Subject: Re: [PATCH 0/8] RSS controller based on process containers (v3.1) Posted by Herbert Poetzl on Fri, 08 Jun 2007 12:03:43 GMT

View Forum Message <> Reply to Message

On Mon, Jun 04, 2007 at 05:25:25PM +0400, Pavel Emelianov wrote:

- > Adds RSS accounting and control within a container.
- > Changes from v3
- > comments across the code
- > git-bisect safe split
- > lost places to move the page between active/inactive lists

> Ported above Paul's containers V10 with fixes from Balbir.

- > RSS container includes the per-container RSS accounting
- > and reclamation, and out-of-memory killer.

>

>

>

- > Each mapped page has an owning container and is linked into its
- > LRU lists just like in the global LRU ones. The owner of the page
- > is the container that touched the page first.
- > As long as the page stays mapped it holds the container, is accounted
- > into its usage and lives in its LRU list. When page is unmapped for
- > the last time it releases the container.
- > The RSS usage is exactly the number of pages in its booth LRU lists,
- > i.e. the nu,ber of pages used by this container.

so there could be two guests, unified (i.e. sharing most of the files as hardlinks), where the first one holds 80% of the resulting pages, and the second one 20%, and thus shows much lower 'RSS' usage as the other one, although it is running the very same processes and providing identical services?

- > When this usage exceeds the limit set some pages are reclaimed from
- > the owning container. In case no reclamation possible the OOM killer
- > starts thinning out the container.

so the system (physical machine) starts reclaiming and probably swapping even when there is no need to do so?

e.g. a system with a single guest, limited to 10k pages, with a working set of 15k pages in different apps would continuously swap (trash?) on an otherwise unused (100k+ pages) system?

- > Thus the container behaves like a standalone machine when it runs
- > out of resources, it tries to reclaim some pages, and if it doesn't
- > succeed, kills some task.

is that really what we want?
I think we can do _better_ than a standalone machine and in many cases we really should ...

```
best,
Herbert
```

```
> Signed-off-by: Pavel Emelianov <xemul@openvz.org>
> The testing scenario may look like this:
> 1. Prepare the containers
> # mkdir -p /containers/rss
> # mount -t container none /containers/rss -o rss
> 2. Make the new group and move bash into it
> # mkdir /containers/rss/0
> # echo $$ > /containers/rss/0/tasks
> Since now we're in the 0 container.
> We can alter the RSS limit
> # echo -n 6000 > /containers/rss/0/rss_limit
> We can check the usage
> # cat /containers/rss/0/rss_usage
> 25
>
> And do other stuff. To check the reclamation to work we need a
> simple program that touches many pages of memory, like this:
>
> #include <stdio.h>
> #include <unistd.h>
> #include <sys/mman.h>
> #include <fcntl.h>
> #ifndef PGSIZE
> #define PGSIZE 4096
> #endif
>
> int main(int argc, char **argv)
> {
       unsigned long pages;
>
       int i;
```

```
char *mem;
>
>
       if (argc < 2) {
>
            printf("Usage: %s <number_of_pages>\n", argv[0]);
>
            return 1;
>
>
      }
>
       pages = strtol(argv[1], &mem, 10);
>
       if (*mem != '\0') {
>
           printf("Bad number %s\n", argv[1]);
>
>
            return 1;
      }
>
>
       mem = mmap(NULL, pages * PGSIZE, PROT_READ | PROT_WRITE,
>
                MAP_PRIVATE | MAP_ANON, 0, 0);
>
      if (mem == MAP_FAILED) {
>
            perror("map");
>
            return 2;
>
      }
>
>
       for (i = 0; i < pages; i++)
>
            mem[i * PGSIZE] = 0;
>
>
       printf("OK\n");
>
       return 0;
>
> }
>
> Containers mailing list
> Containers@lists.linux-foundation.org
> https://lists.linux-foundation.org/mailman/listinfo/containers
Containers mailing list
Containers@lists.linux-foundation.org
```

https://lists.linux-foundation.org/mailman/listinfo/containers