Subject: Re: [ckrm-tech] [PATCH 1/7] containers (V7): Generic container system abstracted from cpusets code Posted by Srivatsa Vaddagiri on Sun, 25 Mar 2007 02:21:09 GMT

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On Sat, Mar 24, 2007 at 06:41:28PM -0700, Paul Jackson wrote:
> > the following code becomes racy with cpuset_exit() ...
> >
        atomic inc(&cs->count);
> >
        rcu assign pointer(tsk->cpuset, cs);
> >
> >
        task_unlock(tsk);
> eh ... so ... ?
> I don't know of any sequence where that causes any problem.
> Do you see one?
Let's say we had two cpusets CS1 amd CS2 (both different from top cpuset).
CS1 has just one task T1 in it (CS1->count = 0) while CS2 has no tasks
in it (CS2->count = 0).
Now consider:
CPU0 (attach_task T1 to CS2) CPU1 (T1 is exiting)
task lock(T1);
oldcs = tsk->cpuset;
[oldcs = CS1]
T1->flags & PF_EXITING? (No)
   T1->flags = PF_EXITING;
atomic_inc(&CS2->count);
   cpuset exit()
      cs = tsk->cpuset; (cs = CS1)
T1->cpuset = CS2;
      T1->cpuset = &top_cpuset;
task unlock(T1);
```

CS2 has one bogus count now (with no tasks in it), which may prevent it from being removed/freed forever.

Not just this, continuing further we have more trouble: CPU0 (attach task T1 to CS2) CPU1 (T1 is exiting) synchronize_rcu() atomic_dec(&CS1->count); [CS1->count=0]if atomic_dec_and_test(&oldcs->count)) [CS1->count = -1]We now have CS1->count negative. Is that good? I am uncomfortable .. We need a task_lock() in cpuset_exit to avoid this race. Regards, vatsa