
Subject: Re: [PATCH 1/3] Signal semantics for /sbin/init
Posted by [Sukadev Bhattiprolu](#) on Mon, 01 Oct 2007 17:47:20 GMT
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Serge E. Hallyn [serue@us.ibm.com] wrote:
| Quoting sukadev@us.ibm.com (sukadev@us.ibm.com):
| > Oleg Nesterov [oleg@tv-sign.ru] wrote:
| > | On 09/13, sukadev@us.ibm.com wrote:
| > | >
| > | > Oleg Nesterov [oleg@tv-sign.ru] wrote:
| > | > | >
| > | > | >> Notes:
| > | > | >>
| > | > | >> - Blocked signals are never ignored, so init still can receive
| > | > | >> a pending blocked signal after sigprocmask(SIG_UNBLOCK).
| > | > | >> Easy to fix, but probably we can ignore this issue.
| > | > | >
| > | > | > I was wrong. This should be fixed right now. I _think_ this is easy,
| > | > | > and I was going to finish this patch yesterday, but - sorry! - I just
| > | > | > can't switch to "kernel mode" these days, I am fighting with some urgent
| > | > | > tasks on my paid job.
| > | > | >
| > | > | > To respect the current init semantic,
| > | > |
| > | > | The current init semantic is broken in many ways ;)
| > | > |
| > | > | > shouldn't we discard any unblockable
| > | > | > signal (STOP and KILL) sent by a process to its pid namespace init process ?
| > | > |
| > | > | Yes. And Patch 1/3 (Oleg's patch) in the set I sent, handles this already
| > | > | (since STOP and KILL are never in the task->blocked list)
| > | > |
| > | > |
| > | > | > Then, all other signals should be handled appropriately by the pid namespace
| > | > | > init.
| > | > |
| > | > |
| > | > | Yes, I think you are probably right, this should be enough in practice. After all,
| > | > | only root can send the signal to /sbin/init.
| > | > |
| > | > | I agree - the assumption that the container-init will handle these
| > | > | other signals, simplifies the kernel implementation for now.
| > | > |
| > | > |
| > | > | On my machine, /proc/1/status shows that init doesn't have a handler for
| > | > | non-ignored SIGUNUSED == 31, though.
| > | > |
| > | > | But who knows? The kernel promises some guarantees, it is not good to break them.

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| > | > | Perhaps some strange non-standard environment may suffer.
| > | > |
| > | > | > We are assuming that the pid namespace init is not doing anything silly and
| > | > | > I guess it's OK if the consequences are only on the its pid namespace and
| > | > | > not the whole system.
| > | > |
| > | > | The sub-namespace case is very easy afaics, we only need the "signal comes from
| > | > | the parent namespace" check, not a problem if we make the decision on the sender's
| > | > | path, like this patch does.
| > | > |
| > | > | Yes, patches 2 and 3 of the set already do the ancestor-ns check. no ?
| > | > |
| > | Yes, I think patches 2-3 are good. But this patch is not. I thought that we
| > | can ignore the "Blocked signals are never ignored" problem, now I am not sure.
| > | It is possible that init temporary blocks a signal which it is not going to
| > | handle.
| > |
| > | Perhaps we can do something like the patch below, but I don't like it. With
| > | this patch, we check the signal handler even if /sbin/init blocks the signal.
| > | This makes the semantics a bit strange for /sbin/init. Hopefully not a problem
| > | in practice, but still not good.
| > |
| > | I think this is one step ahead of what we were discussing last week.
| > | A container-init that does not have a handler for a fatal signal would
| > | survive even if the signal is posted when it is blocked.
| > |
| > |
| > | Unfortunately, I don't know how to make it better. The problem with blocked
| > | signals is that we don't know who is the sender of the signal at the time
| > | when the signal is unblocked.
| > |
| > |
| > | One solution I was thinking of was to possibly queue pending blocked
| > | signals to a container init separately and then requeue them on the
| > | normal queue when signals are unblocked. Its definitely not an easier
| > | solution, but might be less intrusive than the "signal from parent ns
| > | flag" solution.
| > |
| I personally prefer the flag solution just because it will remain
| clear why it is there, whereas understanding why the separate queue
| is there will be harder unless it is named something like
| "child_contaienr_init_blocked_pending_queue".

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In my first version of the "flag" solution, I stored the flag in the sigqueue structure. The problem with that approach was that if the allocation of the sigqueue failed, we would not know if the sender was in parent ns. Note that we post a signal to the process (add to signals->signal set) even if this allocation fails.

By storing the signal info in 'struct pid_namespace' we avoid having to allocate when posting the signal.

I agree that a descriptive name is needed. but since the fields are in 'struct pid_namespace' I was thinking 'child' was not necessary. Maybe 'cinit' instead of 'container init'. Also 'pending' somehow implies a 'queue' - no ?

| But still it may be the way to go. Have you coded up a version of this?

I was playing with a slightly different solution that I could modify for this. But I can code that up in a couple of days. Just wanted to see if it was interesting approach at all.

