Subject: Re: [RFC/PATCH 1/8]: CGroup Files: Add locking mode to cgroups control files

Posted by Matt Helsley on Tue, 13 May 2008 20:38:58 GMT

View Forum Message <> Reply to Message

```
On Tue, 2008-05-13 at 13:01 -0700, Andrew Morton wrote:
> Fear, doubt and resistance!
> On Mon, 12 May 2008 23:37:08 -0700
> menage@google.com wrote:
>> Different cgroup files have different stability requirements of the
>> cgroups framework while the handler is running; currently most
> > subsystems that don't have their own internal synchronization just
> > call cgroup_lock()/cgroup_unlock(), which takes the global cgroup_mutex.
> >
>> This patch introduces a range of locking modes that can be requested
>> by a control file; currently these are all implemented internally by
>> taking cgroup mutex, but expressing the intention will make it simpler
>> to move to a finer-grained locking scheme in the future.
> >
>
> This, umm, doesn't seem to do much to make the kernel a simpler place.
> Do we expect to gain much from this? Hope so... What?
>
> > Index: cgroup-2.6.25-mm1/include/linux/cgroup.h
>> --- cgroup-2.6.25-mm1.orig/include/linux/cgroup.h
>> +++ cgroup-2.6.25-mm1/include/linux/cgroup.h
>> @ @ -200,11 +200,87 @ @ struct cgroup map cb {
>> */
>> #define MAX_CFTYPE_NAME 64
> > +
> > +/* locking modes for control files.
>> + * These determine what level of quarantee the file handler wishes
>> + * cgroups to provide about the stability of control group entities
> > + * for the duration of the handler callback.
>>+ *
>> + * The minimum guarantee is that the subsystem state for this
>> + * subsystem will not be freed (via a call to the subsystem's
> > + * destroy() callback) until after the control file handler
>> + * returns. This guarantee is provided by the fact that the open
>> + * dentry for the control file keeps its parent (cgroup) dentry alive,
>> + * which in turn keeps the cgroup object from being actually freed
>> + * (although it can be moved into the removed state in the
```

```
> > + * meantime). This is suitable for subsystems that completely control
>> + * their own synchronization.
> > + *
>> + * Other possible guarantees are given below.
>> + * XXX_READ bits are used for a read operation on the control file,
>> + * XXX_WRITE bits are used for a write operation on the control file
> Vague handwaving: lockdep doesn't know anything about any of this.
> Whereas if we were more conventional in using separate locks and
> suitable lock types for each application, we would retain full lockdep
> coverage.
>
> > +/*
>> + * CFT_LOCK_ATTACH_(READ|WRITE): This operation will not run
>> + * concurrently with a task movement into or out of this cgroup.
> > + */
>> +#define CFT LOCK ATTACH READ 1
>> +#define CFT LOCK ATTACH WRITE 2
>> +#define CFT LOCK_ATTACH (CFT_LOCK_ATTACH_READ |
CFT LOCK ATTACH WRITE)
> > +
> > +/*
>> + * CFT_LOCK_RMDIR_(READ|WRITE): This operation will not run
>> + * concurrently with the removal of the affected cgroup.
> > + */
>> +#define CFT LOCK RMDIR READ 4
>> +#define CFT LOCK RMDIR WRITE 8
>> +#define CFT LOCK RMDIR (CFT LOCK RMDIR READ | CFT LOCK RMDIR WRITE)
> > +
> > +/*
>> + * CFT_LOCK_HIERARCHY_(READ|WRITE): This operation will not run
