Subject: Re: [PATCH 1/1, v6] cgroup/freezer: add per freezer duty ratio control Posted by jacob.jun.pan on Thu, 10 Feb 2011 23:06:33 GMT

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On Wed, 9 Feb 2011 19:04:42 -0800 Matt Helsley <matthltc@us.ibm.com> wrote:

- > On Tue, Feb 08, 2011 at 05:05:41PM -0800,
- > jacob.jun.pan@linux.intel.com wrote:
- > > From: Jacob Pan < jacob.jun.pan@linux.intel.com>

> >

- >> Freezer subsystem is used to manage batch jobs which can start
- > > stop at the same time. However, sometime it is desirable to let
- > > the kernel manage the freezer state automatically with a given
- > > duty ratio.
- > > For example, if we want to reduce the time that backgroup apps
- > > are allowed to run we can put them into a freezer subsystem and
- > > set the kernel to turn them THAWED/FROZEN at given duty ratio.

> >

- > > This patch introduces two file nodes under cgroup
- >> freezer.duty\_ratio\_pct and freezer.period\_sec

> >

- >> Usage example: set period to be 5 seconds and frozen duty ratio 90%
- >> [root@localhost aoa]# echo 90 > freezer.duty\_ratio\_pct
- >> [root@localhost aoa]# echo 5000 > freezer.period\_ms

>

- > I kept wondering how this was useful when we've got the "cpu"
- > subsystem because for some reason "duty cycle" made me think this was
- > a scheduling policy knob. In fact, I'm pretty sure it is -- it just
- > happens to sometimes reduce power consumption.

>

- > Have you tried using the cpu cgroup subsystem's share to see if it can
- > have a similar effect?

>

- > Can you modify the cpu subsystem to enable this instead of putting it
- > into the cgroup freezer subsystem?

>

I replied in other email. basically, CPU subsystem is for RT only so far. I will give it a try see if it can include non-RT tasks and perform with CFS.

- > The way it oscillates between FROZEN and THAWED also bothers me. The
- > oscillations can be described in millisecond granularity so its
- > possible that reading and manipulating the freezer state from
- > userspace could be largely useless. Also it's not obvious what should
- > happen when the state file is written after the duty cycle has been
- > set (more below).

>

My intention was to have second granularity.

> Perhaps you could fix that up by introducting another state called > "DUTY\_CYCLE" or something. > I did think about that as well. But adding DUTY\_CYCLE state kind of blurs the state machine definition. Since it is can be in THAWED or FROZEN while in DUTY\_CYCLE. But I do need to fix the handling of user direct control of freezer.state while in oscillation. > What's the overhead of using the freezer as a scheduling mechanism at > that granularity? Is it really practical? > I agree at ms granularity the overhead is not practical. Like Arjan said we are looking at much longer time at 20s+, as long as the apps in the freezer can be kept alive :). > What happens to these groups using the duty cycle during suspend and > resume? Presumably they won't be accidentally thawed so long as there > aren't races between the kernel thread(s) and suspend. I don't think > we've ever had a kernel thread that could thaw a frozen task before > (unless it's part of the resume code itself) so I don't think this > race is covered by existing cgroup freezer code. good point, I need to do some investigation and get back to you. > Overall I get the feeling this is a scheduling policy knob that > doesn't "belong" in the cgroup freezer subsystem -- though I don't > have much beyond the above questions and my personal aesthetic sense > to go on :). > I think Rafael is maintaining the cgroup freezer subsystem since it > makes use of the suspend freezer so I've added him to Cc. Thanks for the pointer. As I mentioned in the other reply, cpu cgroup subsystem might be a more natural fit but we may need to overcome the hurdle or non-rt and possible scheduling heuristics. I need to investigate some more. > > Signed-off-by: Jacob Pan < jacob.jun.pan@linux.intel.com> >> --->> Documentation/cgroups/freezer-subsystem.txt | 23 +++++ >> kernel/cgroup\_freezer.c | 132 > > deletions(-)

