## Subject: Re: [PATCH v3 04/13] kmem accounting basic infrastructure Posted by Glauber Costa on Mon, 01 Oct 2012 08:45:11 GMT

View Forum Message <> Reply to Message

```
On 09/30/2012 12:23 PM, Tejun Heo wrote:
> Hello, Glauber.
>
> On Thu, Sep 27, 2012 at 10:30:36PM +0400, Glauber Costa wrote:
>>> But that happens only when pages enter and leave slab and if it still
>>> is significant, we can try to further optimize charging. Given that
>>> this is only for cases where memcg is already in use and we provide a
>>> switch to disable it globally, I really don't think this warrants
>>> implementing fully hierarchy configuration.
>>
>> Not totally true. We still have to match every allocation to the right
>> cache, and that is actually our heaviest hit, responsible for the 2, 3 %
>> we're seeing when this is enabled. It is the kind of path so hot that
>> people frown upon branches being added, so I don't think we'll ever get
>> this close to being free.
> Sure, depening on workload, any addition to alloc/free could be
> noticeable. I don't know. I'll write more about it when replying to
> Michal's message. BTW, __memcg_kmem_get_cache() does seem a bit
> heavy. I wonder whether indexing from cache side would make it
> cheaper? e.g. something like the following.
>
> kmem_cache *__memcg_kmem_get_cache(cachep, gfp)
> {
>
 struct kmem cache *c;
>
  c = cachep->memcg params->caches[percpu read(kmemcg slab idx)];
  if (likely(c))
>
   return c;
  /* try to create and then fall back to cachep */
> }
>
> where kmemcg_slab_idx is updated from sched notifier (or maybe add and
> use current->kmemcg_slab_idx?). You would still need GFP * and
> in interrupt() tests but current->mm and PF KTHREAD tests can be
> rolled into index selection.
>
```

How big would this array be? there can be a lot more kmem\_caches than there are memcgs. That is why it is done from memcg side.