Subject: Re: [RFC] L3 network isolation: broadcast Posted by Daniel Lezcano on Wed, 13 Dec 2006 23:08:11 GMT

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Vlad Yasevich wrote:
> Daniel Lezcano wrote:
>> Hi all,
>>
>> I am trying to find a solution to handle the broadcast traffic on the I3
>> namespace.
>>
>> The broadcast issue comes from the I2 isolation:
>>
>> in udp.c
>>
>> static inline struct sock *udp_v4_mcast_next(struct sock *sk,
        be16 loc port.
      __be32 loc_addr,
>>
        be16 rmt port,
>>
        _be32 rmt_addr,
>>
>>
      int dif)
>> {
>> struct hlist_node *node;
>> struct sock *s = sk;
>> struct net_namespace *ns = current_net_ns;
>> unsigned short hnum = ntohs(loc_port);
>>
>> sk for each from(s, node) {
    struct inet sock *inet = inet sk(s);
>>
    if (inet->num != hnum
                             Ш
>>
       (inet->daddr && inet->daddr != rmt_addr) ||
>>
       (inet->dport != rmt_port && inet->dport) ||
>>
       (inet->rcv_saddr && inet->rcv_saddr != loc_addr) ||
>>
       ipv6_only_sock(s)
>>
       !net_ns_match(sk->sk_net_ns, ns) ||
>>
       (s->sk_bound_dev_if && s->sk_bound_dev_if != dif))
>>
    continue:
>>
    if (!ip_mc_sf_allow(s, loc_addr, rmt_addr, dif))
    continue;
    goto found;
>>
>>
     }
>> s = NULL;
>> found:
     return s;
>>
>> }
>>
>> This is absolutely correct for I2 namespaces because they share the
```

- >> socket hash table. But that is not correct for I3 namespaces because we
- >> want to deliver the packet to each I3 namespaces which have binded to
- >> the broadcast address, so we should avoid checking net_ns_match if we
- >> are in a layer 3 namespace. Doing that we will break the I2 isolation
- >> because an another I2 namespace could have binded to the same broadcast
- >> address.

>

- > A question, if you will... I am still digesting the I2 changes, and I can't
- > remember/find if the broadcasts will be replicated across multiple I2 or not.

Well ... I am not sure (never tested it) but as far as I remember, it is the bridge which should duplicate the packets because it acts as a "hub".

When a packet is received on eth0, it is forwarded to br0 (the bridge) and this one will send the packet to veth0, veth1 and veth2. The packets will follow the normal incoming path for each namespace. So I think the answer is yes, the broadcast is replicated to each I2 namespace.

Dmitry can give more information on that I think.

- > Example:
- > A system has 2 interfaces eth0 and eth1 connected to the same lan/link.
- > Each NIC was isolated to it's own L2 space. Each L2 space configures
- > the its nic with unique IP but in the same subnet. Will both L2s receive
- > a subnet broadcast packet?

Depending on the bridge configuration, I am inclined to say yes if eth0 and eth1 are attached to the bridge, no if they are not attached.

Not attached

Attached

```
-- br0 --
              ---- veth1--|ns l2]--eth0
But again, I am not sure.
> If yes, then below approach will work. If no, then we'll need something else
> since both L2s should get the packet in their own right.
It is a critical path for broadcast and multicast incoming traffic,
should I implement this approach and we try to optimize that later?
>> The solution I see here is:
>>
>> if namespace is I3 then;
>> net_ns match any net_ns registered as listening on this address
>> else
>> net_ns_match
>> fi
>>
>> The registered network namespace is a list shared between brothers I3
>> namespaces. This will add more overhead for sure. Does anyone have
>> comments on that or perhaps a better solution?
>
> -vlad
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