> > diff --git a/Documentation/cgroups/freezer-subsystem.txt > > b/Documentation/cgroups/freezer-subsystem.txt index

```
> > 41f37fe..7f06f05 100644 ---
> > a/Documentation/cgroups/freezer-subsystem.txt +++
>> b/Documentation/cgroups/freezer-subsystem.txt @ @ -100,3 +100,26 @ @
> > things happens: and returns EINVAL)
>> 3) The tasks that blocked the cgroup from entering the
>> "FROZEN" state disappear from the cgroup's set of tasks.
> > +
>> +In embedded systems, it is desirable to manage group of
> > applications +for power saving. E.g. tasks that are not in the
> > foreground may be +frozen unfrozen periodically to save power
> > without affecting user
> nit: probably should be "frozen and unfrozen periodically"
>
>> +experience. In this case, user/management software can attach tasks
>> +into freezer cgroup then specify duty ratio and period that the
> > +managed tasks are allowed to run.
> And presumably the applications either don't care about their power
> consumption, have a bug, or are "malicious" apps -- either way
> assuming cooperation from the applications and knowledgable users
> isn't acceptable.
>
> > +
> > +Usage example:
>> +Assuming freezer cgroup is already mounted, application being
> > managed +are included the "tasks" file node of the given freezer
>> cgroup. +To make the tasks frozen at 90% of the time every 5
> > seconds, do: +
>> +[root@localhost]# echo 90 > freezer.duty_ratio_pct
>> +[root@localhost]# echo 5000 > freezer.period_ms
> > +
>> +After that, the application in this freezer cgroup will only be
> > +allowed to run at the following pattern.
> > +
       | |<-- 90% frozen -->| |
> > +
>>+ |<---- 5 seconds ---->|
> > diff --git a/kernel/cgroup_freezer.c b/kernel/cgroup_freezer.c
> > index e7bebb7..5808f28 100644
>> --- a/kernel/cgroup_freezer.c
>> +++ b/kernel/cgroup_freezer.c
>> @ @ -21,6 +21,7 @ @
>> #include ux/uaccess.h>
>> #include ux/freezer.h>
>> #include ux/seq file.h>
> > +#include ux/kthread.h>
```

```
> >
>> enum freezer state {
>> CGROUP_THAWED = 0,
>> @ @ -28,12 +29,28 @ @ enum freezer_state {
>> CGROUP_FROZEN,
>> };
> >
>> +enum duty_ratio_params {
>> + FREEZER DUTY RATIO = 0,
>> + FREEZER PERIOD,
> > +};
> > +
>> +struct freezer_duty {
>> + u32 ratio; /* percentage of time frozen */
> > + u32 period_pct_ms; /* one percent of the period in
> > miliseconds */ +};
> > +
>> struct freezer {
>> struct cgroup_subsys_state css;
>> enum freezer state state;
>> + struct freezer duty duty;
>> + struct task struct *fkh;
>> spinlock_t lock; /* protects _writes_ to state */
>> };
> >
>> +static struct task_struct *freezer_task;
>> +static int try_to_freeze_cgroup(struct cgroup *cgroup, struct
>> freezer *freezer); +static void unfreeze cgroup(struct cgroup
> > *cgroup, struct freezer *freezer); +
>> static inline struct freezer *cgroup_freezer(
>> struct cgroup *cgroup)
>> {
>> @ @ -63,6 +80,31 @ @ int cgroup_freezing_or_frozen(struct task_struct
> > *task) return result;
>> }
>> +static DECLARE_WAIT_QUEUE_HEAD(freezer_wait);
> > +
> > +static int freezer_kh(void *data)
> nit: What's "kh"? "Kernel Handler"?
I meant kernel thread:)
> > +{
>> + struct cgroup *cgroup = (struct cgroup *)data;
> > + struct freezer *freezer = cgroup_freezer(cgroup);
> > +
> > + do {
```

```
>> + if (freezer->duty.ratio < 100 &&
> > freezer->duty.ratio > 0 &&
>> + freezer->duty.period_pct_ms) {
>> + if (try_to_freeze_cgroup(cgroup, freezer))
>> + pr_info("cannot freeze\n");
>> + msleep(freezer->duty.period_pct_ms *
>> + freezer->duty.ratio);
>> + unfreeze_cgroup(cgroup, freezer);
>> + msleep(freezer->duty.period_pct_ms *
> > + (100 - freezer->duty.ratio));
>>+ } else {
>> + sleep on(&freezer wait);
>> + pr_debug("freezer thread wake up\n");
>>+ }
>> + } while (!kthread_should_stop());
> > + return 0;
> > +}
> Seems to me you could avoid the thread-per-cgroup overhead and the
> sleep-loop code by using one timer-per-cgroup. When the timer expires
> you freeze/thaw the cgroup associated with the timer, setup the next
> wakeup timer, and use only one kernel thread to do it all. If you
> use workqueues you might even avoid the single kernel thread.
> Seems to me like that'd be a good fit for embedded devices.
will try schedule_delayed_work() as Kirill suggested.
> > +
>> /*
>> * cgroups write string() limits the size of freezer state strings
> to
>> * CGROUP_LOCAL_BUFFER_SIZE