>> + * concurrently with a cgroup creation or removal in this hierarchy,
> > + * or a bind/move/unbind for this subsystem.
> > + */
> > +#define CFT_LOCK_HIERARCHY_READ 16
>> +#define CFT LOCK HIERARCHY WRITE 32
>> +#define CFT_LOCK_HIERARCHY (CFT_LOCK_HIERARCHY_READ |
CFT LOCK HIERARCHY WRITE)
> > +
> > +/*
>> + * CFT_LOCK_CGL_(READ|WRITE): This operation is called with
>> + * cgroup_lock() held; it will not run concurrently with any of the
>> + * above operations in any cgroup/hierarchy. This should be considered
>> + * to be the BKL of cgroups - it should be avoided if you can use
>> + * finer-grained locking
> > + */
```

```
>> +#define CFT LOCK CGL READ 64
>> +#define CFT LOCK CGL WRITE 128
>> +#define CFT_LOCK_CGL (CFT_LOCK_CGL_READ | CFT_LOCK_CGL_WRITE)
> > +
>> +#define CFT_LOCK_FOR_READ (CFT_LOCK_ATTACH_READ | \
        CFT_LOCK_RMDIR_READ | \
        CFT LOCK HIERARCHY READ | \
> > +
        CFT_LOCK_CGL_READ)
> > +
> > +
>> +#define CFT LOCK FOR_WRITE (CFT_LOCK_ATTACH_WRITE | \
        CFT LOCK RMDIR WRITE | \
> > +
        CFT LOCK HIERARCHY WRITE I \
> > +
        CFT_LOCK_CGL_WRITE)
>>+
> > +
>> struct cftype {
>> /* By convention, the name should begin with the name of the
>> * subsystem, followed by a period */
>> char name[MAX CFTYPE NAME];
>> int private;
> > +
> > + /*
>> + * Determine what locks (if any) are held across calls to
>> + * read X/write X callback. See lockmode definitions above
>> + */
>> + int lockmode;
> > +
>> int (*open) (struct inode *inode, struct file *file);
>> ssize_t (*read) (struct cgroup *cgrp, struct cftype *cft,
      struct file *file,
> >
> > Index: cgroup-2.6.25-mm1/kernel/cgroup.c
>> --- cgroup-2.6.25-mm1.orig/kernel/cgroup.c
>> +++ cgroup-2.6.25-mm1/kernel/cgroup.c
>> @ @ -1327,38 +1327,65 @ @ enum cgroup_filetype {
>> FILE RELEASE AGENT.
>> }:
> >
>> -static ssize t cgroup write X64(struct cgroup *cgrp, struct cftype *cft,
>> - struct file *file.
      const char user *userbuf,
      size t nbytes, loff t *unused ppos)
> > +
>> +/**
> > + * cgroup_file_lock(). Helper for cgroup read/write methods.
>> + * @cgrp: the cgroup being acted on
>> + * @cft: the control file being written to or read from
>> + * *write: true if the access is a write access.
```

```
>>+*
> > + * Takes any necessary locks as requested by the control file's
>> + * 'lockmode' field; checks (after locking if necessary) that the
>> + * control group is not in the process of being destroyed.
>> + * Currently all the locking options are implemented in the same way,
>> + * by taking cgroup mutex. Future patches will add finer-grained
> > + * locking.
> > + *
>> + * Calls to cgroup file lock() should always be paired with calls to
>> + * cgroup file unlock(), even if cgroup file lock() returns an error.
> > + */
> > +
>> +static int cgroup_file_lock(struct cgroup *cgrp, struct cftype *cft, int write)
> > - char buffer[64];
>> - int retval = 0:
> > - char *end:
>> + int mask = write ? CFT LOCK FOR WRITE : CFT LOCK FOR READ;
>> + BUILD_BUG_ON(CFT_LOCK_FOR_READ != (CFT_LOCK_FOR_WRITE >> 1));
> >
> > - if (!nbytes)
>> - return -EINVAL;
>> - if (nbytes >= sizeof(buffer))
>> - return -E2BIG;
>> - if (copy_from_user(buffer, userbuf, nbytes))
>> - return -EFAULT;
> > + if (cft->lockmode & mask)
