## Subject: Re: [PATCH] namespaces: introduce sys\_hijack (v4) Posted by Cedric Le Goater on Fri, 12 Oct 2007 09:30:33 GMT

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Serge E. Hallyn wrote:
> Quoting Serge E. Hallyn (serue@us.ibm.com):
>> >From 945fe66259cd0cfdc2fe846287b7821e329a558c Mon Sep 17 00:00:00 2001
>> From: sergeh@us.ibm.com <hallyn@kernel.(none)>
>> Date: Tue, 9 Oct 2007 08:30:30 -0700
>> Subject: [PATCH] namespaces: introduce sys hijack (v4)
>>
>> Move most of do fork() into a new do fork task() which acts on
>> a new argument, task, rather than on current. do_fork() becomes
>> a call to do_fork_task(current, ...).
>>
>> Introduce sys_hijack (for x86 only so far). It is like clone, but
>> in place of a stack pointer (which is assumed null) it accepts a
>> pid. The process identified by that pid is the one which is
>> actually cloned. Some state - include the file table, the signals
>> and sighand (and hence tty), and the ->parent are taken from the
>> calling process.
>>
>> The effect is a sort of namespace enter. The following program
>> uses sys hijack to 'enter' all namespaces of the specified pid.
>> For instance in one terminal, do
>>
>> mount -t cgroup -ons /cgroup
>> hostname
    gemu
>>
>> ns exec -u /bin/sh
    hostname serge
         echo $$
>>
          1073
>>
    cat /proc/$$/cgroup
>>
      ns:/node 1073
>>
>> In another terminal then do
>>
>> hostname
    qemu
>> cat /proc/$$/cgroup
    ns:/
>> hijack 1073
    hostname
>>
      serge
>>
    cat /proc/$$/cgroup
>>
      ns:/node 1073
>>
>>
```

>> sys\_hijack is arch-dependent and is only implemented for i386 so far.