Subject: Re: [PATCH v2 06/11] memcg: kmem controller infrastructure Posted by KAMEZAWA Hiroyuki on Fri, 17 Aug 2012 02:36:26 GMT

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```
(2012/08/13 17:28), Glauber Costa wrote:
>>> + * Needs to be called after memcg_kmem_new_page, regardless of success or
>>> + * failure of the allocation, if @page is NULL, this function will revert the
>>> + * charges. Otherwise, it will commit the memcg given by @handle to the
>>> + * corresponding page cgroup.
>>> + */
>>> +static always inline void
>>> +memcg_kmem_commit_page(struct page *page, struct mem_cgroup *handle, int order)
>>>> +{
>>> + if (memcg_kmem_on)
>>> + __memcg_kmem_commit_page(page, handle, order);
>>>> +}
>> Doesn't this 2 functions has no short-cuts?
> Sorry kame, what exactly do you mean?
I meant avoinding function call. But please ignore, I missed following patches.
>> if (memcg_kmem_on && handle)?
> I guess this can be done to avoid a function call.
>
>> Maybe free() needs to access page_cgroup...
> Can you also be a bit more specific here?
Please ignore, I misunderstood the usage of free_accounted_pages().
>>> +bool __memcg_kmem_new_page(gfp_t gfp, void *_handle, int order)
>>>> +{
>>> + struct mem_cgroup *memcg;
>>>> + struct mem_cgroup **handle = (struct mem_cgroup **)_handle;
>>> + bool ret = true;
>>> + size t size;
>>> + struct task struct *p;
>>>> +
>>> + *handle = NULL;
>>> + rcu read lock():
>>> + p = rcu_dereference(current->mm->owner);
>>> + memcg = mem_cgroup_from_task(p);
>>>> + if (!memcg_kmem_enabled(memcg))
>>> + goto out;
>>>> +
```

```
>>> + mem_cgroup_get(memcg);
>>>> +
>> This mem_cgroup_get() will be a potentioal performance problem.
>> Don't you have good idea to avoid accessing atomic counter here?
>> I think some kind of percpu counter or a feature to disable "move task"
>> will be a help.
>
>
>
>>> + pc = lookup_page_cgroup(page);
>>> + lock page cgroup(pc);
>>> + pc->mem_cgroup = memcg;
>>> + SetPageCgroupUsed(pc);
>>> + unlock_page_cgroup(pc);
>>>> +}
>>>> +
>>> +void __memcg_kmem_free_page(struct page *page, int order)
>>>> +{
>>> + struct mem_cgroup *memcg;
>>> + size t size;
>>> + struct page cgroup *pc;
>>>> +
>>> + if (mem_cgroup_disabled())
>>> + return;
>>>> +
>>> + pc = lookup_page_cgroup(page);
>>> + lock page cgroup(pc);
>>> + memcg = pc->mem cgroup;
>>> + pc->mem_cgroup = NULL;
>> shouldn't this happen after checking "Used" bit?
>> Ah, BTW, why do you need to clear pc->memcg?
> As for clearing pc->memcg, I think I'm just being overzealous. I can't
> foresee any problems due to removing it.
> As for the Used bit, what difference does it make when we clear it?
I just want to see the same logic used in mem_cgroup_uncharge_common().
Hmm, at setting pc->mem cgroup, the things happens in
 set pc->mem_cgroup
 set Used bit
order. If you clear pc->mem_cgroup
 unset Used bit
 clear pc->mem_cgroup
seems reasonable.
```

```
>>> + if (!PageCgroupUsed(pc)) {
>>> + unlock_page_cgroup(pc);
>>> + return;
>>>> + }
>>> + ClearPageCgroupUsed(pc);
>>> + unlock_page_cgroup(pc);
>>>> +
>>> + /*
>>> + * Checking if kmem accounted is enabled won't work for uncharge, since
>>> + * it is possible that the user enabled kmem tracking, allocated, and
>>> + * then disabled it again.
>>>> + *
>>> + * We trust if there is a memcg associated with the page, it is a valid
>>>> + * allocation
>>> + */
>>>> + if (!memcg)
>>> + return;
>>>> +
>>> + WARN_ON(mem_cgroup_is_root(memcg));
>>> + size = (1 << order) << PAGE SHIFT;
>>> + memcg_uncharge_kmem(memcg, size);
>>> + mem_cgroup_put(memcg);
>> Why do we need ref-counting here? kmem res_counter cannot work as
>> reference?
> This is of course the pair of the mem_cgroup_get() you commented on
> earlier. If we need one, we need the other. If we don't need one, we
> don't need the other =)
> The guarantee we're trying to give here is that the memcg structure will
> stay around while there are dangling charges to kmem, that we decided
> not to move (remember: moving it for the stack is simple, for the slab
> is very complicated and ill-defined, and I believe it is better to treat
> all kmem equally here)
> So maybe we can be clever here, and avoid reference counting at all
> times. We call mem_cgroup_get() when the first charge occurs, and then
> go for mem_cgroup_put() when our count reaches 0.
> What do you think about that?
```

I think that should work. I don't want to add not-optimized atomic counter ops in this very hot path.

>

```
>>>> +#ifdef CONFIG MEMCG KMEM
>>> +int memcg_charge_kmem(struct mem_cgroup *memcg, gfp_t gfp, s64 delta)
>>>> +{
>> What does 'delta' means?
>>
> I can change it to something like nr_bytes, more informative.
>>> + struct res_counter *fail_res;
>>> + struct mem cgroup * memcg;
>>>> + int ret:
>>> + bool may oom;
>>> + bool nofail = false:
>>>> +
>>> + may_oom = (gfp & __GFP_WAIT) && (gfp & __GFP_FS) &&
         !(gfp & __GFP_NORETRY);
>>>> +
>>> + ret = 0;
>>>> +
>>> + if (!memcg)
>>> + return ret;
>>>> +
>>> + memcg = memcg;
>>> + ret = __mem_cgroup_try_charge(NULL, gfp, delta / PAGE_SIZE,
         &_memcg, may_oom);
>>>> +
>>>> +
>>> + if (ret == -EINTR) {
>>> + nofail = true;
>>> + /*
>>> + * mem cgroup try charge() chosed to bypass to root due to
>>> + * OOM kill or fatal signal. Since our only options are to
>>> + * either fail the allocation or charge it to this cgroup, do
>>> + * it as a temporary condition. But we can't fail. From a
>>> + * kmem/slab perspective, the cache has already been selected,
>>> + * by mem_cgroup_get_kmem_cache(), so it is too late to change
>>> + * our minds
>>> + */
>>> + res_counter_charge_nofail(&memcg->res, delta, &fail_res);
>>> + if (do swap account)
>>> + res_counter_charge_nofail(&memcg->memsw, delta,
>>>> +
           &fail res):
>>> + ret = 0;
>> Hm, you returns 0 and this charge may never be uncharged....right?
>>
>
> Can't see why. By returning 0 we inform our caller that the allocation
> succeeded. It is up to him to undo it later through a call to uncharge.
>
Hmm, okay. You trust callers.
```

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