>> + mutex lock(&cgroup mutex);
>> + if (cgroup is removed(cgrp))
>> + return -ENODEV:
> > + return 0;
> > +}
> > +
> > +/**
>> + * caroup file unlock(): undoes the effect of caroup file lock()
> > +
>> +static void cgroup_file_unlock(struct cgroup *cgrp, struct cftype *cft,
          int write)
> > +
> > +{
>> + int mask = write ? CFT_LOCK_FOR_WRITE : CFT_LOCK_FOR_READ;
> > + if (cft->lockmode & mask)
>> + mutex_unlock(&cgroup_mutex);
> > +}
> >
> > - buffer[nbytes] = 0; /* nul-terminate */
>> - strstrip(buffer);
```

```
>> +static ssize_t cgroup_write_X64(struct cgroup *cgrp, struct cftype *cft,
> > +
       const char *buffer)
> > +{
> > + char *end;
>> if (cft->write u64) {
>> u64 val = simple_strtoull(buffer, &end, 0);
>> if (*end)
      return -EINVAL;
>>
>> - retval = cft->write u64(cgrp, cft, val);
>> + return cft->write u64(cgrp, cft, val);
>> } else {
>> s64 val = simple strtoll(buffer, &end, 0);
>> if (*end)
      return -EINVAL;
> >
>> - retval = cft->write_s64(cgrp, cft, val);
>> + return cft->write_s64(cgrp, cft, val);
>> }
> > - if (!retval)
>> - retval = nbytes;
> > - return retval;
>> }
> >
>> static ssize_t cgroup_common_file_write(struct cgroup *cgrp,
>> @ @ -1426,47 +1453,82 @ @ out1:
>> return retval;
>> }
> >
>> -static ssize t cgroup file write(struct file *file, const char user *buf,
>> +static ssize t cgroup file write(struct file *file, const char user *userbuf,
        size_t nbytes, loff_t *ppos)
> >
>> {
>> struct cftype *cft = __d_cft(file->f_dentry);
>> struct cgroup *cgrp = __d_cgrp(file->f_dentry->d_parent);
>> - if (!cft || cgroup_is_removed(cgrp))
>> - return -ENODEV:
>> - if (cft->write)
>> - return cft->write(cgrp, cft, file, buf, nbytes, ppos);
>> - if (cft->write u64 || cft->write s64)
>> - return cgroup write X64(cgrp, cft, file, buf, nbytes, ppos);
>> - if (cft->trigger) {
> > - int ret = cft->trigger(cgrp, (unsigned int)cft->private);
>> - return ret ? ret : nbytes;
> > + ssize_t retval;
> > + char static_buffer[64];
> > + char *buffer = static_buffer;
>> + ssize t max bytes = sizeof(static buffer) - 1;
> > + if (!cft->write && !cft->trigger) {
```

```
> > + if (!nbytes)
>> + return -EINVAL;
>> + if (nbytes >= max_bytes)
>> + return -E2BIG;
>> + if (nbytes >= sizeof(static_buffer)) {
> afaict this can't happen - we would have already returned -E2BIG?
>> + /* +1 for nul-terminator */
>> + buffer = kmalloc(nbytes + 1, GFP KERNEL);
>> + if (buffer == NULL)
>> + return -ENOMEM;
>>+ }
>> + if (copy_from_user(buffer, userbuf, nbytes)) {
>>+ retval = -EFAULT;
>> + goto out_free;
>>+ }
>> + buffer[nbytes] = 0; /* nul-terminate */
>> + strstrip(buffer); /* strip -just- trailing whitespace */
>> - return -EINVAL;
> > -}
> I'm trying to work out what protects static_buffer?
```

One of us must be having a brain cramp because it looks to me like the buffer doesn't need protection -- it's on the stack. It's probably me but I'm just not seeing how this use is unsafe..

> Why does it need to be static anyway? 64 bytes on-stack is OK.

Uh, it is on stack. It doesn't use the C keyword "static". It's just poorly-named.

<snip>

Cheers,

-Matt Helsley

Containers mailing list

Containers@lists.linux-foundation.org

https://lists.linux-foundation.org/mailman/listinfo/containers