Subject: Re: [RFC] cpuset update_cgroup_cpus_allowed Posted by Paul Jackson on Tue, 16 Oct 2007 02:34:39 GMT

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> currently against an older kernel

ah .. which older kernel?

I tried it against the broken out 2.6.23-rc8-mm2 patch set, inserting it before the task-containersv11-* patches, but that blew up on me - three rejected hunks.

Any chance of getting this against a current cgroup (aka container) kernel?

Could you use the diff --show-c-function option when composing patches - they're easier to read that way - thanks.

```
+ if (!retval) {
+ cpus_allowed = cpuset_cpus_allowed(p);
+ if (!cpus subset(new mask, cpus allowed)) {
+ /*
   * We must have raced with a concurrent cpuset
   * update. Just reset the cpus_allowed to the
   * cpuset's cpus_allowed
   */
   new_mask = cpus_allowed;
```

This narrows the race, perhaps sufficiently, but I don't see that it guarantees closure. Memory accesses to two different locations are not guaranteed to be ordered across nodes, as best I recall. The second line above, that rereads the cpuset cpus_allowed, could get an old value, in essence.

cpuset update task sched_setaffinity task

A. write cpuset [Q] V. read cpuset [Q]

B. read task [P] W. check ok

C. write task [P] X. write task [P]

Y. reread cpuset [Q]

Z. check ok again

Two memory locations:

[P] the cpus_allowed mask in the task_struct of the task doing the sched_setaffinity call.

[Q] the cpus allowed mask in the cpuset of the cpuset to which the sched setaffinity task is attached.

Even though, from the perspective of location [P], both B. and C. happened before X., still from the perspective of location [Q] the rereading in Y. could return the value the cpuset cpus_allowed had before the write in A. This could result in a task running with a cpus_allowed that was totally outside its cpusets cpus_allowed.

I will grant that this is a narrow window. I won't loose much sleep over it.

> - uses a priority heap to pick the processes to act on, based on start time

This adds a fair bit of code and complexity, relative to my patch. This I do loose more sleep over. There has to be a compelling reason for doing this.

The point that David raises, regarding the interaction of this with hotplug, seems to be a compelling reason for doing -somethingdifferent than my patch proposal.

I don't know yet if it compels us to this much code, however.

Any chance you could provide a patch that works against cgroups?

I won't rest till it's the best ... Programmer, Linux Scalability Paul Jackson <pj@sqi.com> 1.925.600.0401

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