid on memsw.
-static void vfree_work(struct work_struct *work)
+static void free_work(struct work_struct *work)
 struct mem_cgroup *memcg;
+ int size = sizeof(struct mem cgroup);
 memcg = container of(work, struct mem cgroup, work freeing);
vfree(memcg);
+ if (size < PAGE_SIZE)
+ kfree(memcg);
+ else
+ vfree(memcg);
-static void vfree rcu(struct rcu head *rcu head)
```

```
+static void free rcu(struct rcu head *rcu head)
 struct mem_cgroup *memcg;
 memcg = container_of(rcu_head, struct mem_cgroup, rcu_freeing);
- INIT_WORK(&memcg->work_freeing, vfree_work);
+ INIT WORK(&memcg->work freeing, free work);
 schedule_work(&memcg->work_freeing);
}
@ @ -4868,10 +4873,7 @ @ static void __mem_cgroup_free(struct mem_cgroup *memcg)
 free mem cgroup per zone info(memcg, node);
 free_percpu(memcg->stat);
- if (sizeof(struct mem_cgroup) < PAGE_SIZE)</li>
kfree_rcu(memcg, rcu_freeing);
- else
call_rcu(&memcg->rcu_freeing, vfree_rcu);
+ call rcu(&memcg->rcu freeing, free rcu);
}
static void mem cgroup get(struct mem cgroup *memcg)
1.7.7.6
```

Subject: [PATCH v7 2/2] decrement static keys on real destroy time Posted by Glauber Costa on Fri, 25 May 2012 09:32:08 GMT View Forum Message <> Reply to Message

We call the destroy function when a cgroup starts to be removed, such as by a rmdir event.

However, because of our reference counters, some objects are still inflight. Right now, we are decrementing the static_keys at destroy() time, meaning that if we get rid of the last static_key reference, some objects will still have charges, but the code to properly uncharge them won't be run.

This becomes a problem specially if it is ever enabled again, because now new charges will be added to the staled charges making keeping it pretty much impossible.

We just need to be careful with the static branch activation: since there is no particular preferred order of their activation, we need to make sure that we only start using it after all call sites are active. This is achieved by having a per-memcg flag that is only updated after static_key_slow_inc() returns.

At this time, we are sure all sites are active.

This is made per-memcg, not global, for a reason: it also has the effect of making socket accounting more consistent. The first memcg to be limited will trigger static_key() activation, therefore, accounting. But all the others will then be accounted no matter what. After this patch, only limited memcgs will have its sockets accounted.

```
[v2: changed a tcp limited flag for a generic proto limited flag ]
[v3: update the current active flag only after the static key update ]
[v4: disarm static keys() inside free work ]
[v5: got rid of tcp_limit_mutex, now in the static_key interface ]
[v6: changed active and activated to a flags field, as suggested by akpm ]
[v7: merged more comments from akpm ]
Signed-off-by: Glauber Costa <glommer@parallels.com>
CC: Tejun Heo <tj@kernel.org>
CC: Li Zefan < lizefan@huawei.com>
CC: Kamezawa Hiroyuki <kamezawa.hiroyu@jp.fujitsu.com>
CC: Johannes Weiner <hannes@cmpxchq.org>
CC: Michal Hocko <mhocko@suse.cz>
CC: Andrew Morton <akpm@linux-foundation.org>
include/net/sock.h
                     | 22 ++++++++++++++++++
mm/memcontrol.c
                      3 files changed, 78 insertions(+), 9 deletions(-)
diff --git a/include/net/sock.h b/include/net/sock.h
index b3ebe6b..d6a8ae3 100644
--- a/include/net/sock.h
+++ b/include/net/sock.h
@ @ -46,6 +46,7 @ @
#include linux/list nulls.h>
#include linux/timer.h>
#include linux/cache.h>
+#include linux/bitops.h>
#include linux/lockdep.h>
#include linux/netdevice.h>
#include linux/skbuff.h> /* struct sk buff */
@ @ -907,12 +908,23 @ @ struct proto {
#endif
};
+ * Bits in struct cq proto.flags
+ */
```

```
+enum cg proto flags {
+ /* Currently active and new sockets should be assigned to cgroups */
+ MEMCG_SOCK_ACTIVE,
+ /* It was ever activated; we must disarm static keys on destruction */
+ MEMCG SOCK ACTIVATED.
+};
+
struct cg_proto {
 void (*enter_memory_pressure)(struct sock *sk);
 struct res counter *memory allocated; /* Current allocated memory. */
 struct percpu counter *sockets allocated; /* Current number of sockets. */
 int *memory pressure;
 long *sysctl_mem;
+ unsigned long flags;
 * memcg field is used to find which memcg we belong directly
 * Each memcg struct can hold more than one cg proto, so container of
@ @ -928,6 +940,16 @ @ struct cg proto {
extern int proto register(struct proto *prot, int alloc slab);
extern void proto unregister(struct proto *prot);
+static inline bool memog proto active(struct cg proto *cg proto)
+{
+ return test_bit(MEMCG_SOCK_ACTIVE, &cg_proto->flags);
+}
+static inline bool memcg_proto_activated(struct cg_proto *cg_proto)
+ return test_bit(MEMCG_SOCK_ACTIVATED, &cg_proto->flags);
+}
#ifdef SOCK REFCNT DEBUG
static inline void sk_refcnt_debug_inc(struct sock *sk)
diff --git a/mm/memcontrol.c b/mm/memcontrol.c
index 0b4b4c8..788be2e 100644
--- a/mm/memcontrol.c
+++ b/mm/memcontrol.c
@ @ -404,6 +404,7 @ @ void sock_update_memcg(struct sock *sk)
if (mem cgroup sockets enabled) {
 struct mem_cgroup *memcg;
+ struct cg_proto *cg_proto;
 BUG_ON(!sk->sk_prot->proto_cgroup);
@ @ -423,9 +424,10 @ @ void sock update memcg(struct sock *sk)
```