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|  
| thanks,  
| -serge  
|  
| > i.e suppose we have:  
| >  
| > struct pid_namespace {  
| > ...  
| > struct sigpending cinit_blocked_pending;  
| > struct sigpending cinit_blocked_shared_pending;  
| > }  
| >  
| > Signals from ancestor ns are queued as usual on task->pending and  
| > task->signal->shared_pending. They don't need any special handling.  
| >  
| > Only signals posted to a container-init from within its namespace  
| > need special handling (as in: ignore unhandled fatal signals from  
| > same namespace).  
| >  
| > If the container-init has say SIGUSR1 blocked, and a descendant of  
| > container-init posts SIGUSR1 to container-init, queue the SIGUSR1  
| > in pid_namespace->cinit_blocked_pending.  
| >  
| > When container-init unblocks SIGUSR1, check if there was a pending  
| > SIGUSR1 from same namespace (i.e check ->cinit_blocked_pending list).  
| > If there was and container-init has a handler for SIGUSR1, post SIGUSR1  
| > on task->pending queue and let the container-init handle SIGUSR1.  
| >  
| > If there was a SIGUSR1 posted to container init and there is no handler  
| > for SIGUSR1, then just ignore the SIGUSR1 (since it would be fatal  
| > otherwise).  
| >  
| > I chose 'struct pid_namespace' for the temporary queue, since we need
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| > the temporary queues only for container-inits (not for all processes).
| > And having it allocated ahead of time, ensures we can queue the signal
| > even under low-memory conditions.
| >
| > Just an idea at this point.
| >
| > |
| > | What do you think? Can we live with this oddity? Otherwise, we have to add
| > | something like the "the signal is from the parent namespace" flag, and I bet
| > | this is not trivial to implement correctly.
| >
| > I think its reasonable to place some restrictions on container-init
| > processes, so, yes, I think the oddity is fine for now (i.e at least
| > until someone needs a different behavior).
| >
| > BTW, I ran some tests on this patch and they seem to work as expected :)
| > Will run some more tests today.
| >
| > |
| > | Oleg.
| > |
| > | --- t/kernel/signal.c~IINITSIGS 2007-08-28 19:15:28.000000000 +0400
| > | +++ t/kernel/signal.c 2007-09-17 19:20:24.000000000 +0400
| > | @@ -39,11 +39,35 @@
| > |
| > | static struct kmem_cache *sigqueue_cachep;
| > |
| > | +static int sig_init_ignore(struct task_struct *tsk)
| > | +{
| > | + // Currently this check is a bit racy with exec(),
| > | + // we can _simplify_ de_thread and close the race.
| > | + if (likely(!is_init(tsk->group_leader)))
| > | + return 0;
| > |
| > | -static int sig_ignored(struct task_struct *t, int sig)
| > | + // ----- Multiple pid namespaces -----
| > | + // if (current is from tsk's parent pid_ns && !in_interrupt())
| > | + // return 0;
| > | +
| > | + return 1;
| > | +}
| > | +
| > | +static int sig_task_ignore(struct task_struct *tsk, int sig)
| > | {
| > | - void __user * handler;
| > | + void __user * handler = tsk->sighand->action[sig-1].sa.sa_handler;
| > | +
| > | + if (handler == SIG_IGN)

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|> | + return 1;
|> | +
|> | + if (handler != SIG_DFL)
|> | + return 0;
|> |
|> | + return sig_kernel_ignore(sig) || sig_init_ignore(tsk);
|> | +}
|> | +
|> | +static int sig_ignored(struct task_struct *t, int sig)
|> | +{
|> | /*
|> |  * Tracers always want to know about signals..
|> |  */
|> | @@ -55,13 +79,10 @@ static int sig_ignored(struct task_struct
|> |  * signal handler may change by the time it is
|> |  * unblocked.
|> |  */
|> | - if (sigismember(&t->blocked, sig))
|> | + if (sigismember(&t->blocked, sig) && !sig_init_ignore(t))
|> |     return 0;
|> |
|> | - /* Is it explicitly or implicitly ignored? */
|> | - handler = t->sighand->action[sig-1].sa.sa_handler;
|> | - return handler == SIG_IGN ||
|> | - (handler == SIG_DFL && sig_kernel_ignore(sig));
|> | + return sig_task_ignore(t, sig);
|> | }
|> |
|> | /*
|> | @@ -554,6 +575,9 @@ static void handle_stop_signal(int sig,
|> |  */
|> |     return;
|> |
|> | + if (sig_init_ignore(p))
|> | + return;
|> | +
|> | if (sig_kernel_stop(sig)) {
|> |     /*
|> |      * This is a stop signal. Remove SIGCONT from all queues.
|> |      */
|> |     @@ -1822,14 +1846,6 @@ relock:
|> |     if (sig_kernel_ignore(signr)) /* Default is nothing. */
|> |         continue;
|> |
|> | - /*
|> | -  * Init of a pid space gets no signals it doesn't want from
|> | -  * within that pid space. It can of course get signals from
|> | -  * its parent pid space.
|> | -  */

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```
|> | - if (current == child_reaper(current))
|> | - continue;
|> | -
|> | if (sig_kernel_stop(signr)) {
|> |   if (current->signal->flags & SIGNAL_GROUP_EXIT)
|> |     continue;
|> | @@ -2308,8 +2324,7 @@ int do_sigaction(int sig, struct k_sigac
|> |   * (for example, SIGCHLD), shall cause the pending signal to
|> |   * be discarded, whether or not it is blocked"
|> |   */
|> | - if (act->sa.sa_handler == SIG_IGN ||
|> |   (act->sa.sa_handler == SIG_DFL && sig_kernel_ignore(sig))) {
|> | + if (sig_task_ignore(current, sig)) {
|> |   struct task_struct *t = current;
|> |   sigemptyset(&mask);
|> |   sigaddset(&mask, sig);
|> |
|> | _____
|> Containers mailing list
|> Containers@lists.linux-foundation.org
|> https://lists.linux-foundation.org/mailman/listinfo/containers
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