Subject: [PATCH v3 2/2] decrement static keys on real destroy time Posted by Glauber Costa on Thu, 26 Apr 2012 21:24:23 GMT

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We call the destroy function when a cgroup starts to be removed, such as by a rmdir event.

However, because of our reference counters, some objects are still inflight. Right now, we are decrementing the static_keys at destroy() time, meaning that if we get rid of the last static_key reference, some objects will still have charges, but the code to properly uncharge them won't be run.

This becomes a problem specially if it is ever enabled again, because now new charges will be added to the staled charges making keeping it pretty much impossible.

We just need to be careful with the static branch activation: since there is no particular preferred order of their activation, we need to make sure that we only start using it after all call sites are active. This is achieved by having a per-memcg flag that is only updated after static_key_slow_inc() returns. At this time, we are sure all sites are active.

This is made per-memcg, not global, for a reason: it also has the effect of making socket accounting more consistent. The first memcg to be limited will trigger static_key() activation, therefore, accounting. But all the others will then be accounted no matter what. After this patch, only limited memcgs will have its sockets accounted.

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[v2: changed a tcp limited flag for a generic proto limited flag ]
[v3: update the current active flag only after the static_key update ]
[v4: disarm_static_keys() inside free_work ]
Signed-off-by: Glauber Costa <glommer@parallels.com>
include/net/sock.h
                      9 +++++
mm/memcontrol.c
                     | 31 +++++++++++++++++
3 files changed, 101 insertions(+), 9 deletions(-)
diff --git a/include/net/sock.h b/include/net/sock.h
index b3ebe6b..c5a2010 100644
--- a/include/net/sock.h
+++ b/include/net/sock.h
@@ -914,6 +914,15 @@ struct cg proto {
int *memory pressure;
```

```
long *sysctl_mem;
+ * active means it is currently active, and new sockets should
  * be assigned to cgroups.
+ * activated means it was ever activated, and we need to
  * disarm the static keys on destruction
+ bool activated;
+ bool active;
+ /*
 * memcg field is used to find which memcg we belong directly
 * Each memcg struct can hold more than one cg_proto, so container_of
 * won't really cut.
diff --git a/mm/memcontrol.c b/mm/memcontrol.c
index b0076cc..53a0815 100644
--- a/mm/memcontrol.c
+++ b/mm/memcontrol.c
@ @ -404,6 +404,7 @ @ void sock update memcg(struct sock *sk)
if (mem_cgroup_sockets_enabled) {
 struct mem_cgroup *memcg;
+ struct cg_proto *cg_proto;
 BUG_ON(!sk->sk_prot->proto_cgroup);
@ @ -423,9 +424,10 @ @ void sock_update_memcg(struct sock *sk)
 rcu read lock();
 memcg = mem_cgroup_from_task(current);
- if (!mem cgroup is root(memcg)) {
+ cg_proto = sk->sk_prot->proto_cgroup(memcg);
+ if (!mem_cgroup_is_root(memcg) && cg_proto->active) {
  mem_cgroup_get(memcg);
- sk->sk_cgrp = sk->sk_prot->proto_cgroup(memcg);
+ sk->sk carp = ca proto:
 rcu read unlock();
@ @ -442,6 +444,14 @ @ void sock release memcg(struct sock *sk)
}
+static void disarm_static_keys(struct mem_cgroup *memcg)
+{
+#ifdef CONFIG_INET
+ if (memcg->tcp_mem.cg_proto.activated)
+ static key slow dec(&memcg socket limit enabled);
```