```
rcu read lock();
 memcg = mem cgroup from task(current);
- if (!mem_cgroup_is_root(memcg)) {
+ cg_proto = sk->sk_prot->proto_cgroup(memcg);
+ if (!mem_cgroup_is_root(memcg) && memcg_proto_active(cg_proto)) {
  mem_cgroup_get(memcg);
- sk->sk cgrp = sk->sk prot->proto cgroup(memcg);
+ sk->sk_cgrp = cg_proto;
 rcu read unlock();
@ @ -454,6 +456,19 @ @ EXPORT SYMBOL(tcp proto cgroup);
#endif /* CONFIG INET */
#endif /* CONFIG_CGROUP_MEM_RES_CTLR_KMEM */
+#if defined(CONFIG_INET) && defined(CONFIG_CGROUP_MEM_RES_CTLR_KMEM)
+static void disarm sock keys(struct mem cgroup *memcg)
+ if (!memcg proto activated(&memcg->tcp mem.cg proto))
+ return;
+ static key slow dec(&memcg socket limit enabled);
+}
+#else
+static void disarm sock keys(struct mem cgroup *memcg)
+}
+#endif
static void drain all stock async(struct mem cgroup *memcg);
static struct mem cgroup per zone *
@ @ -4836,6 +4851,18 @ @ static void free work(struct work struct *work)
 int size = sizeof(struct mem_cgroup);
 memcg = container_of(work, struct mem_cgroup, work_freeing);
+ /*
+ * We need to make sure that (at least for now), the jump label
+ * destruction code runs outside of the cgroup lock. This is because
+ * get online cpus(), which is called from the static branch update,
+ * can't be called inside the cgroup lock, cpusets are the ones
+ * enforcing this dependency, so if they ever change, we might as well.
+ * schedule work() will guarantee this happens. Be careful if you need
+ * to move this code around, and make sure it is outside
+ * the caroup lock.
+ disarm sock keys(memcg);
 if (size < PAGE SIZE)
```

```
kfree(memcg);
diff --git a/net/ipv4/tcp_memcontrol.c b/net/ipv4/tcp_memcontrol.c
index 1517037..b6f3583 100644
--- a/net/ipv4/tcp_memcontrol.c
+++ b/net/ipv4/tcp_memcontrol.c
@ @ -74,9 +74,6 @ @ void tcp_destroy_cgroup(struct mem_cgroup *memcg)
 percpu_counter_destroy(&tcp->tcp_sockets_allocated);
 val = res counter read u64(&tcp->tcp memory allocated, RES LIMIT);
if (val != RESOURCE MAX)
static_key_slow_dec(&memcg_socket_limit_enabled);
EXPORT_SYMBOL(tcp_destroy_cgroup);
@ @ -107,10 +104,33 @ @ static int tcp_update_limit(struct mem_cgroup *memcg, u64 val)
 tcp->tcp prot mem[i] = min t(long, val >> PAGE SHIFT,
      net->ipv4.sysctl tcp mem[i]);
- if (val == RESOURCE MAX && old lim!= RESOURCE MAX)

    static key slow dec(&memcg socket limit enabled);

- else if (old_lim == RESOURCE_MAX && val != RESOURCE_MAX)
static_key_slow_inc(&memcg_socket_limit_enabled);
+ if (val == RESOURCE_MAX)
+ clear_bit(MEMCG_SOCK_ACTIVE, &cg_proto->flags);
+ else if (val != RESOURCE_MAX) {
+ /*
  * The active bit needs to be written after the static key
  * update. This is what guarantees that the socket activation
  * function is the last one to run. See sock update memcg() for
  * details, and note that we don't mark any socket as belonging
  * to this memcg until that flag is up.
  * We need to do this, because static keys will span multiple
  * sites, but we can't control their order. If we mark a socket
  * as accounted, but the accounting functions are not patched in
  * yet, we'll lose accounting.
  * We never race with the readers in sock update memcg(),
   * because when this value change, the code to process it is not
  * patched in yet.
  * The activated bit is used to guarantee that no two writers
  * will do the update in the same memcg. Without that, we can't
  * properly shutdown the static key.
 */
+
+ if (!test and set bit(MEMCG SOCK ACTIVATED, &cg proto->flags))
```

