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Subject: Re: design of user namespaces

Posted by [serue](#) on Mon, 30 Jun 2008 21:13:25 GMT

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Quoting Eric W. Biederman (ebiederm@xmission.com):

> "Serge E. Hallyn" <serue@us.ibm.com> writes:

>

> > There are. But one key point is that the namespace ids are not  
> > cryptographic keys. They don't need to be globally unique, or even 100%  
> > unique on one system (though that gets too subtle).

> >

> > I was wanting to keep them tiny - or at least variably sized - so they  
> > could be stored with each inode.

>

> Sure. I think they make a lot of sense. idmapd uses domain names  
> for this purpose. At the moment I just don't think they are necessary.

> Like veth isn't necessary for network namespaces. Ultimately we  
> will use these identifiers all of the time but it doesn't mean  
> the generic code has to deal with them.

>

>

> >> In particular: Is this user allowed to use this ID?

> >

> > One way to address that is actually by having a system-wide tool with  
> > CAP\_SET\_USERNS=fp which enforces a userid-to-unsid mapping. The host  
> > sysadmin creates a table, say, 500:0 may use unsid 10-15. 500:0 (let's  
> > call him hallyn) doesn't have CAP\_SET\_USERNS permissions himself, but  
> > can run /bin/set\_unsid, which runs with CAP\_SET\_USERNS and ensures that  
> > hallyn uses only unsids 10-15.

> >

> > It's not ideal. I'd rather have some sort of fancy collision-proof  
> > global persistent id system, but again I think it's important that if  
> > 500:0 creates userns 1 and 400:1 creates userns 2, and 0:2 creates a  
> > file, that the file be persistently marked as belonging to  
> > (500:0,400:1,0:2), distinct from another file created by  
> > (500:0,400:1,1000:2). Which means these things have to be stored  
> > per-inode, meaning they can't be too large.

>

> So my suggestion was something like this:

> mount -o nativemount,uidns=1 /

>

> Then the filesystem performs magic to ask if the owner of user namespace  
> is allowed to use uidns 1. That magic would consult a config file like:

>

> [domains]

> local1.mydomain 1

> local2.mydomain 2

> local3.mydomain 3

>  
> [users]  
> bob local1.mydomain  
> bob local3.mydomain  
> nancy local2.mydomain  
>  
> Or something like that. Reporting which users are allowed to use  
> which userid namespaces, and the mapping of those userid namespaces  
> to something compact for storing in the filesystem.  
>  
> The magic could be an upcall to userspace.  
> The magic could be loading the configuration file at mount time.  
> The magic could be storing the config file in the filesystem  
> and having special commands to access it like quotas.  
>  
> The very important points are that it is a remount of an existing mount  
> so that we don't have to worry about corrupted filesystem attacks, and  
> that authentication is performed at mount time.

Conceptually that (making corrupted fs attacks a non-issue) is wonderful. Practically, I may be missing something: When you say remount, it seems you must either mean a bind mount or a remount. If remount, then that will want to change superblock flags. If the child users(+child mntns) does a real remount, then that will change the flags for the parent ns as well, right?

If instead we do a bind mount we don't have that problem, but then the fs can't be the one doing the user namespace work.

I'm probably missing something...

> I just think that once we get to the point of specifying the parameters at  
> mount time. There is no need for generic kernel configuration of a  
> uidns name.

thanks,  
-serge

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