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Subject: Re: LSM and Containers

Posted by [Peter Dolding](#) on Thu, 25 Oct 2007 04:31:11 GMT

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> You're sometimes hard to parse, but here are a few basic facts within  
> which to constrain our discussions:

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> 1. LSMs are a part of the kernel. The entire kernel is in the  
> same trusted computing base

> 2. containers all run on the same kernel

> 3. whether an lsm is compromised, or a tty driver, or anything  
> else which is in the TCB, all containers are compromised

> 4. it is very explicitly NOT a goal to hide from a container

> the fact that it is in a container. So your 'cannot tell how

> deep they are' is not a goal.

You are missing what I am saying.

I am breaking the LSM in two. Kernel level comprise that is out of my hands.

If you wished in mine the controller could be user space run outside kernel does not have to stay in the TCB. Since its not having any direct control on the approval or rejection of access. Its only queried when needed like when a process breaches the rules applied to it or when processes start or changing security. So breaching it turning it off and so on. This is to reduce the risk of someone trying to customize a LSM can cause a security flaw of hell level(Inside kernel space).

With Intel's and Amd new memory table splitting yes the controller could be put away from the main kernel in time with the rest of the container only data. Yes less bits trusted. This includes when sending the container between machines the controller with it current state could be sent with it. If the interface to controller is standard of course. Ok apparmor and selinux will cope with your system but some of the more complex state based need to transfer state as well this should be in the controller bit of a LSM and not mixed up with the state of the overall machine either.

Over all this is making it simpler to do advanced container things like sending between servers.

The permission enforcement parts stay the same no matter the controller in use. Breaching the enforcement parts still remain a problem. Allows running many controllers unlike LSM's where they can fight and on way around that cleanly. Just like LSM two controllers controlling the same space would still fight so guess how many each container take 1. Since a different controller has to be give its own zone. So people cannot complain about not being able to run there

preferred controller if they want a different one they have to give it a different zone. Yes this is getting on top of a lot of long term problems of people wanting a or b LSM. Stuff it have both and be happy.

Its also removing the single overall kill switch. And replacing them with a per container kill switches. So someone does what will normally turn X controller off and it does. Yet still the outside security is applied to everything run in the container. Security is lowered but not off.

The controllers don't need to keep track of how deep they are. They just don't need to know. From there point of view the world ends from there startup security settings. Only time the controller would have to be give notice is if those outside security settings were changed so it could update the programs under it to the way it wanted. Even that the security settings if they were lower would have just been brute forced applied. Basically its have you done containers running containers is you model Serge E. Hallyn. This avoids have to much around with flags for the next container in. Only thing in mine are you doing is changing the path back to the controller from the security enforcing modules. Enforcing modules process not doing permissions were told to allow I report that to x security container that reports it to contoller. Note the security container only hands out what its allowed nothing more the controller is not trusted. No processing to work out what security files it should be using since the controller only knows the files it should be using.

Yes there is overhead in my system. But were able the system is avoiding traveling in a perfectly functioning system controller could be only being bothered when new processors were being created the rest of the time the enforcement modules just do there job as they were configured.

Traveling down a tree of stacked LSM is going to cause massive lag. Ie knowing that something is below you. Not knowing reduces travel and time to resolve. And allows security to be controlled over the complete container from one spot. The problem with running many LSM again having to tweek them all to get it locked down on container. Ie enabling some disabling others so simple to loss track and have big problems. No problem here the security model is applied to the container and if you want to lock its outside down you just do and the controller inside has to cope. Yes it will be possible to take too much away from a security container. Ok this security container should not have access to that device slam not a question. With running many LSM the question what one does that container own to so I do turn that device off and not have another LSM turn it back on.

Besides having selinux and apparmor installed side by side is going to lag the system due to overlapping hooks. This what I am avoiding. The enforcement modules are only hooked in once no reason at all to overlap them. This overlapping is what makes running many LSM problematic.

The outside security on a container will stay on even if the host LSM is turned off in the design I am look at. It was set when the container started and would have to be directly updated. The outside security defined to a container would be everything selinux could do to a program and then some. Just like selinux opt in.

Over lapping I am handling differently in this design. The outside security set on the container covers overlaps. Far more flex able option. Using directory filtering and other options are on the table as well as far more complex cap options. My design depends on the common protective modules. Not the bit that says this is Mac or Role based or State or so on.

Basically apparmor and selinux would be sitting on the same engine that everyone can use to build solid controllers. If you want a new enforcement feature you would add it for everyone. Since altering the controller would be toothless only can tweak what already exists since you can never be sure its even in kernel space no direct tweaking no hooking no bad stuff with controllers.

There is a security advantage from the design change even for people not using containers or LSM. Everything in system can apply LSM style restrictions lower than what it has to what it wants to run even if a controller is present or not. The controller is to provide the security model in use. With the enforcement system you could almost do any security model with user space code. The difference being not catching process starts to apply security to them. This is useful for programs running untrusted content web browsers that java program running from web does not need all the access I am cut it back. Application level security something LSM overlook and it not wise to use different interfaces for it.

Yes these are major changes in design. Yes they will force clear designs and more sharing so everyone can take advantage of advancements. As well should kill off lots of problems of poor quality LSM's and LSM limitations. Basically my design and LSM cannot live side by side effectively.

Peter Dolding

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