```
+ static_key_slow_inc(&memcg_socket_limit_enabled);
+ set bit(MEMCG SOCK ACTIVE, &cg proto->flags);
+ }
 return 0;
}
1.7.7.6
Subject: Re: [PATCH v7 1/2] Always free struct memcg through schedule_work()
Posted by Michal Hocko on Fri, 25 May 2012 09:50:08 GMT
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On Fri 25-05-12 13:32:07, Glauber Costa wrote:
> Right now we free struct memcg with kfree right after a
> rcu grace period, but defer it if we need to use vfree() to get
> rid of that memory area. We do that by need, because we need vfree
> to be called in a process context.
> This patch unifies this behavior, by ensuring that even kfree will
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> of the complicated and quite far-reaching cgroup lock (that can't be
> held when calling neither the cpu hotplug.lock nor the jump label mutex)
> Signed-off-by: Glauber Costa <glommer@parallels.com>
> Acked-by: Kamezawa Hiroyuki <kamezawa.hiroyu@jp.fujitsu.com>
Acked-by: Michal Hocko <mhocko@suse.cz>
Just one comment below
> CC: Tejun Heo <tj@kernel.org>
> CC: Li Zefan < lizefan@huawei.com>
> CC: Johannes Weiner < hannes@cmpxchg.org>
> CC: Michal Hocko <mhocko@suse.cz>
> CC: Andrew Morton <akpm@linux-foundation.org>
> mm/memcontrol.c | 24 ++++++++++
> 1 files changed, 13 insertions(+), 11 deletions(-)
>
> diff --git a/mm/memcontrol.c b/mm/memcontrol.c
```

> @ @ -4826,23 +4826,28 @ @ out_free:

> index 932a734..0b4b4c8 100644

> --- a/mm/memcontrol.c
> +++ b/mm/memcontrol.c

```
> }
>
> /*
> - * Helpers for freeing a vzalloc()ed mem_cgroup by RCU,
> + * Helpers for freeing a kmalloc()ed/vzalloc()ed mem_cgroup by RCU,
  * but in process context. The work_freeing structure is overlaid
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> -static void vfree work(struct work struct *work)
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> struct mem cgroup *memcg;
> + int size = sizeof(struct mem_cgroup);
> memcg = container_of(work, struct mem_cgroup, work_freeing);
> - vfree(memcg);
> + if (size < PAGE SIZE)
What about
if (is_vmalloc_addr(memcg))
> + kfree(memcg);
> + else
> + vfree(memcg);
> }
Michal Hocko
SUSE Labs
SUSE LINUX s.r.o.
Lihovarska 1060/12
190 00 Praha 9
Czech Republic
```

Subject: Re: [PATCH v7 1/2] Always free struct memcg through schedule_work() Posted by Glauber Costa on Fri, 25 May 2012 09:51:22 GMT

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```
On 05/25/2012 01:50 PM, Michal Hocko wrote:

> On Fri 25-05-12 13:32:07, Glauber Costa wrote:

>> Right now we free struct memcg with kfree right after a

>> rcu grace period, but defer it if we need to use vfree() to get

>> rid of that memory area. We do that by need, because we need vfree

>> to be called in a process context.

>>

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>> happen in a separate thread. The goal is to have a stable place to

>> call the upcoming jump label destruction function outside the realm
```

```
>> of the complicated and quite far-reaching cgroup lock (that can't be
>> held when calling neither the cpu hotplug.lock nor the jump label mutex)
>>
>> Signed-off-by: Glauber Costa<glommer@parallels.com>
>> Acked-by: Kamezawa Hiroyuki<kamezawa.hiroyu@jp.fujitsu.com>
> Acked-by: Michal Hocko<mhocko@suse.cz>
> Just one comment below
>> CC: Tejun Heo<tj@kernel.org>
>> CC: Li Zefan<a href="mailto:lizefan@huawei.com">>> CC: Li Zefan</a><a href="mailto:lizefan@huawei.com">>> CC: Li Zefan</a>
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>> + * Helpers for freeing a kmalloc()ed/vzalloc()ed mem_cgroup by RCU,
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>> -static void vfree_work(struct work_struct *work)
>> +static void free_work(struct work_struct *work)
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>> + int size = sizeof(struct mem_cgroup);
>>
    memcg = container of(work, struct mem cgroup, work freeing);
>>
>> - vfree(memcg);
>> + if (size< PAGE_SIZE)
> What about
> if (is_vmalloc_addr(memcg))
>> + kfree(memcg);
>> + else
>> + vfree(memcg);
```

```
>> }
```

Could be, but I believe this one is already in Andrew's tree from last submission (might be wrong)

Subject: Re: [PATCH v7 2/2] decrement static keys on real destroy time Posted by davem on Sun, 27 May 2012 05:10:49 GMT

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From: Glauber Costa <glommer@parallels.com>

Date: Fri, 25 May 2012 13:32:08 +0400

> Signed-off-by: Glauber Costa <glommer@parallels.com>

Acked-by: David S. Miller <davem@davemloft.net>