```
+#endif
+}
+
#ifdef CONFIG INET
struct cg_proto *tcp_proto_cgroup(struct mem_cgroup *memcq)
@ @ -452,6 +462,11 @ @ struct cg_proto *tcp_proto_cgroup(struct mem_cgroup *memcg)
EXPORT SYMBOL(tcp proto cgroup);
#endif /* CONFIG INET */
+#else
+static inline void disarm static keys(struct mem cgroup *memcg)
+}
#endif /* CONFIG_CGROUP_MEM_RES_CTLR_KMEM */
static void drain all stock async(struct mem cgroup *memcg);
@ @ -4836,6 +4851,18 @ @ static void free work(struct work struct *work)
 int size = sizeof(struct mem_cgroup);
 memcg = container of(work, struct mem cgroup, work freeing);
+ /*
+ * We need to make sure that (at least for now), the jump label
+ * destruction code runs outside of the cgroup lock. It is in theory
+ * possible to call the cgroup destruction function outside of that
+ * lock, but it is not yet done. rate limiting plus the deferred
+ * interface for static branch destruction guarantees that it will
+ * run through schedule work(), therefore, not holding any cgroup
+ * related lock (this is, of course, until someone decides to write
+ * a schedule_work cgroup :p )
+ */
+ disarm_static_keys(memcg);
 if (size < PAGE_SIZE)
 kfree(memcg):
 else
diff --git a/net/ipv4/tcp memcontrol.c b/net/ipv4/tcp memcontrol.c
index 1517037..7790008 100644
--- a/net/ipv4/tcp memcontrol.c
+++ b/net/ipv4/tcp memcontrol.c
@ @ -54,6 +54,8 @ @ int tcp_init_cgroup(struct mem_cgroup *memcg, struct cgroup_subsys *ss)
 cg proto->sysctl mem = tcp->tcp prot mem;
 cg_proto->memory_allocated = &tcp->tcp_memory_allocated;
 cg proto->sockets allocated = &tcp->tcp sockets allocated:
+ cg_proto->active = false;
+ cq proto->activated = false;
 cg proto->memcg = memcg;
```

```
return 0:
@ @ -74,12 +76,43 @ @ void tcp_destroy_cgroup(struct mem_cgroup *memcg)
 percpu_counter_destroy(&tcp->tcp_sockets_allocated);
 val = res_counter_read_u64(&tcp->tcp_memory_allocated, RES_LIMIT);
- if (val != RESOURCE MAX)
- static key slow dec(&memcg socket limit enabled);
EXPORT_SYMBOL(tcp_destroy_cgroup);
+/*
+ * This is to prevent two writes arriving at the same time
+ * at kmem.tcp.limit_in_bytes.
+ * There is a race at the first time we write to this file:
+ * - cq proto->activated == false for all writers.
+ * - They all do a static key slow inc().
+ * - When we are finally read to decrement the static keys,
    we'll do it only once per activated cgroup. So we won't
+ * be able to disable it.
    Also, after the first caller increments the static_branch
    counter, all others will return right away. That does not mean,
    however, that the update is finished.
+ * Without this mutex, it would then be possible for a second writer
    to get to the update site, return
    When a user updates limit of 2 cgroups at once, following happens.
    CPU A CPU B
+ * if (cg_proto->activated) if (cg->proto_activated)
+ * static_key_inc() static_key_inc()
    => set counter 0->1 => set counter 1->2,
       return immediately.
    => hold mutex => cg_proto->activated = true.
    => overwrite imps.
+ * This race was described by Kamezawa Hiroyuki.
+ */
+static DEFINE_MUTEX(tcp_set_limit_mutex);
static int tcp_update_limit(struct mem_cgroup *memcg, u64 val)
{
```

```
struct net *net = current->nsproxy->net ns;
@@ -107,10 +140,33 @@ static int tcp_update_limit(struct mem_cgroup *memcg, u64 val)
 tcp->tcp_prot_mem[i] = min_t(long, val >> PAGE_SHIFT,
      net->ipv4.sysctl_tcp_mem[i]);
- if (val == RESOURCE_MAX && old_lim != RESOURCE_MAX)
- static key slow dec(&memcg socket limit enabled):
- else if (old_lim == RESOURCE_MAX && val != RESOURCE_MAX)
- static key slow inc(&memcg socket limit enabled);
+ if (val == RESOURCE MAX)
+ cq proto->active = false;
+ else if (val != RESOURCE MAX) {
+ /*
  * ->activated needs to be written after the static key update.
  * This is what guarantees that the socket activation function
   * is the last one to run. See sock_update_memcg() for details,
    and note that we don't mark any socket as belonging to this
    memcg until that flag is up.
+
  * We need to do this, because static keys will span multiple
  * sites, but we can't control their order. If we mark a socket
+
  * as accounted, but the accounting functions are not patched in
  * yet, we'll lose accounting.
+
  * We never race with the readers in sock_update_memcg(), because
  * when this value change, the code to process it is not patched in
  * yet.
+
  */
+ mutex lock(&tcp set limit mutex);
+ if (!cq proto->activated) {
+ static key slow inc(&memcg socket limit enabled);
+ cg_proto->activated = true;
+ }
+ mutex_unlock(&tcp_set_limit_mutex);
+ cg_proto->active = true;
+ }
 return 0;
}
1.7.7.6
```