>> @ @ -150,7 +192,12 @ @ static struct cgroup_subsys_state
> > *freezer_create(struct cgroup_subsys *ss, static void
>> freezer_destroy(struct cgroup_subsys *ss, struct cgroup *cgroup)
>> - kfree(cgroup freezer(cgroup));
>> + struct freezer *freezer;
>> + freezer = cgroup freezer(cgroup);
>> + if (freezer->fkh)
> > + kthread_stop(freezer->fkh);
>> + kfree(freezer);
>> }
> >
>> /*
>> @ @ -282,6 +329,16 @ @ static int freezer read(struct cgroup *cgroup,
```

```
> > struct cftype *cft, return 0;
>> }
> >
>> +static u64 freezer_read_duty_ratio(struct cgroup *cgroup, struct
> > cftype *cft) +{
>> + return cgroup_freezer(cgroup)->duty.ratio;
> > +}
> > +
>> +static u64 freezer read period(struct cgroup *cgroup, struct
> > cftype *cft) +{
>> + return cgroup_freezer(cgroup)->duty.period_pct_ms * 100;
> > +
>> static int try_to_freeze_cgroup(struct cgroup *cgroup, struct
> > freezer *freezer) {
>> struct cgroup_iter it;
>> @ @ -368,12 +425,85 @ @ static int freezer write(struct cgroup
> > *cgroup, return retval;
>> }
> >
>> +#define FREEZER KH PREFIX "freezer "
>> +static int freezer write param(struct cgroup *cgroup, struct
> > cftype *cft,
>>+ u64 val)
> > +{
>> + struct freezer *freezer;
> > + char thread_name[32];
> + int ret = 0;
> > +
>> + freezer = cgroup_freezer(cgroup);
>> + if (!cgroup_lock_live_group(cgroup))
>> + return -ENODEV;
>> + switch (cft->private) {
>> + case FREEZER DUTY RATIO:
>>+ if (val >=100 || val < 0) {
>>+ ret = -EINVAL;
>> + goto exit;
>>+ }
>> + freezer->duty.ratio = val;
> Why can't val == 100? At that point it's always THAWED and no kernel
> thread is necessary (just like at 0 it's always FROZEN and no kernel
> thread is necessary).
the val is percentage of time FROZEN. in that case user can just change
freezer.state to FROZEN.
```

```
>
>> + break;
> > + case FREEZER_PERIOD:
> > + if (val)
> > + do_div(val, 100);
>> + freezer->duty.period_pct_ms = val;
>
> Wrong indent level at least. Possible bug?
> Shouldn't you disallow duty.period pct ms being set to 0? Then
> userspace can pin a kernel thread at 100% cpu just doing freeze/thaws
> couldn't it?
I will fix that, no need to check val != 0.
> > + break;
>> + default:
>> + BUG();
> > + }
> > +
>> + /* start/stop management kthread as needed, the rule is
> > that
>> + * if both duty ratio and period values are zero, then no
> > management
>> + * kthread is created. when both are non-zero, we create a
> > kthread
>> + * for the cgroup. When user set zero to duty ratio and
> > period again
>> + * the kthread is stopped.
>> + */
>> + if (freezer->duty.ratio && freezer->duty.period pct ms) {
> > + if (!freezer->fkh) {
>> + snprintf(thread name, 32, "%s%s",
> > FREEZER_KH_PREFIX,
> > + cgroup->dentry->d_name.name);
>> + freezer->fkh = kthread_run(freezer_kh,
> > (void *)cgroup,
        thread name);
>> + if (IS_ERR(freezer_task)) {
>> + pr_err("create %s failed\n",
> > thread name);
       ret = PTR ERR(freezer task);
> > +
       goto exit;
> > +
>>+ }
>>+ } else
> > + wake_up(&freezer_wait);
>> + } else if ((!freezer->duty.ratio
>> || !freezer->duty.period_pct_ms) &&
> > + freezer->fkh) {
>> + kthread stop(freezer->fkh);
```

```
>> + freezer->fkh = NULL;
> > + }
> > +
> > +exit:
> > + cgroup_unlock();
>> + return ret;
> > +}
> > +
>> static struct cftype files[] = {
>> {
>> .name = "state",
>> .read seg string = freezer read,
     .write_string = freezer_write,
> It's not clear what should happen when userspace writes the state
> file after writing a duty_ratio_pct.
> If the new state file write takes priority then:
> Writing THAWED to the state should set duty ratio pct to 100.
> Writing FROZEN to the state should set it to 0.
> This means existing code will get the behavior it expects.
>
> Else, if you want duty_ratio_pct to take priority then you ought to
> make the state file read-only when duty_ratio_pct is set. Otherwise
> existing userspace code will happily chug along without noticing that
> their groups aren't doing what they expected. This is also another
> good reason to introduce a new state as suggested above (with the
> tenative name "DUTY CYCLE").
I like the former logic, where freezer.state takes precedence. As i
mentioned before, my concern is that DUTY CYCLE state overlaps THAWED
and FROZEN states.
Thanks.
Containers mailing list
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