Subject: Re: [ckrm-tech] [PATCH 0/2] resource control file system - aka containers on top of nsproxy!

Posted by Srivatsa Vaddagiri on Mon, 05 Mar 2007 17:34:01 GMT

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On Sat, Mar 03, 2007 at 06:32:44PM +0100, Herbert Poetzl wrote:

- > Yes, perhaps this overloads nsproxy more than what it was intended for.
- > > But, then if we have to to support resource management of each
- > > container/vserver (or whatever group is represented by nsproxy),
- >> then nsproxy seems the best place to store this resource control
- > > information for a container.

>

- > well, the thing is, as nsproxy is working now, you
- > will get a new one (with a changed subset of entries)
- > every time a task does a clone() with one of the
- > space flags set, which means, that you will end up
- > with quite a lot of them, but resource limits have
- > to address a group of them, not a single nsproxy
- > (or act in a deeply hierarchical way which is not
- > there atm, and probably will never be, as it simply
- > adds too much overhead)

Thats why nsproxy has pointers to resource control objects, rather than embedding resource control information in nsproxy itself.

>From the patches:

```
struct nsproxy {
+#ifdef CONFIG_RCFS
+ struct list_head list;
+ void *ctlr_data[CONFIG_MAX_RC_SUBSYS];
+#endif
}
```

This will let different nsproxy structures share the same resource control objects (ctlr\_data) and thus be governed by the same parameters.

Where else do you think the resource control information for a container should be stored?

- > > It should have the same perf overhead as the original
- > > container patches (basically a double dereference -
- > > task->containers/nsproxy->cpuset required to get to the
- > > cpuset from a task).

>

> on every limit accounting or check? I think that

> is quite a lot of overhead ...

tsk->nsproxy->ctlr\_data[cpu\_ctlr->id]->limit (4 dereferences) is what we need to get to the cpu b/w limit for a task.

If cpu\_ctlr->id is compile time decided, then that would reduce it to 3.

But I think if CPU scheduler schedules tasks from same container one after another (to the extent possible that is), then other derefences (->ctlr\_data[] and ->limit) should be fast, as they should be in the cache?

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Regards, vatsa

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