Subject: Re: [RFC] Transactional CGroup task attachment Posted by Daisuke Nishimura on Mon, 14 Jul 2008 12:36:09 GMT

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On Mon, 14 Jul 2008 16:54:44 +0900, KAMEZAWA Hiroyuki <kamezawa.hiroyu@jp.fujitsu.com> wrote: > On Mon, 14 Jul 2008 15:28:22 +0900 > Daisuke Nishimura <nishimura@mxp.nes.nec.co.jp> wrote: > > On Fri, 11 Jul 2008 09:20:58 +0900, KAMEZAWA Hiroyuki <kamezawa.hiroyu@jp.fujitsu.com> wrote: >>> Thank you for your effort. >>> > > On Wed, 9 Jul 2008 23:46:33 -0700 >> "Paul Menage" <menage@google.com> wrote: > > > 3) memory >>>> >>> Curently the memory cgroup only uses the mm->owner's cgroup at charge >>> time, and keeps a reference to the cgroup on the page. However, >>> patches have been proposed that would move all non-shared (page count >>> == 1) pages to the destination cgroup when the mm->owner moves to a >>> new cgroup. Since it's not possible to prevent page count changes >>> without locking all mms on the system, even this transaction approach >>> can't really give guarantees. However, something like the following >>> would probably be suitable. It's very similar to the memrlimit >>> approach, except for the fact that we have to handle the fact that the >>> number of pages we finally move might not be exactly the same as the >>> number of pages we thought we'd be moving. >>>> >>> prepare_attach_sleep() { >>> down read(&mm->mmap sem); >>> if (mm->owner != state->task) return 0; >>> count = count_unshared_pages(mm); >>> // save the count charged to the new cgroup >>> state->subsys[memcgroup_subsys_id] = (void *)count; >>> if ((ret = res_counter_charge(state->dest, count)) { >>>> up_read(&mm->mmap_sem); >>>> } >>> return ret; >>>>} >>>> > > > commit_attach() { >>> if (mm->owner == state->task) { final_count = move_unshared_pages(mm, state->dest); >>>> res_counter_uncharge(state->src, final_count); >>>> count = state->subsys[memcgroup_subsys_id]; >>>> res_counter_force_charge(state->dest, final_count - count); >>>> >>>> }

```
>>> up_read(&mm->mmap_sem);
>>>>}
>>>>
>>> abort_attach_sleep() {
>>> if (mm->owner == state->task) {
         count = state->subsys[memcgroup_subsys_id];
>>>>
         res_counter_uncharge(state->dest, count);
>>>>
>>>> }
>>> up_read(&mm->mmap_sem);
>>>>}
>>>>
>>>
>>> At frist look, maybe works well. we need some special codes (to move resource)
> > but that's all.
>>>
>> My small concern is a state change between prepare_attach_sleep() ->
>> commit attach(). Hmm...but as you say, we cannot do down write(mmap sem).
>> Maybe inserting some check codes to mem_cgroup_charge() to stop charge while
>>> move is the last thing we can do.
> > I have two comments.
>> - I think page reclaiming code decreases the memory charge
>> without holding mmap_sem(e.g. try_to_unmap(), __remove_mapping()).
>> Shouldn't we handle these cases?
>
> I think decreasing is not problem, here.
> I don't like handle mmap->sem by some unclear way. I'd like to add some flag to
> mm struct or page struct to stop(skip/avoid) charge/uncharge while task move.
It would be a good idea.
>> - When swap controller is merged, I should implement
>> prepare_attach_nosleep() which holds swap_lock.
> >
> just making add_to_swap() fail during move is not enough?
This can only avoid increasing, I think.
I thought it would be better to avoid decreasing too,
iust because some special handling on uncharged usage
```

would be needed in rollback or commit.

Anyway, I think it depends on how to implement move and rollback, and I will consider more.

Thank you for your suggestion.

Thanks, Daisuke Nishimura.

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