Subject: Re: [PATCH 0/5] blk-throttle: writeback and swap IO control Posted by KAMEZAWA Hiroyuki on Thu, 24 Feb 2011 00:40:39 GMT

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On Wed, 23 Feb 2011 19:10:33 -0500 > On Thu, Feb 24, 2011 at 12:14:11AM +0100, Andrea Righi wrote: >> On Wed, Feb 23, 2011 at 10:23:54AM -0500, Vivek Goyal wrote: >>> Agreed. Granularity of per inode level might be accetable in many >>> cases. Again, I am worried faster group getting stuck behind slower >>> > qroup. >>>> >>>> I am wondering if we are trying to solve the problem of ASYNC write throttling >>>> at wrong layer. Should ASYNC IO be throttled before we allow task to write to >>> page cache. The way we throttle the process based on dirty ratio, can we >>> > just check for throttle limits also there or something like that.(I think >>>> that's what you had done in your initial throttling controller implementation?) >>>> >>> Right. This is exactly the same approach I've used in my old throttling >>> controller: throttle sync READs and WRITEs at the block layer and async >>> WRITEs when the task is dirtying memory pages. >>> This is probably the simplest way to resolve the problem of faster group >>> getting blocked by slower group, but the controller will be a little bit >>> more leaky, because the writeback IO will be never throttled and we'll >>> see some limited IO spikes during the writeback. >> Yes writeback will not be throttled. Not sure how big a problem that is. >> - We have controlled the input rate. So that should help a bit. >> - May be one can put some high limit on root cgroup to in blkio throttle >>> controller to limit overall WRITE rate of the system. >> - For SATA disks, try to use CFQ which can try to minimize the impact of >>> WRITE. >>> It will atleast provide consistent bandwindth experience to application. > > > > Right. > > >>> >>> However, this is always >>> a better solution IMHO respect to the current implementation that is >>> affected by that kind of priority inversion problem. >>>> >>> I can try to add this logic to the current blk-throttle controller if >>> you think it is worth to test it.

>>>

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>>> At this point of time I have few concerns with this approach.
>>>
>>> - Configuration issues. Asking user to plan for SYNC ans ASYNC IO
      separately is inconvenient. One has to know the nature of workload.
>>>
>> - Most likely we will come up with global limits (atleast to begin with),
>>> and not per device limit. That can lead to contention on one single
>>> lock and scalability issues on big systems.
>>>
>> Having said that, this approach should reduce the kernel complexity a lot.
>> So if we can do some intelligent locking to limit the overhead then it
>>> will boil down to reduced complexity in kernel vs ease of use to user. I
>> guess at this point of time I am inclined towards keeping it simple in
> > > kernel.
>>>
> >
>> BTW, with this approach probably we can even get rid of the page
> > tracking stuff for now.
>
> Agreed.
>> If we don't consider the swap IO, any other IO
> > operation from our point of view will happen directly from process
>> context (writes in memory + sync reads from the block device).
> Why do we need to account for swap IO? Application never asked for swap
> IO. It is kernel's decision to move soem pages to swap to free up some
> memory. What's the point in charging those pages to application group
> and throttle accordingly?
>
```

I think swap I/O should be controlled by memcg's dirty_ratio. But, IIRC, NEC guy had a requirement for this...

I think some enterprise cusotmer may want to throttle the whole speed of swapout I/O (not swapin)...so, they may be glad if they can limit throttle the I/O against a disk partition or all I/O tagged as 'swapio' rather than some cgroup name.

But I'm afraid slow swapout may consume much dirty ratio and make things worse;)

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> > However, I'm sure we'll need the page tracking also for the blkio
>> controller soon or later. This is an important information and also the
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> > proportional bandwidth controller can take advantage of it.

> >

- > Yes page tracking will be needed for CFQ proportional bandwidth ASYNC
- > write support. But until and unless we implement memory cgroup dirty
- > ratio and figure a way out to make writeback logic cgroup aware, till
- > then I think page tracking stuff is not really useful.

>

I think Greg Thelen is now preparing patches for dirty_ratio.

Thanks,

-Kame

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