Subject: Re: [PATCH 1/2] namespaces: introduce sys hijack (v10) Posted by Casey Schauffer on Tue, 27 Nov 2007 22:54:26 GMT

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--- "Serge E. Hallyn" <serue@us.ibm.com> wrote:
> Quoting Stephen Smalley (sds@tycho.nsa.gov):
> > On Tue, 2007-11-27 at 10:11 -0600, Serge E. Hallyn wrote:
>> Quoting Crispin Cowan (crispin@crispincowan.com):
>>> Just the name "sys hijack" makes me concerned.
>>>>
>>> This post describes a bunch of "what", but doesn't tell us about "why"
>>> we would want this. What is it for?
>>> Please see my response to Casey's email.
>>>
>>> And I second Casey's concern about careful management of the privilege
>>> required to "hijack" a process.
>>>
>> Absolutely. We're definately still in RFC territory.
>>> Note that there are currently several proposed (but no upstream) ways to
>> accomplish entering a namespace:
>>>
>>> 1. bind_ns() is a new pair of syscalls proposed by Cedric. An
>>> nsproxy is given an integer id. The id can be used to enter
>>> an nsproxy, basically a straight current->nsproxy = target_nsproxy;
>>>
>>> 2. I had previously posted a patchset on top of the nsproxy
>>> cgroup which allowed entering a nsproxy through the ns cgroup
>>> interface.
> > >
>>> There are objections to both those patchsets because simply switching a
>>> task's nsproxy using a syscall or file write in the middle of running a
>>> binary is quite unsafe. Eric Biederman had suggested using ptrace or
>> something like it to accomplish the goal.
>> Just using ptrace is however not safe either. You are inheriting *all*
>>> of the target's context, so it shouldn't be difficult for a nefarious
>> container/vserver admin to trick the host admin into running something
>>> which gives the container/vserver admin full access to the host.
>> I don't follow the above - with ptrace, you are controlling a process
> > already within the container (hence in theory already limited to its
>> container), and it continues to execute within that container. What's
> > the issue there?
> Hmm, yeah, I may have overspoken - I'm not good at making up exploits
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- > but while I see it possible to confuse the host admin by setting bogus > environment, I guess there may not be an actual exploit. > Still after the fork induced through ptrace, we'll have to execute a > file out of the hijacked process' namespaces and path (unless we get > *really* 'exotic'). With hijack, execution continues under the caller's > control, which I do much prefer. > The remaining advantages of hijack over ptrace (beside "using ptrace for > that is crufty") are > 1. not subject to pid wraparound (when doing hijack cgroup or hijack_ns) > > 2. ability to enter a namespace which has no active processes > These also highlight selinux issues. In the case of hijacking an > empty cgroup, there is no security context (because there is no task) so > the context of 'current' will be used. In the case of hijacking a > populated cgroup, a task is chosen "at random" to be the hijack source. > So there are two ways to look at deciding which context to use. Since
- I wouldn't be surprised if you've been over this a dozen times already, but why hijack an existing process instead of injecting a new one with completely specified attributes? That way you don't distinguish between an empty cgroup and a propulated one and you're not at the mercy of the available hijackees. I know that I would be much less uncomfortable with that schenario.

> mislabeled, so we might want the task in the hijacked task's context.

> control continues in the original acting process' context, we might > want the child to continue in its context. However if the process > creates any objects in the virtual server, we don't want them

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> Sigh. So here's why I thought I'd punt on selinux at least until I had
> a working selinux-enforced container/vserver :)
>
> -serge
> PS: I'm certainly open to the suggestion that the kernel patch in the
> end us as crufty as using ptrace.
>
>>> That's where the hijack idea came from. Yes, I called it hijack to make
>>> sure alarm bells went off:) bc it's definately still worrisome. But at
>>> this point I believe it is the safest solution suggested so far.
> >
>> --
> > Stephen Smalley
> > National Security Agency
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>>

>>-

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