## Subject: Re: [RFC] Default child of a cgroup Posted by Balbir Singh on Thu, 31 Jan 2008 18:09:12 GMT

View Forum Message <> Reply to Message

```
Srivatsa Vaddagiri wrote:
```

```
> Hi,
> As we were implementing multiple-hierarchy support for CPU
> controller, we hit some oddities in its implementation, partly related
> to current cgroups implementation. Peter and I have been debating on the
> exact solution and I thought of bringing that discussion to lkml.
>
> Consider the cgroup filesystem structure for managing cpu resource.
>
> # mount -t cgroup -ocpu,cpuacct none /cgroup
> # mkdir /cgroup/A
> # mkdir /cgroup/B
  # mkdir /cgroup/A/a1
> will result in:
>
> /cgroup
    |----<tasks>
    |-----<cpuacct.usage>
>
     |----<cpu.shares>
>
>
>
    |----[A]
        |----<tasks>
>
        |----<cpuacct.usage>
>
        |----<cpu.shares>
>
>
        |---[a1]
>
            |----<tasks>
>
            |----<cpuacct.usage>
>
            |----<cpu.shares>
>
>
>
    |----[B]
>
        |----<tasks>
>
        |----<cpuacct.usage>
>
        |----<cpu.shares>
>
>
>
>
> Here are some questions that arise in this picture:
>
> 1. What is the relationship of the task-group in A/tasks with the
   task-group in A/a1/tasks? In otherwords do they form siblings
```

of the same parent A?

>

I consider them to be the same relationship between directories and files. A/tasks are siblings of A/a1 and A/other children, \*but\* the entities of interest are A and A/a1.

- > 2. Somewhat related to the above question, how much resource should the
- > task-group A/a1/tasks get in relation to A/tasks? Is it 1/2 of parent
- > A's share or 1/(1 + N) of parent A's share (where N = number of tasks
- > in A/tasks)?

>

I propose that it gets 1/2 of the bandwidth, here is why

- 1. Assume that a task in A/tasks forks 1000 children, what happens to the bandwidth of A/a1's tasks then? We have no control over how many tasks can be created on A/tasks as a consequence of moving one task to A/tasks. Doing it the other way would mean, that A/a1/tasks will get 1/1001 of the bandwidth (sounds very unfair and prone to Denial of Service/Fairness)
- > 3. What should A/cpuacct.usage reflect? CPU usage of A/tasks? Or CPU usage
- > of all siblings put together? It can reflect only one, in which case
- > user has to manually derive the other component of the statistics.

>

It should reflect the accumulated usage of A's children and the tasks in A.

- > It seems to me that tasks in A/tasks form what can be called the
- > "default" child group of A, in which case:

>

- > 4. Modifications to A/cpu.shares should affect the parent or its default
- > child group (A/tasks)?

-

- > To avoid these ambiguities, it may be good if cgroup create this
- > "default child group" automatically whenever a cgroup is created?
- > Something like below (not the absence of tasks file in some directories
- > now):

>

I think the concept makes sense, but creating a default child is going to be confusing, as it is not really a child of A.

```
> /cgroup
> |
> |-----<cpuacct.usage>
> |-----<cpu.shares>
```

```
>
      |---[def_child]
>
        |----<tasks>
>
        |----<cpuacct.usage>
>
         |----<cpu.shares>
>
>
>
    |----[A]
>
>
         |----<cpuacct.usage>
>
         |----<cpu.shares>
>
>
        |---[def_child]
>
            |----<tasks>
>
            |----<cpuacct.usage>
>
            |----<cpu.shares>
>
>
>
        |---[a1]
>
>
            |----<cpuacct.usage>
>
            |----<cpu.shares>
>
>
           |---[def_child]
>
                |---<tasks>
>
                |---<cpuacct.usage>
>
                |---<cpu.shares>
>
>
>
    |----[B]
>
>
         |----<cpuacct.usage>
>
         |----<cpu.shares>
>
>
        |---[def_child]
>
            |----<tasks>
>
            |----<cpuacct.usage>
>
            |----<cpu.shares>
>
>
> Note that user cannot create subdirectories under def child with this
> scheme! I am also not sure what impact this will have on other resources
> like cpusets ..
>
```

Which means we'll need special logic in the cgroup filesystem to handle def\_child. Not a very good idea.

> Thoughts?

>

Warm Regards, Balbir Singh Linux Technology Center IBM, ISTL

Containers mailing list Containers@lists.linux-foundation.org https://lists.linux-foundation.org/mailman/listinfo/containers